



CÓRDOBA UNIVERSITY

SUPERIOR POLYTECHNIC SCHOOL

DEPARTMENT OF  
COMPUTER SCIENCE AND NUMERICAL ANALYSIS



# DECLARATIVE PROGRAMMING

COMPUTER ENGINEERING  
COMPUTATION ESPECIALITY

FOURTH YEAR

FIRST FOUR-MONTH PERIOD



**Subject 1.- Introduction to Scheme language**

**First part:**  
**Scheme**

**Subject 1.- Introduction to Scheme language**

**Subject 2.- Expressions and Functions**

**Subject 3.- Conditional Predicates and Sentences**

**Subject 4.- Iteration and Recursion**

**Subject 5.- Compound Data Types**

**Subject 6.- Data Abstraction**

**Subject 7.- Reading and Writing**

**Second part:**  
**Prolog**

**Subject 8.- Introduction to Prolog language**

**Subject 9.- Basic Elements of Prolog**

**Subject 10.- Lists**

**Subject 11.- Re-evaluation and the "cut"**

**Subject 12.- Input and Output**

First part: **Scheme**

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## Contents

1. Fundamental Characteristics of Functional Programming
2. Historic Summary of Scheme

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1. Fundamental Characteristics of Functional Programming
2. Historic Summary of Scheme

## 1. **Fundamental Characteristics of Functional Programming**

- ✓ **Functional** Programming is a **subtype** of **Declarative** Programming

# 1. Fundamental Characteristics of Functional Programming

## ✓ Declarative Programming (1 / 2)

### ➤ Objective: Problem description

**“What”** problem must be resolved?

#### ▪ Notice:

- It does **not** mind **“how”** the problem is resolved
- It **avoids** the implementation features.

# 1. Fundamental Characteristics of Functional Programming

## ✓ Declarative Programming (2 / 2)

### ➤ Features

- Expressivity
- Extensible: 10% - 90% rule
- Protection
- Mathematic Elegance

### ➤ Types:

- **Functional** *or* Applicative Programming:
  - Lisp, **Scheme**, Haskell, ...
- **Logic** Programming: **Prolog**

## 1. Fundamental Characteristics of Functional Programming

### ✓ Principle of the **"Pure"** Functional Programming

*"The **expression value** only **depends on its sub-expressions values**, if such sub-expressions exist".*

### ✓ **Non collateral effects**

The value of **"a + b"** **only** depends on **"a"** and **"b"**.

### ✓ The **function** term is used in its **mathematical** sense.

### ✓ **No instructions**: programming **without** assignments

➤ The **impure** Functional programming **allows** the

**"assignment instruction"**

# 1. Fundamental Characteristics of Functional Programming

- ✓ **Program structure in Functional Programming**
  - The **program** is a function **composed** of simpler functions
  - **Function execution:**
    - **Receives the input data:** functions arguments or parameters
    - **Evaluates the expressions**
    - **Returns the Result:** computed value of the function

# 1. Fundamental Characteristics of Functional Programming

## ✓ Type of Functional Languages

- Most of them are **interpreted** languages
- Some of them have **compiled** versions

## ✓ Memory management

### ➤ Implicit memory management:

- Memory management is a task of the interpreter.
- The programmer must **not** worry about memory management.

### ➤ **Garbage collection**: task of the interpreter.

**In short**: the programmer must only worry about the **Problem description**

## Contents

1. Fundamental Characteristics of Functional Programming
2. **Historic Summary of Scheme**

## 2. **Historic Summary of Scheme**

- ✓ LISP
- ✓ Compilation versus Interpretation
- ✓ Lexical (or static) versus dynamical scope
- ✓ Origin of Scheme

## 2. **Historic Summary of Scheme**

- ✓ **LISP**
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## 2. **Historic Summary of Scheme**

### ✓ **LISP**

- **John McCarthy** (MIT)
- **“Advice Taker”** program:
  - Theoretical basis: Logic Mathematics
  - Objective: Deduction and Inferences
- **LISP: LISt Processing** (1956 - 1958)
  - Second historic language of **Artificial Intelligence** (after IPL)
  - At present time, second historic language **in use** (after Fortran)
  - LISP is based on Lambda Calculus (**Alonzo Church**)
- **Scheme** is a **dialect** of **LISP**

## 2. Historic Summary of Scheme

### ✓ LISP

#### ➤ Functional Programming Characteristics

- **Recursion**
- **Lists**
- **Implicit** memory management
- Interactive and **interpreted** programs
- **Symbolic** Programming
- **Dynamically** scoped for **non** local variables

## 2. **Historic Summary of Scheme**

### ✓ **LISP**

- **LISP's contributions:**
  - **Built - in functions**
  - **Garbage collection**
  - **Definition Formal Language: **LISP** itself**

## 2. Historic Summary of Scheme

### ✓ LISP

#### ➤ Applications: **Artificial Intelligence** Programs

- Theorem verification and testing
- Symbolic differentiation and integration
- Search Problems
- Natural Language Processing
- Computer Vision
- Robotics
- Knowledge Representation Systems
- Expert Systems
- **And so on**

## 2. **Historic Summary of Scheme**

### ✓ **LISP**

#### ➤ **Dialects (1 /2)**

- **Mac LISP** (Man and computer or Machine - aided cognition): **East** Coast Version
- **Inter LISP** (Interactive LISP): **West** Coast Version
  - Bolt, Beranek y Newman Company (BBN)
  - Research Center of Xerox at Palo Alto (Texas)
  - **LISP Machine**

## 2. Historic Summary of Scheme

### ✓ LISP

#### ➤ Dialects (2 / 2)

- **Mac LISP** (Man and computer or Machine - aided cognition): East Coast Version
  - C-LISP: Massachusetts University
  - Franz LISP: California University (Berkeley).  
**Compiled version.**
  - NIL (New implementation of LISP): MIT.
  - PSL (Portable Standard LISP): Utah University
  - **Scheme**: MIT.
  - T (True): Yale University.
  - Common LISP

## 2. **Historic Summary of Scheme**

- ✓ LISP
- ✓ **Compilation versus Interpretation**
- ✓ Lexical (or static) versus dynamical scope
- ✓ Origin of Scheme

## 2. Historic Summary of Scheme

### ✓ **Compilation versus interpretation**

#### ➤ **Compilation:**

- The **source code** (high level) is **transformed** into **executable code** (low level), which can be independently run.

## 2. **Historic Summary of Scheme**

- ✓ **Compilation versus interpretation**

- **Compilation**

**Source code** →

**Compiler**

## 2. Historic Summary of Scheme

### ✓ **Compilation** versus **interpretation**

#### ➤ **Compilation**

**Source code** →

**Compiler**



**Compilation  
errors**

## 2. Historic Summary of Scheme

✓ **Compilation** versus **interpretation**

➤ **Compilation**



## 2. Historic Summary of Scheme

### ✓ **Compilation** versus **interpretation**

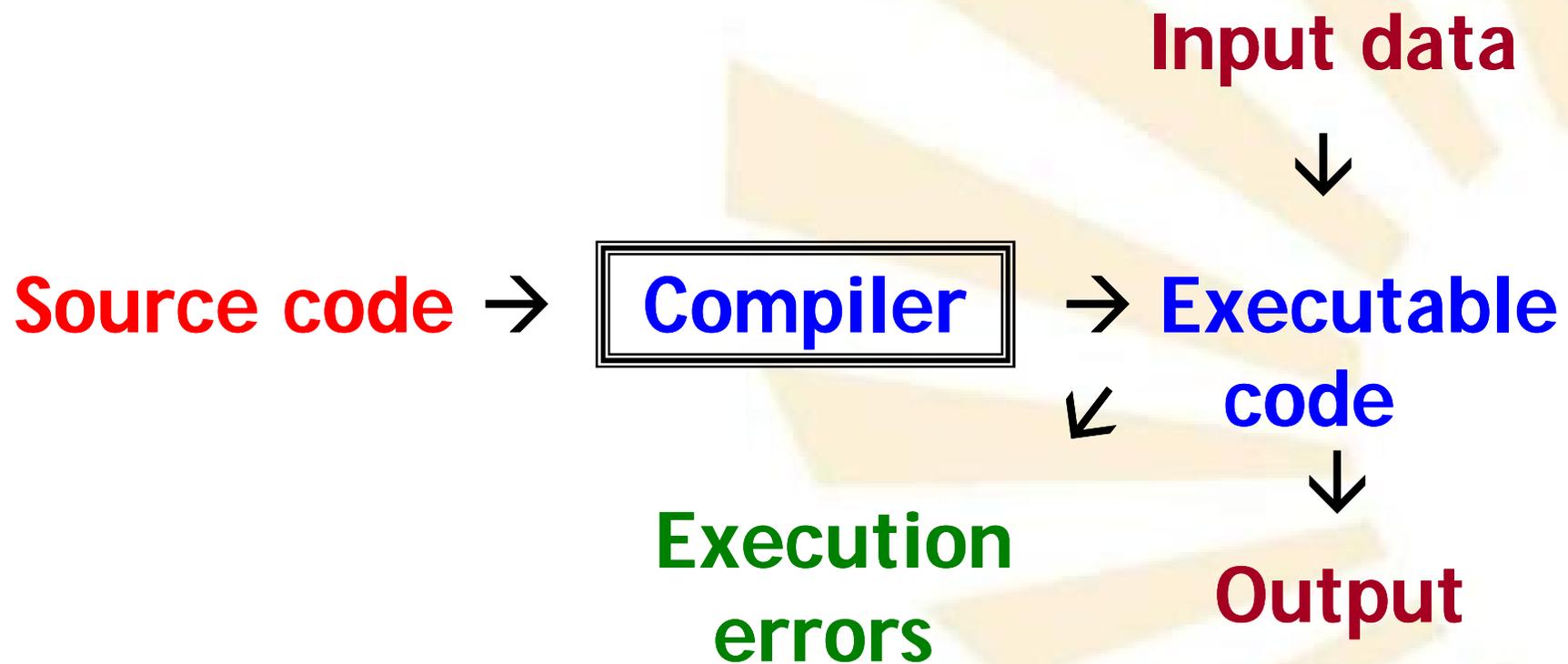
#### ➤ **Compilation**



## 2. Historic Summary of Scheme

### ✓ **Compilation** versus **interpretation**

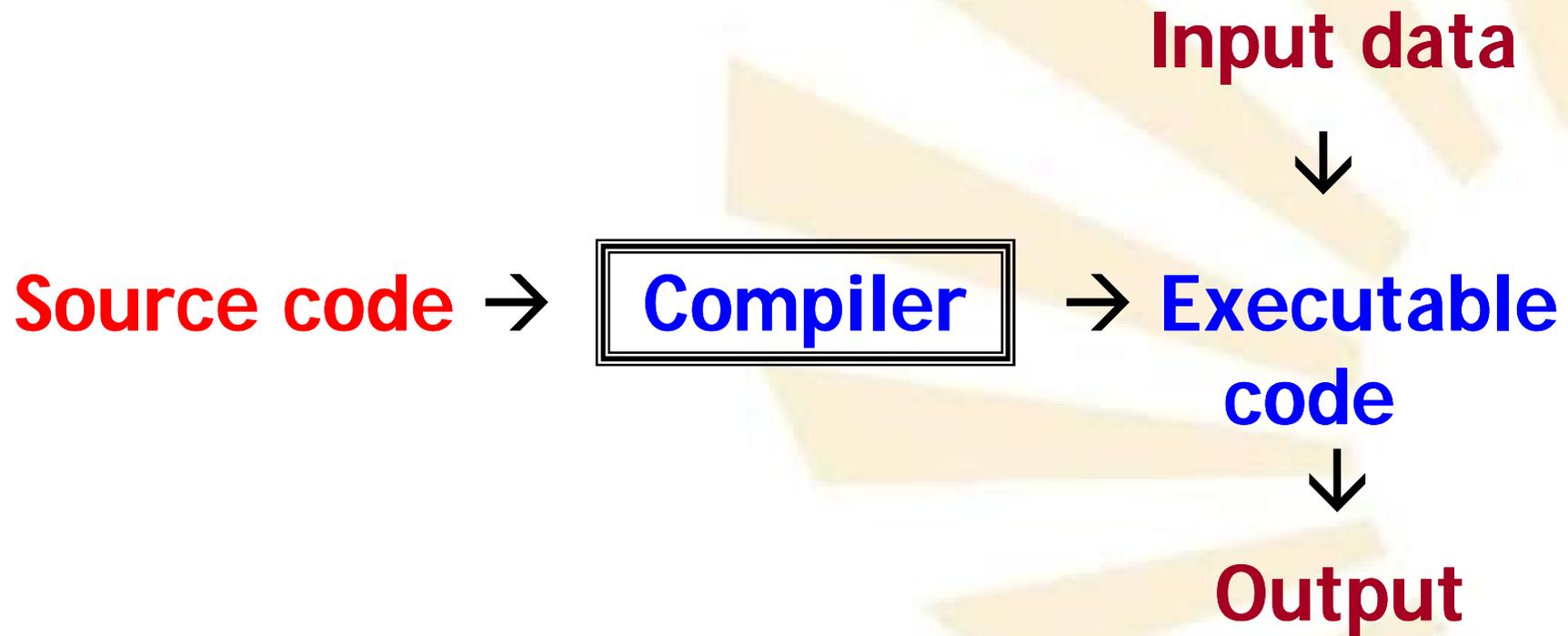
#### ➤ **Compilation**



## 2. Historic Summary of Scheme

### ✓ **Compilation** versus **interpretation**

#### ➤ **Compilation**



## 2. **Historic Summary of Scheme**

✓ **Compilation versus interpretation**

➤ **Interpretation**

## 2. **Historic Summary of Scheme**

- ✓ **Compilation versus interpretation**

- **Interpretation** or **simulation**: consists of a cycle of three stages

## 2. **Historic Summary of Scheme**

### ✓ **Compilation versus interpretation**

➤ **Interpretation** or simulation: consists of a cycle of three stages

1. **Analysis**: the source code is analysed to determine the following correct sentence to be run.

## 2. **Historic Summary of Scheme**

### ✓ **Compilation versus interpretation**

➤ **Interpretation** or simulation: consists of a cycle of three stages

1. **Analysis**: the source code is analysed to determine the following correct sentence to be run.
2. **Generation**: the sentence is transformed into executable code.

## 2. **Historic Summary of Scheme**

### ✓ **Compilation versus interpretation**

➤ **Interpretation** or simulation: consists of a cycle of **three** stages

1. **Analysis**: the source code is analysed to determine the following correct sentence to be run.
2. **Generation**: the sentence is transformed into executable code.
3. **Execution**: the executable code is run.

## 2. **Historic Summary of Scheme**

✓ **Compilation versus interpretation**

➤ **Interpretation**

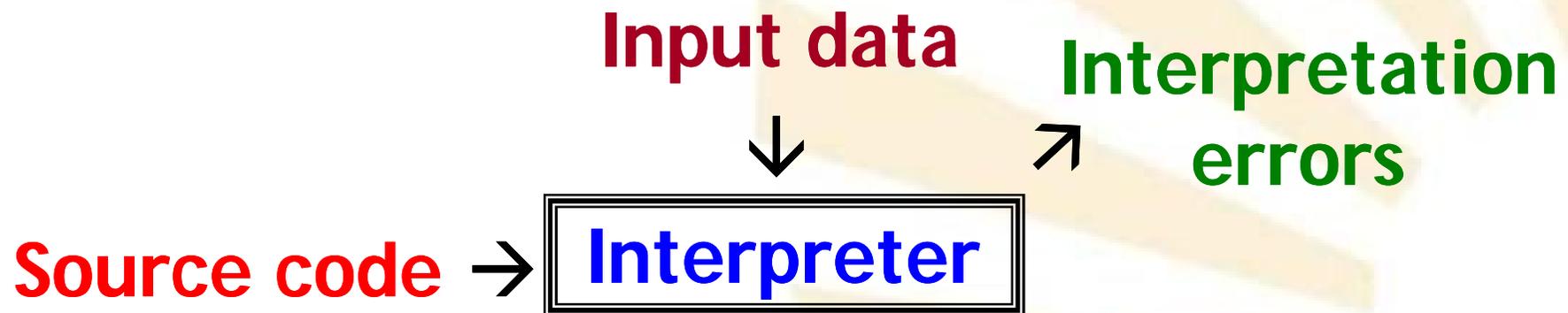
**Source code** →

**Interpreter**

## 2. Historic Summary of Scheme

✓ **Compilation versus interpretation**

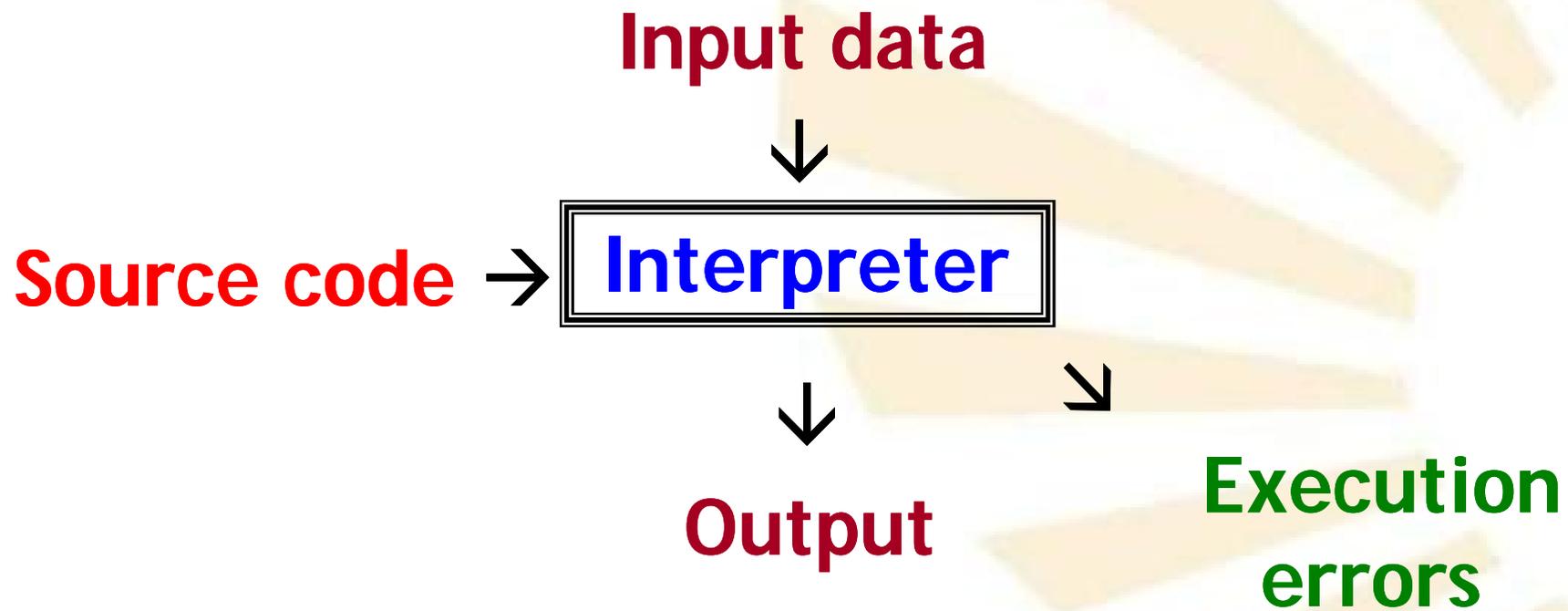
➤ **Interpretation**



## 2. Historic Summary of Scheme

✓ **Compilation** versus **interpretation**

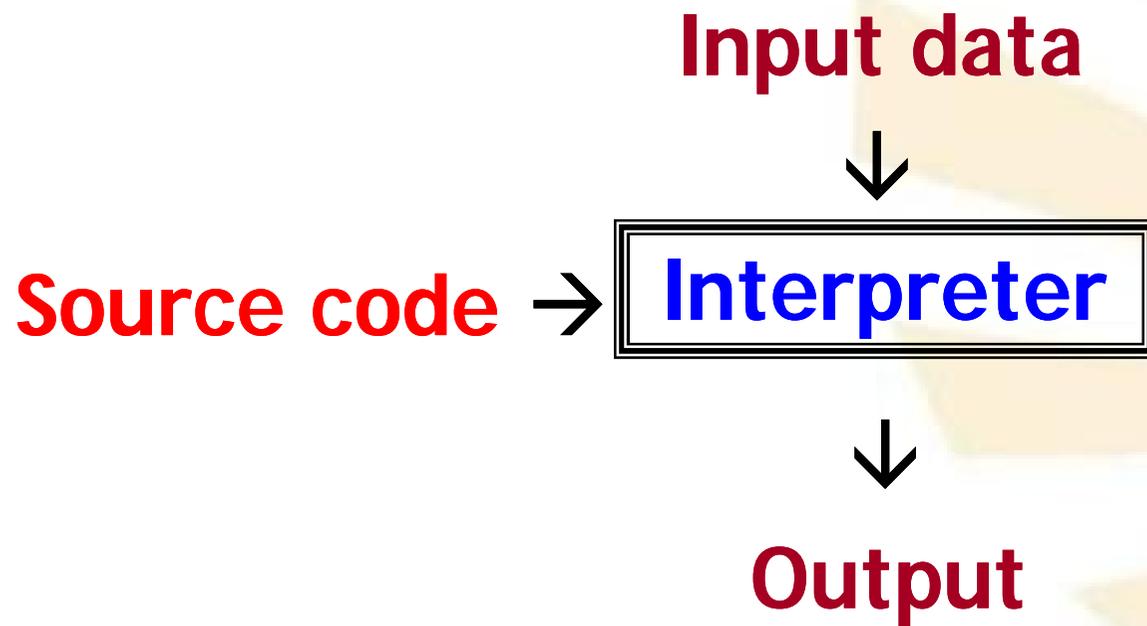
➤ **Interpretation**



## 2. Historic Summary of Scheme

✓ **Compilation** versus **interpretation**

➤ **Interpretation**



## 2. Historic Summary of Scheme

### ✓ **Compilation** versus **interpretation**

#### ■ **Compilation**

- Independent
- Memory necessities
- Efficient
- Global
- No interaction
- **Closed** code during execution

#### ■ **Interpretation**

- Dependent
- No memory necessities
- Less efficient
- Local
- Interaction
- **Open** code during execution

## 2. **Historic Summary of Scheme**

- ✓ LISP
- ✓ Compilation versus Interpretation
- ✓ **Lexical (or static) versus dynamical scope**
- ✓ Origin of Scheme

## 2. Historic Summary of Scheme

### ✓ Lexical (or static) versus dynamical scope

- The **scope rules** determine the **declaration** of **non** local identifiers
- **Non** local identifiers:
  - **Variables** or **functions** which can be **used** in a function or procedure but are **not** declared in that function or procedure
- **Two types**
  - **Lexical or static scope**
    - **With** “blocks structure”: Pascal, **Scheme**
    - **Without** “blocks structure”: C, Fortran
  - **Dynamical scope:**
    - **Always with** “blocks structure”: Lisp, SNOBOL, APL

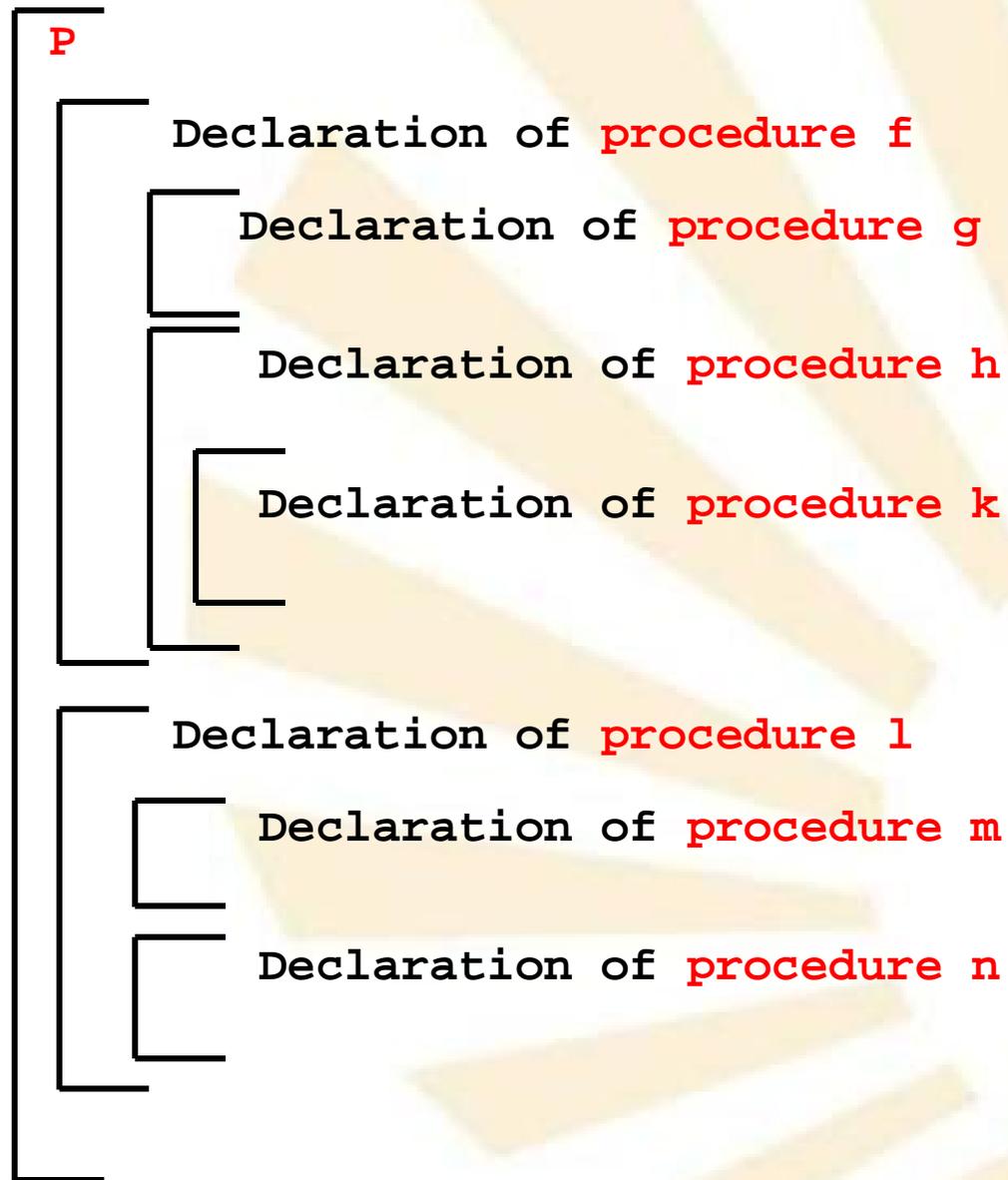
## 2. Historic Summary of Scheme

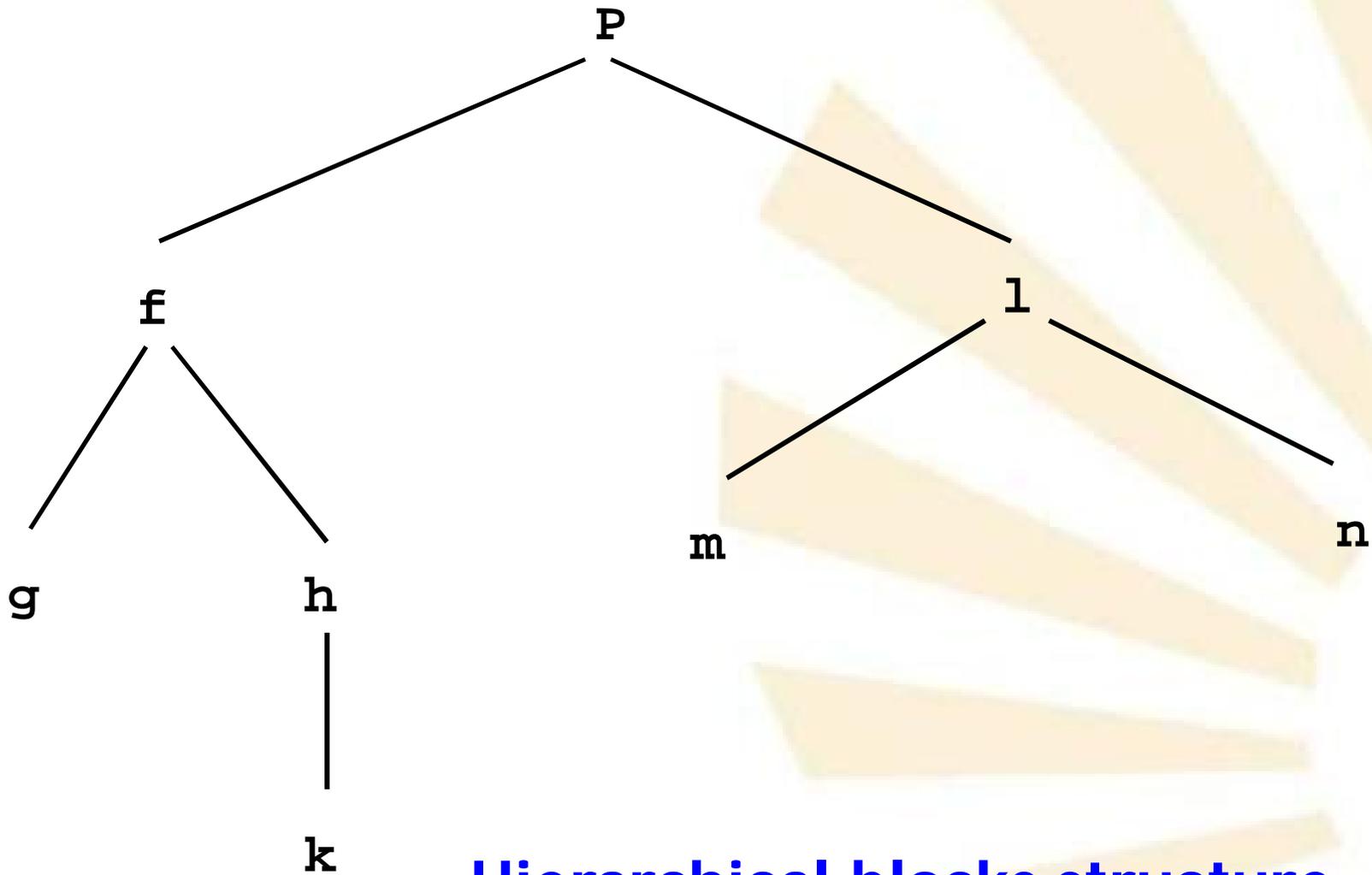
### ✓ Lexical (or static) versus dynamical scope

#### ➤ Block structure

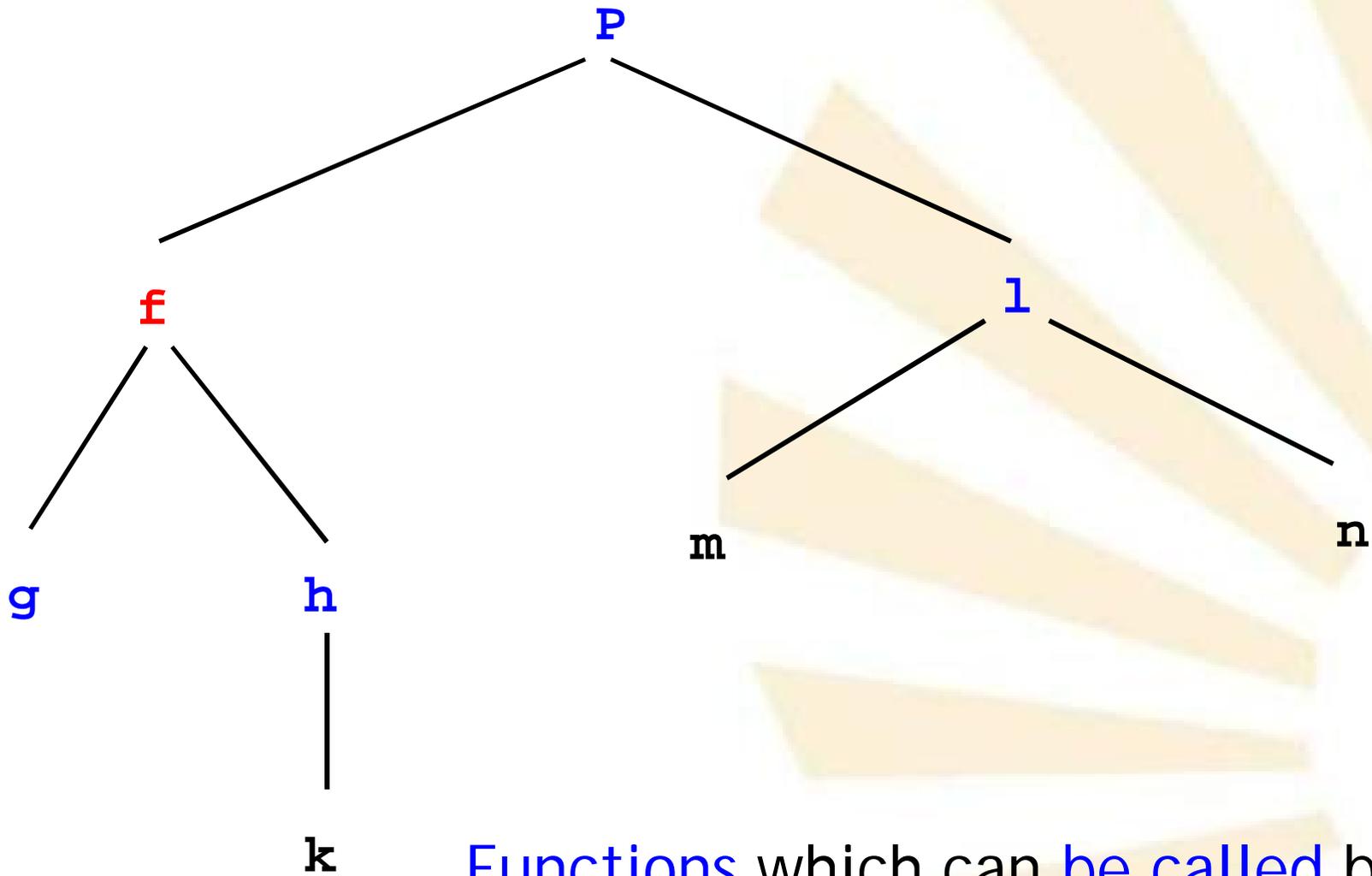
- A procedure or function can **call**
  - Itself
  - Its children (but **not** its grandchildren...)
  - Its brothers (but **not** its nephews)
  - Its father, grandfather, great-grandfather, ...
  - The brothers of its father, grandfather, ...
- A procedure or function can **be called** by
  - Itself
  - Its father (but **not** by its grandfather, ...)
  - Its children, grandchildren, great-grandchildren, ...
  - Its brothers and their children, grandchildren, ...

## Example of blocks structure

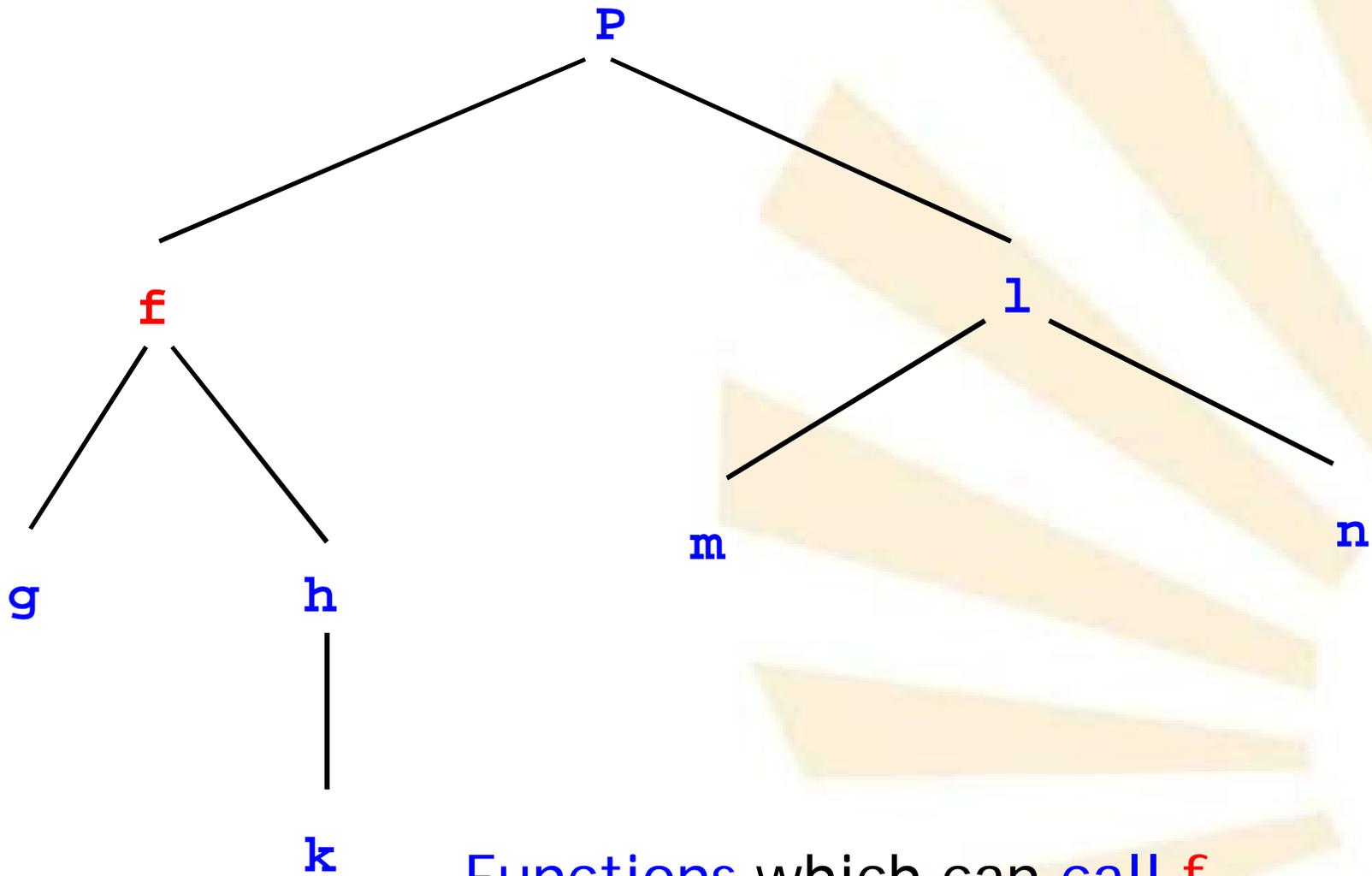




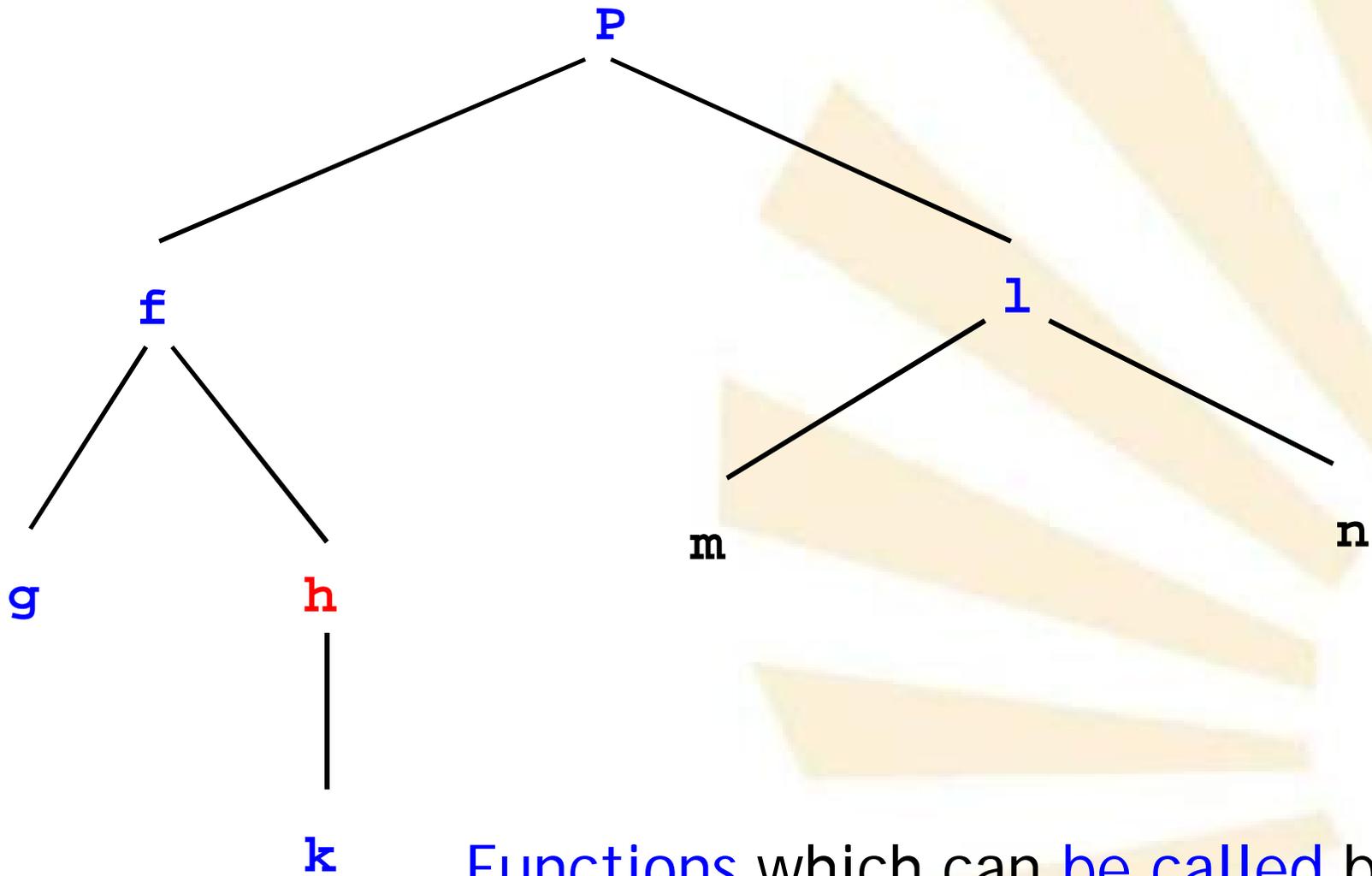
## Hierarchical blocks structure

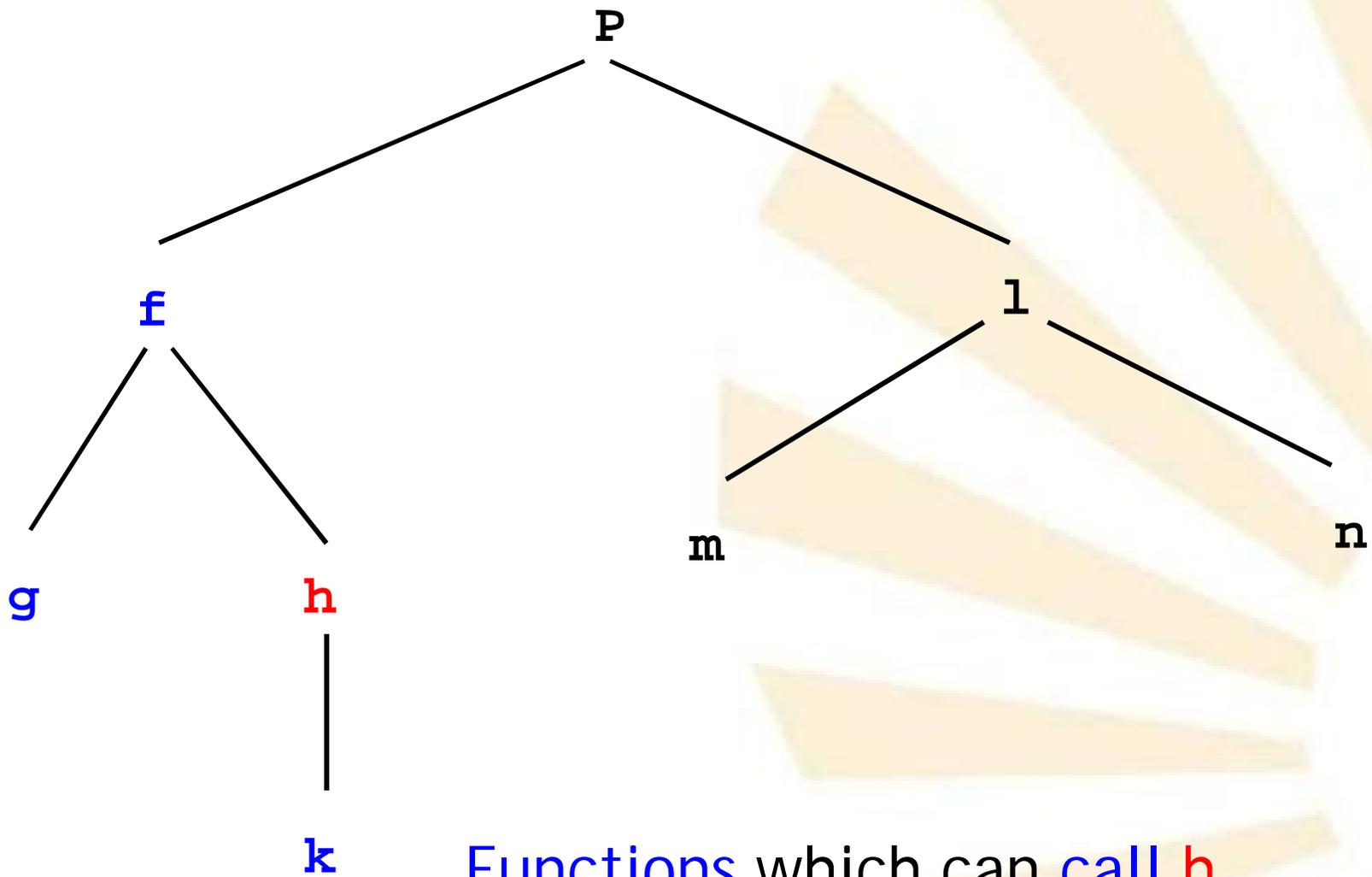


Functions which can be called by **f**



Functions which can call **f**





## 2. Historic Summary of Scheme

- ✓ **Lexical (or static) versus dynamical scope**

- **Lexical** or static scope

- The **declaration** of a **non** local identifier **depends on** the **closest lexical context**
- **The closest nesting rules**

## 2. Historic Summary of Scheme

- ✓ **Lexical (or static) versus dynamical scope**

- **Lexical** or static scope

- The **declaration** of a **non** local identifier **depends on** the **closest lexical context**:

You only have to **read** the program  
to determine the declaration of an identifier.

## 2. Historic Summary of Scheme

### ✓ Lexical (or static) versus dynamical scope

#### ➤ Lexical or static scope

#### ▪ The closest nesting rules:

- The **scope** of a procedure (\*) **f** includes the procedure **f**.
- If a **non** local identifier **x** is used in **f** then the declaration of **x** must be found in the **closest** procedure **g** which includes **f**
- **Notice** (\*) : procedure, function or block

Example.

Lexical scope

with

“block structure”

```
Declaration of procedure h
  Declaration of variable x (x1)
  Declaration of variable y (y1)
  Declaration of variable z (z1)
  Declaration of procedure g
    Declaration of variable x (x2)
    Declaration of variable y (y2)
    Declaration of procedure f
      Declaration of variable x (x3)
      Use of x (→ x3)
      Use of y (→ y2)
      Use of z (→ z1)
    Use of x (→ x2)
    Use of y (→ y2)
    Use of z (→ z1)
    Call to f
  Use of x (→ x1)
  Use of y (→ y1)
  Use of z (→ z1)
  Call to g
```

## 2. Historic Summary of Scheme

- ✓ **Lexical (or static) versus dynamical scope**

- **Lexical or static scope**

- **Without** block structure:

- If **x** is **not** local for a **specific** function then it is **not** local for **all** functions

## Example in C:

without

“block structure”

```
int x; /* x1 */
int y; /* y1 */
int z; /* z1 */
```

```
main()
```

```
{
    int x; /* x2 */
    int y; /* y2 */

    /* Use of x → x2 */
    /* Use of y → y2 */
    /* Use of z → z1 */
    /* Call to f */
    f ();
}
```

```
f()
```

```
{
    int x; /* x3 */
    /* Use of x → x3 */
    /* Use of y → y1 */
    /* Use of z → z1 */
}
```

*Global variables  
are **not**  
recommended*

## 2. Historic Summary of Scheme

### ✓ Lexical (or static) versus dynamical scope

#### ➤ **Dynamical** scope:

- The **declaration** of an **identifier** depends on the execution of the program
- The closest activation rules

## 2. Historic Summary of Scheme

### ✓ Lexical (or static) versus dynamical scope

#### ➤ **Dynamical** scope:

- The **declaration** of an **identifier** depends on the **execution of the program**

You have to **run** the program  
to determine the declaration of an identifier

## 2. Historic Summary of Scheme

### ✓ Lexical (or static) versus dynamical scope

#### ➤ **Dynamical** scope:

##### ▪ **The closest activation rules:**

- The **scope** of a procedure (\*) **f** includes the procedure **f**.
- If a **non** local identifier **x** is used in the **activation** of **f** then the declaration of **x** must be found in the **closest active** procedure **g** with a declaration of **x**
- **Notice** (\*) : procedure, function or block

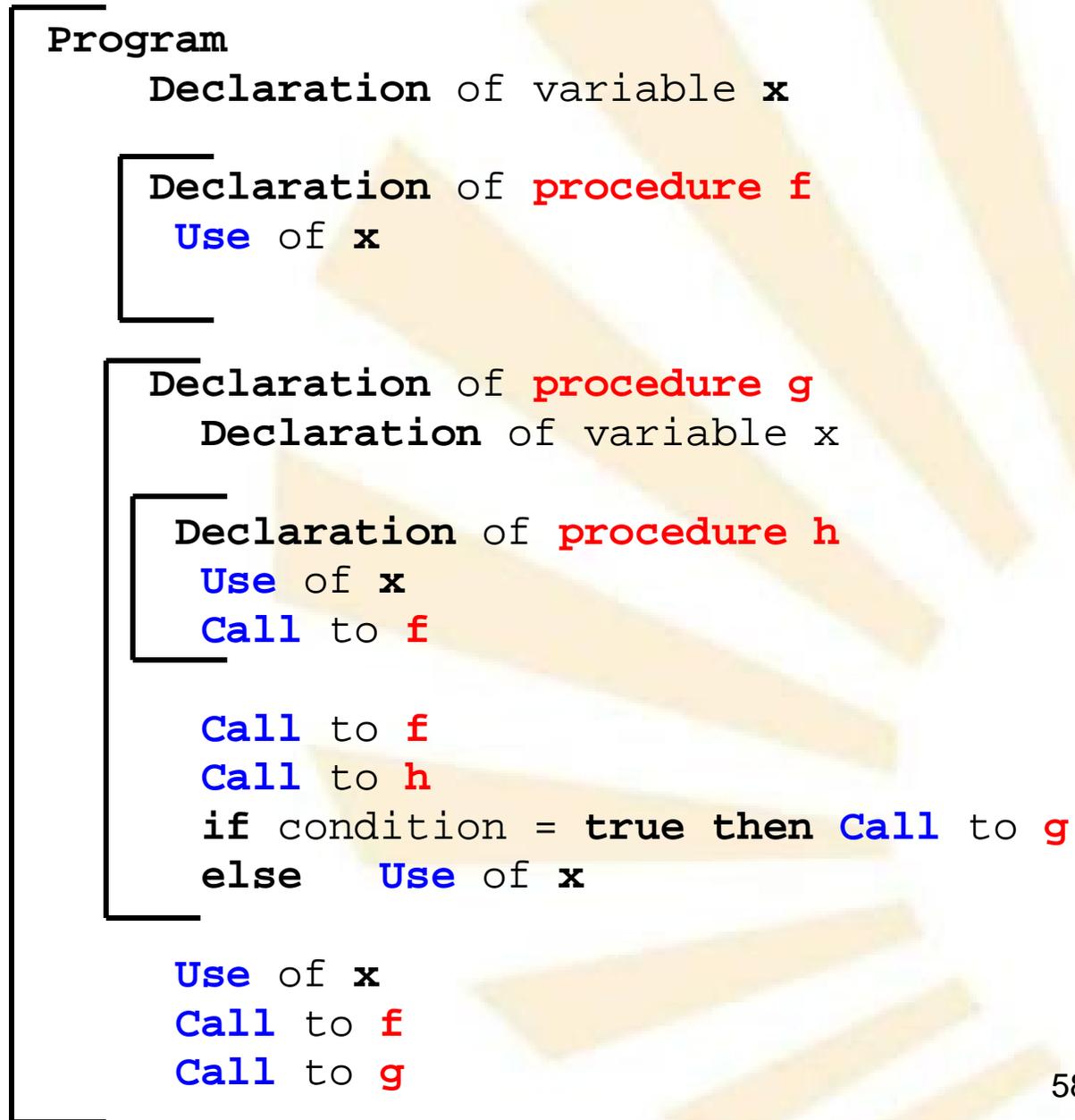
## 2. Historic Summary of Scheme

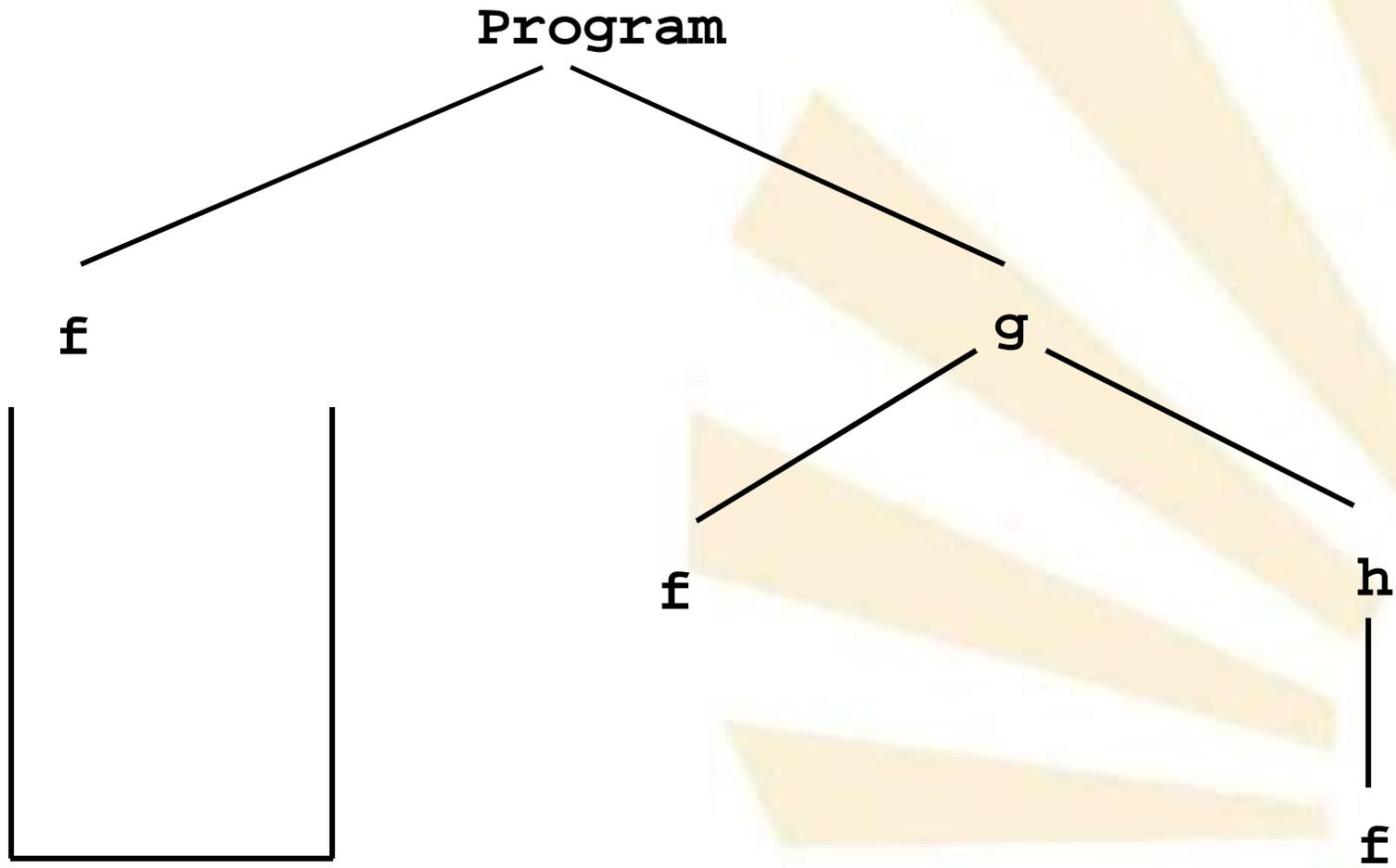
### ✓ Lexical (or static) versus dynamical scope

#### ➤ Notice:

- The **dynamical scope** allows that an **identifier** can be associated to **different declarations** during the program execution

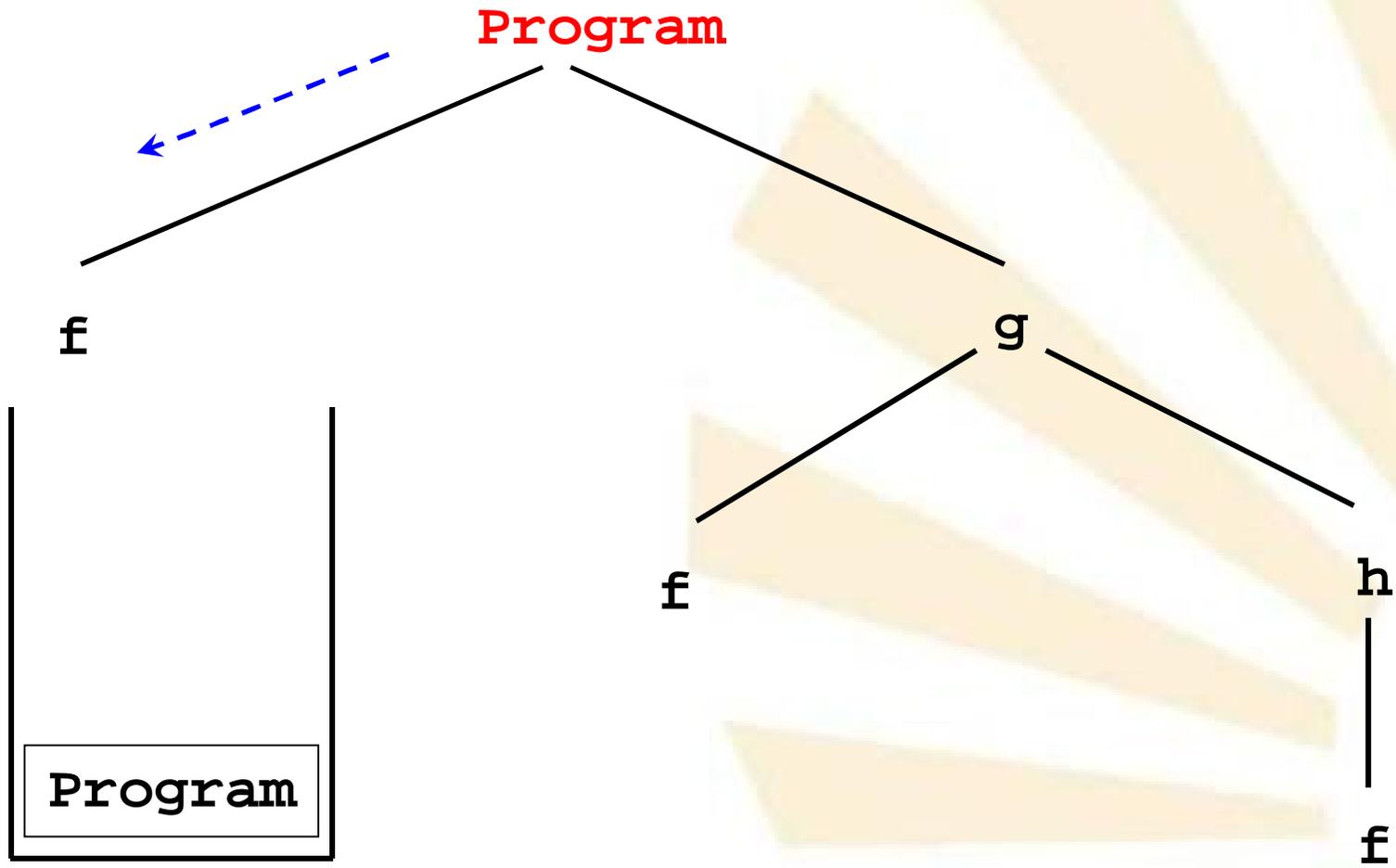
**Example:**  
**Lexical**  
**versus**  
**Dynamical**  
**scope**





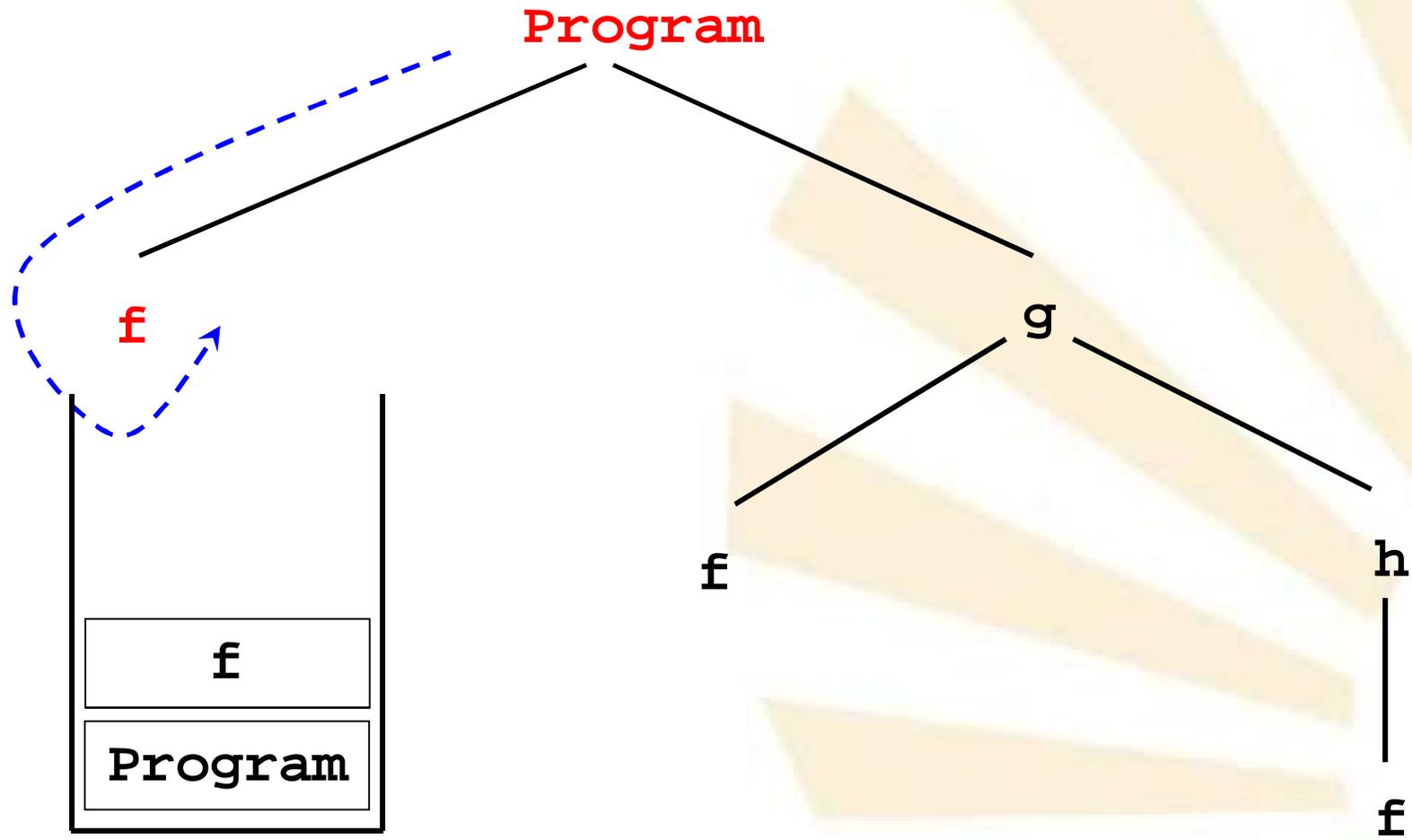
**Activation Stack**

**Activation Tree**



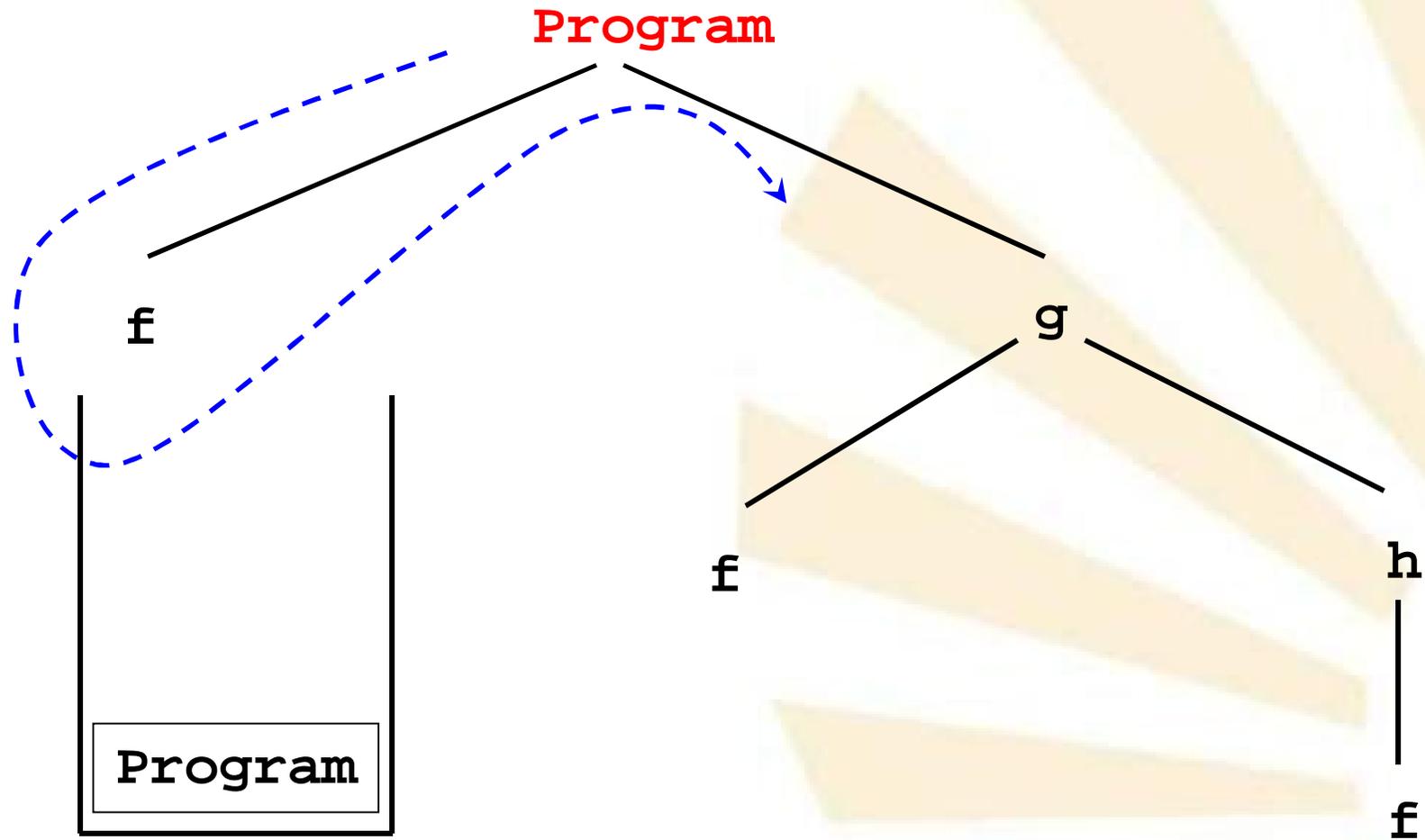
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**Activation Tree**



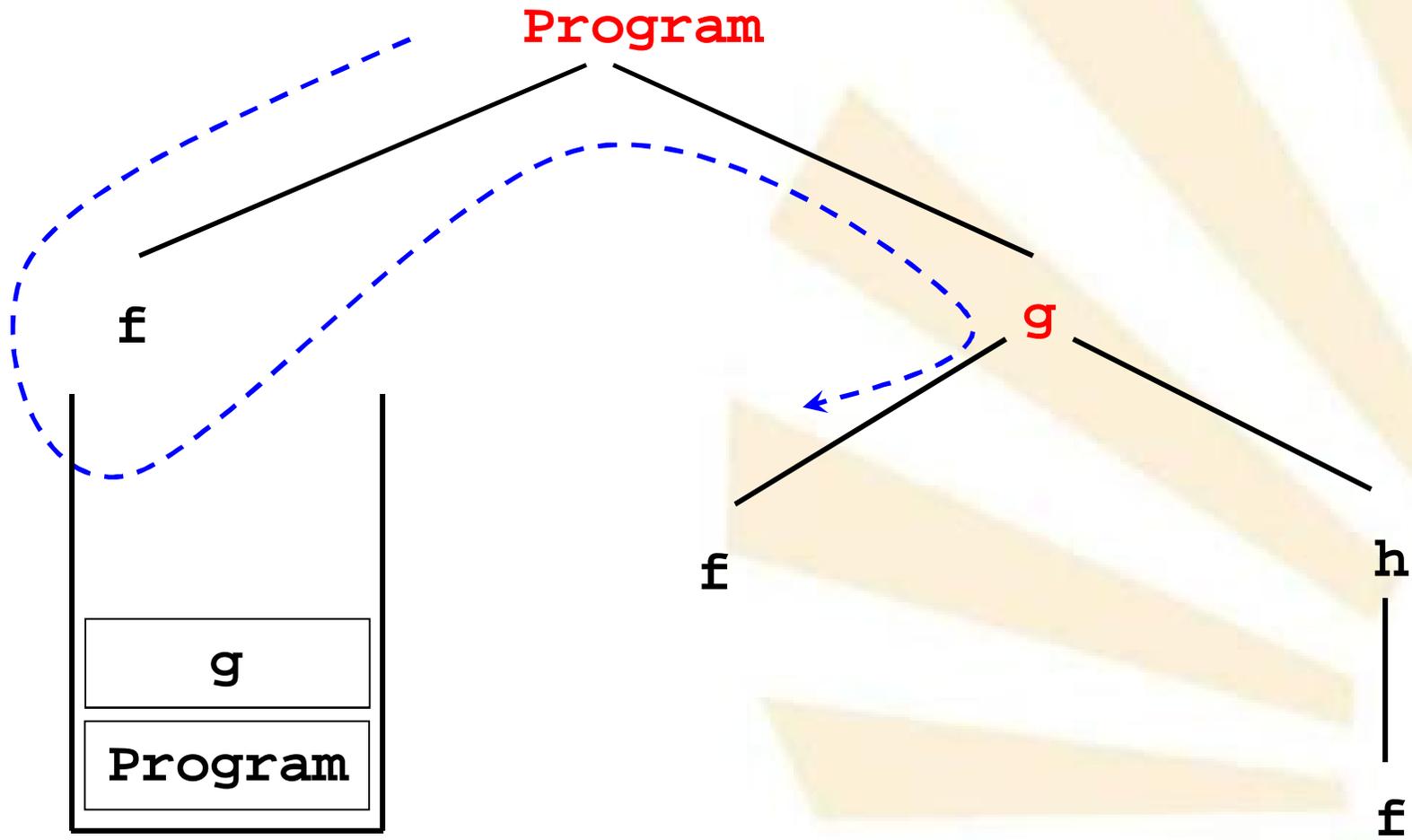
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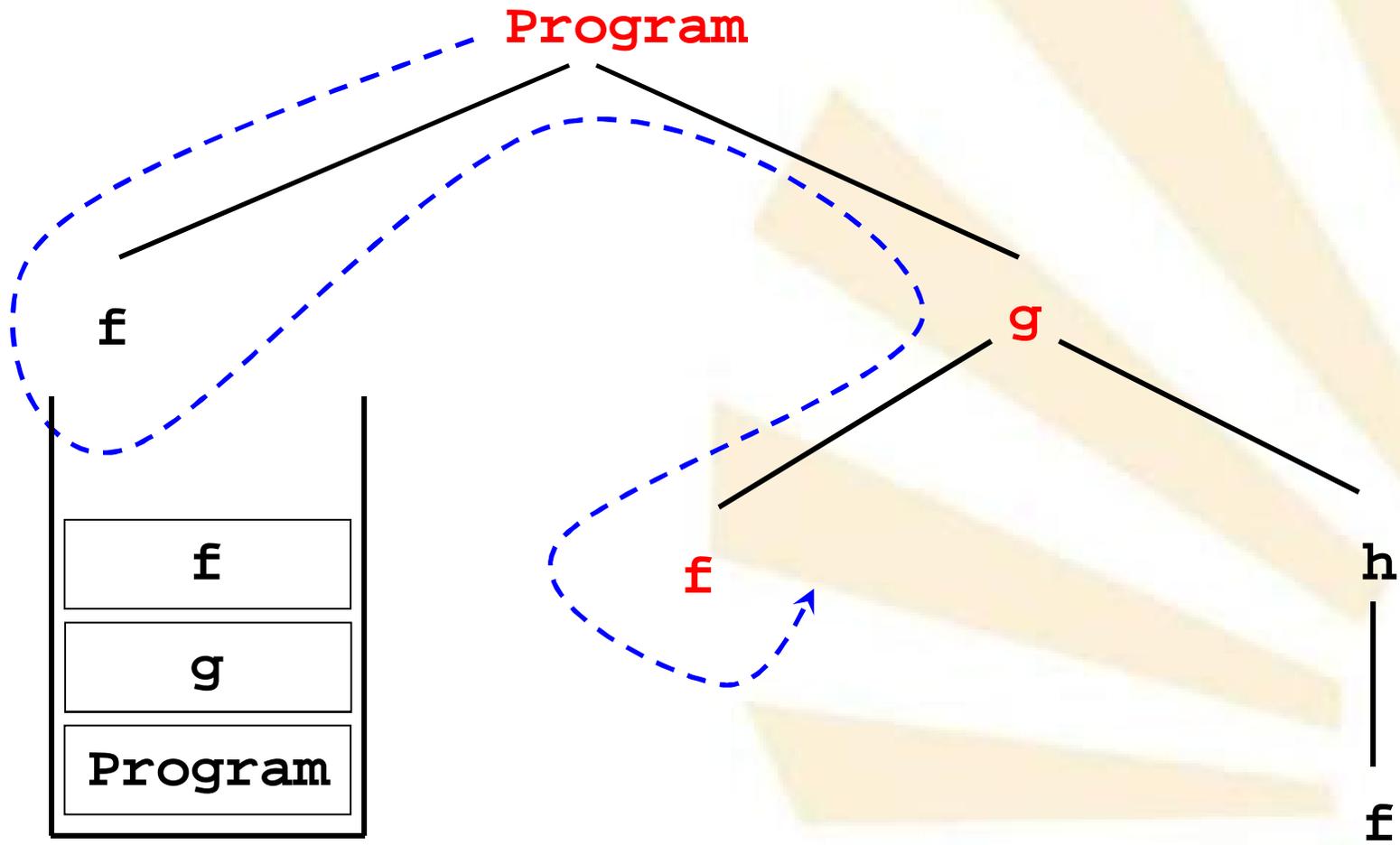
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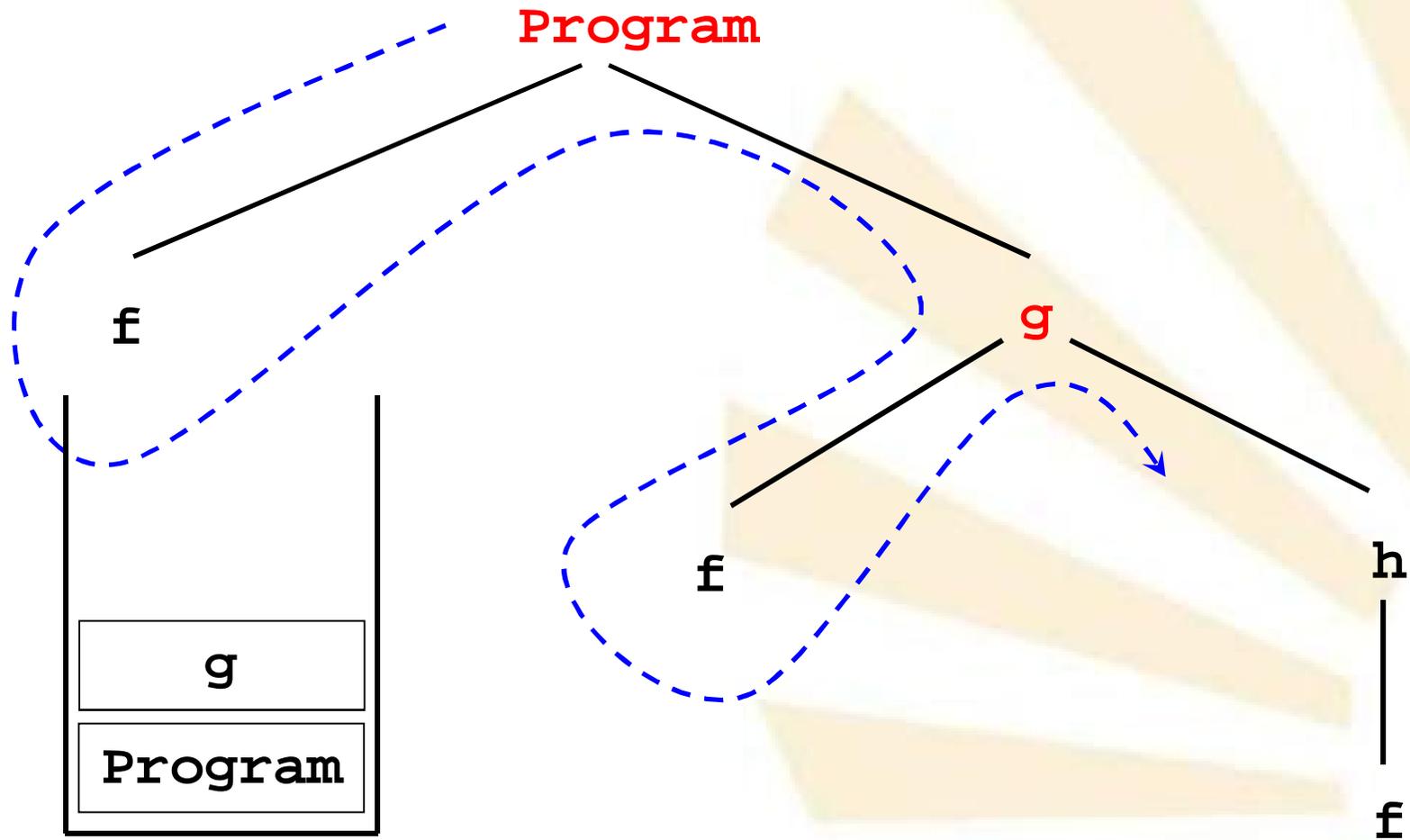
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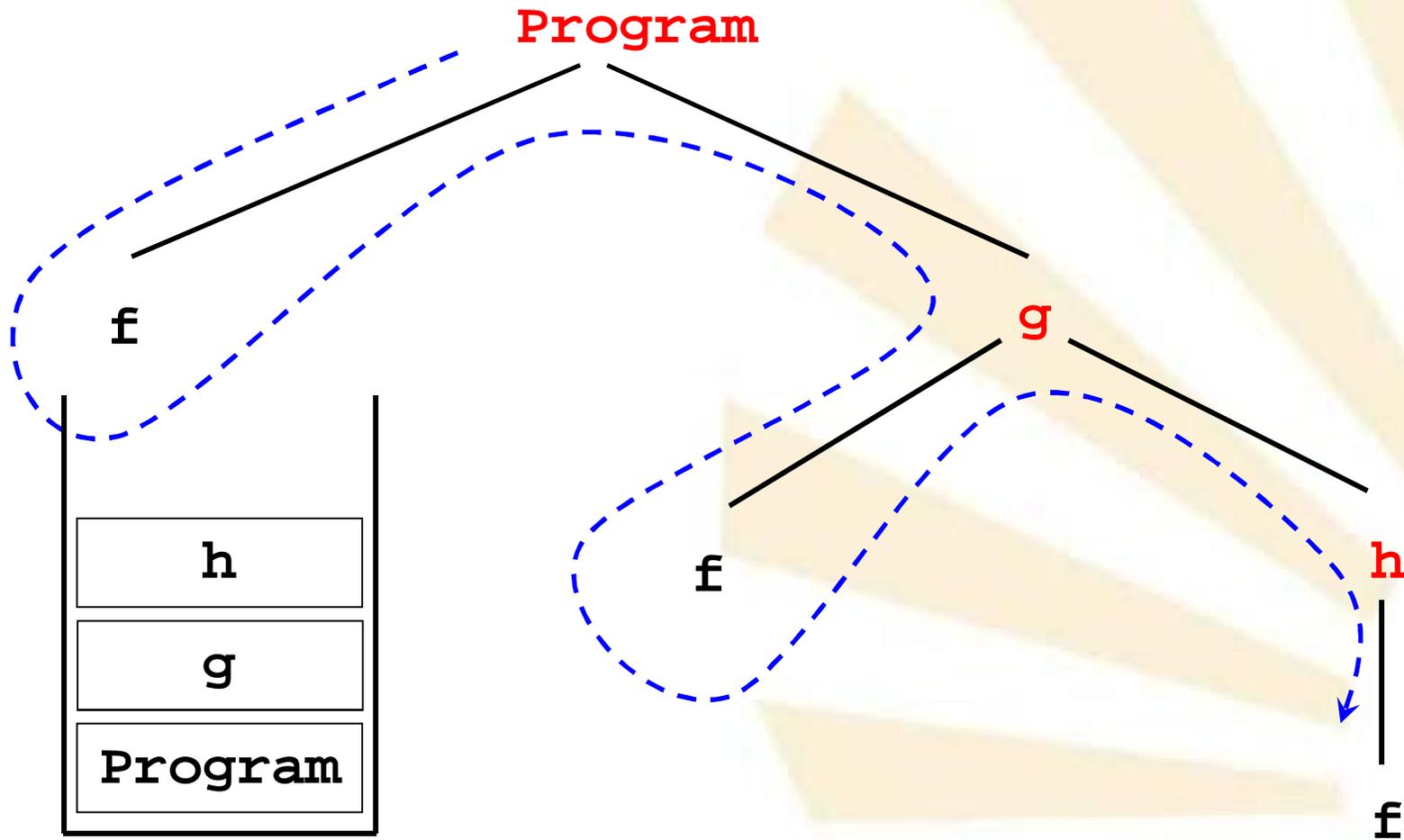
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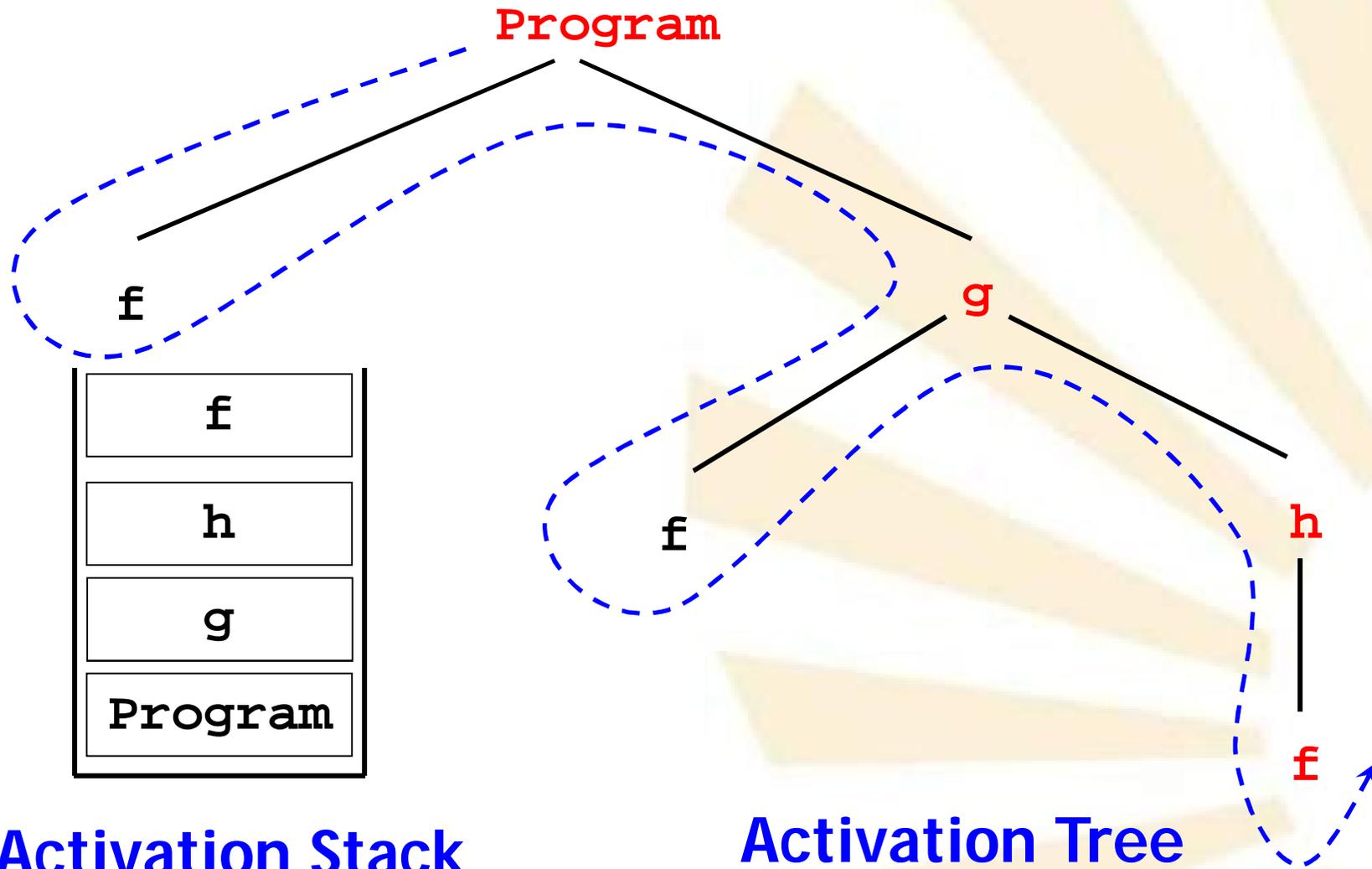
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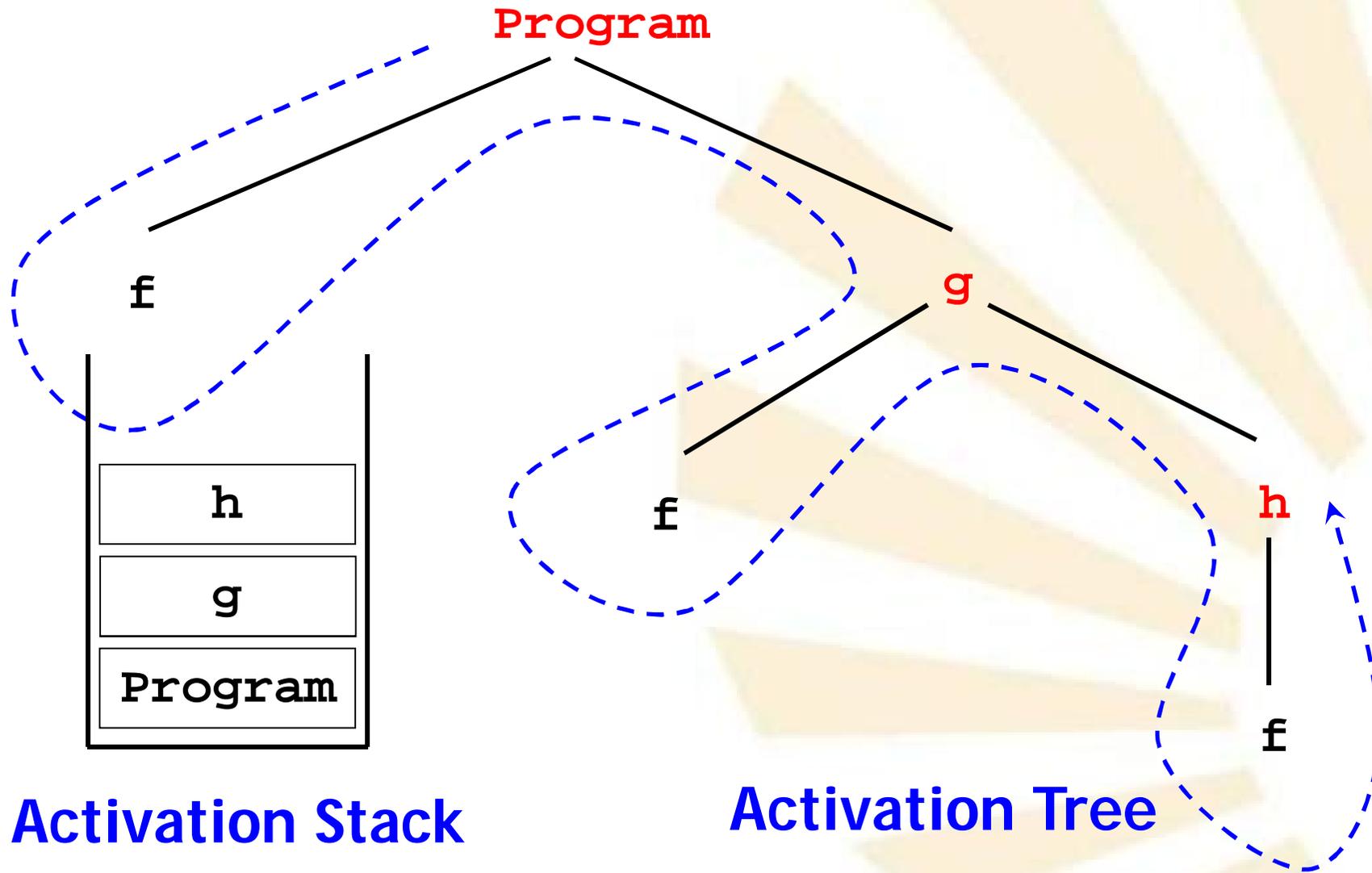
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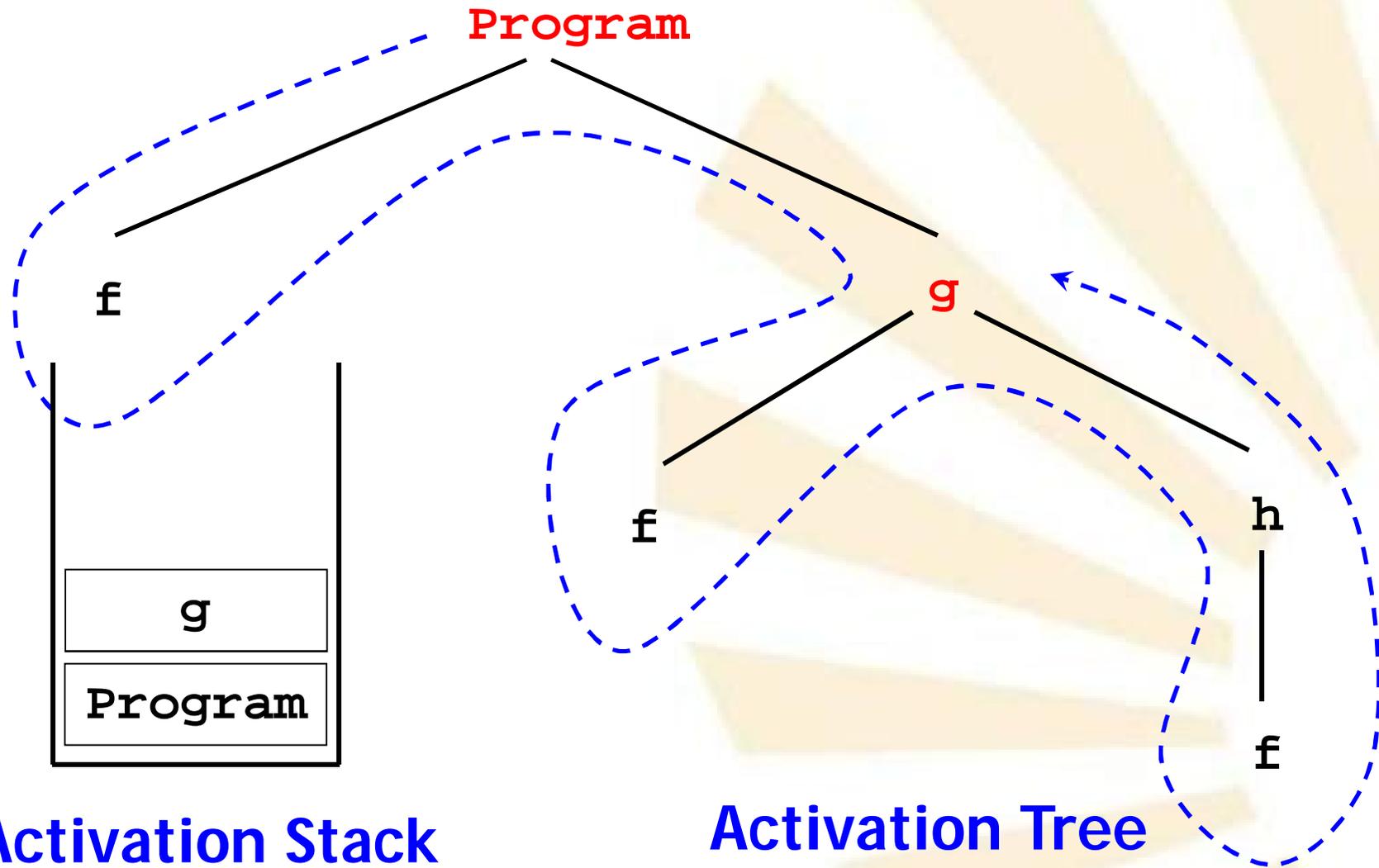
**Activation Tree**



Activation Stack

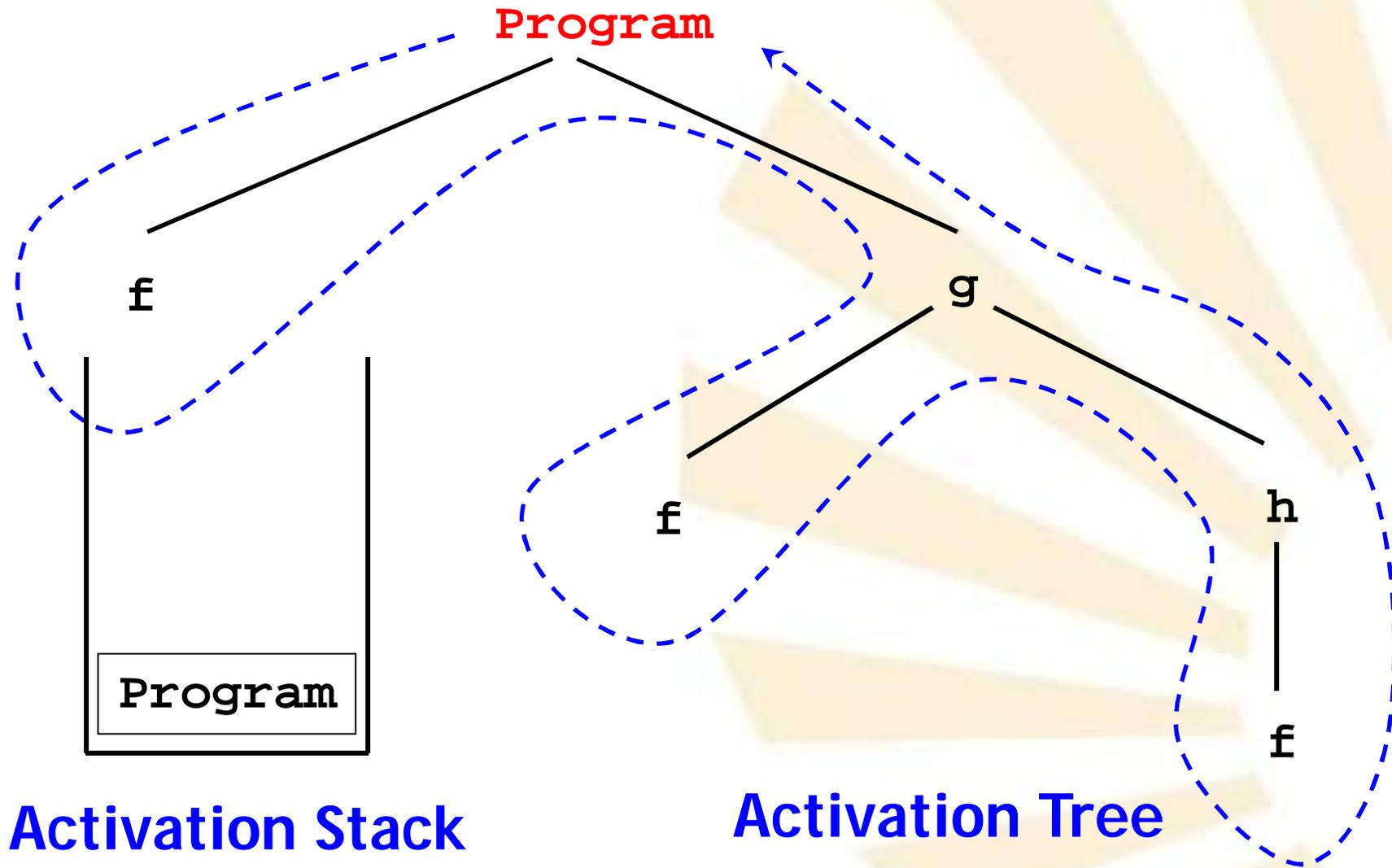
Activation Tree





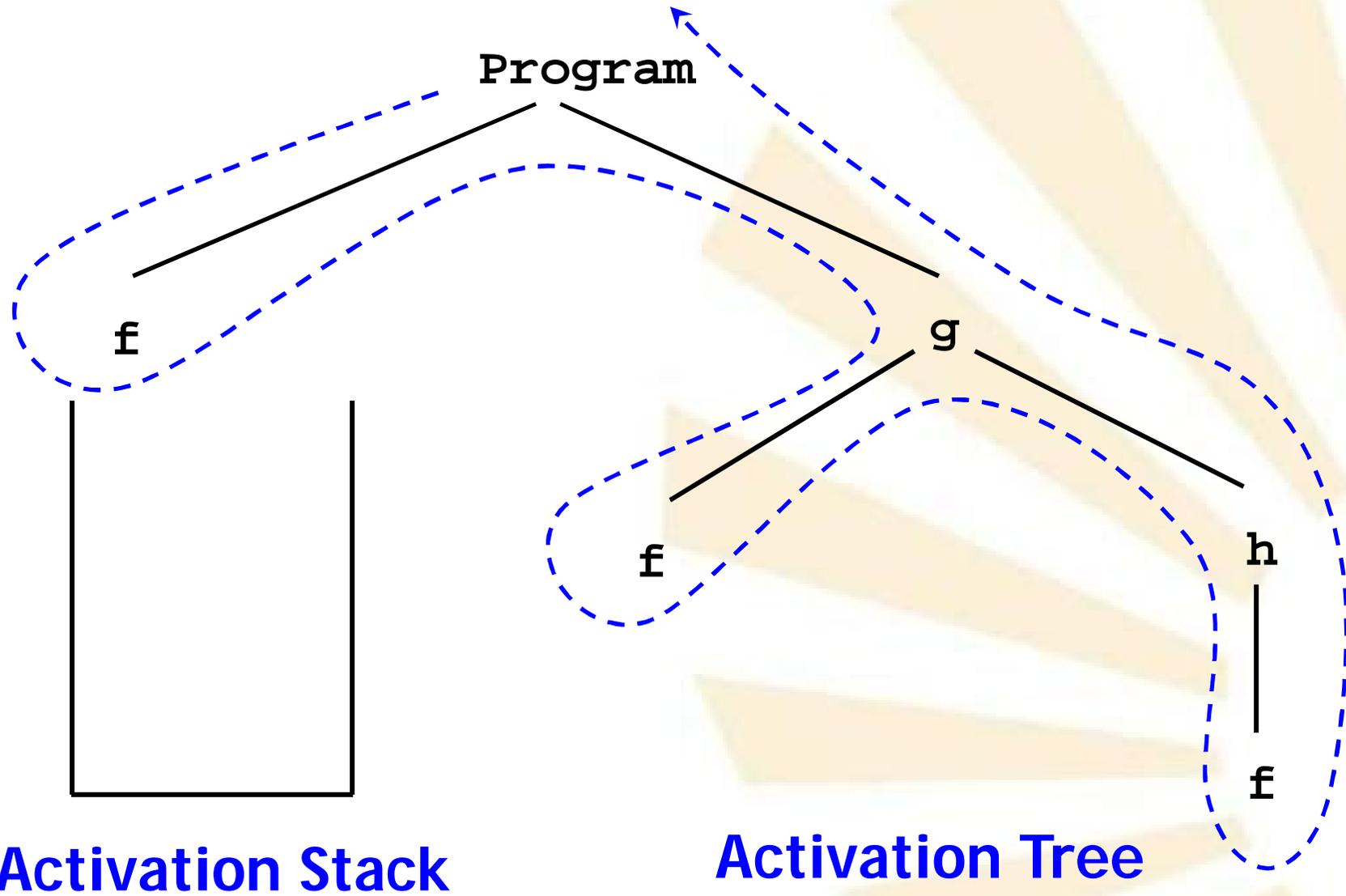
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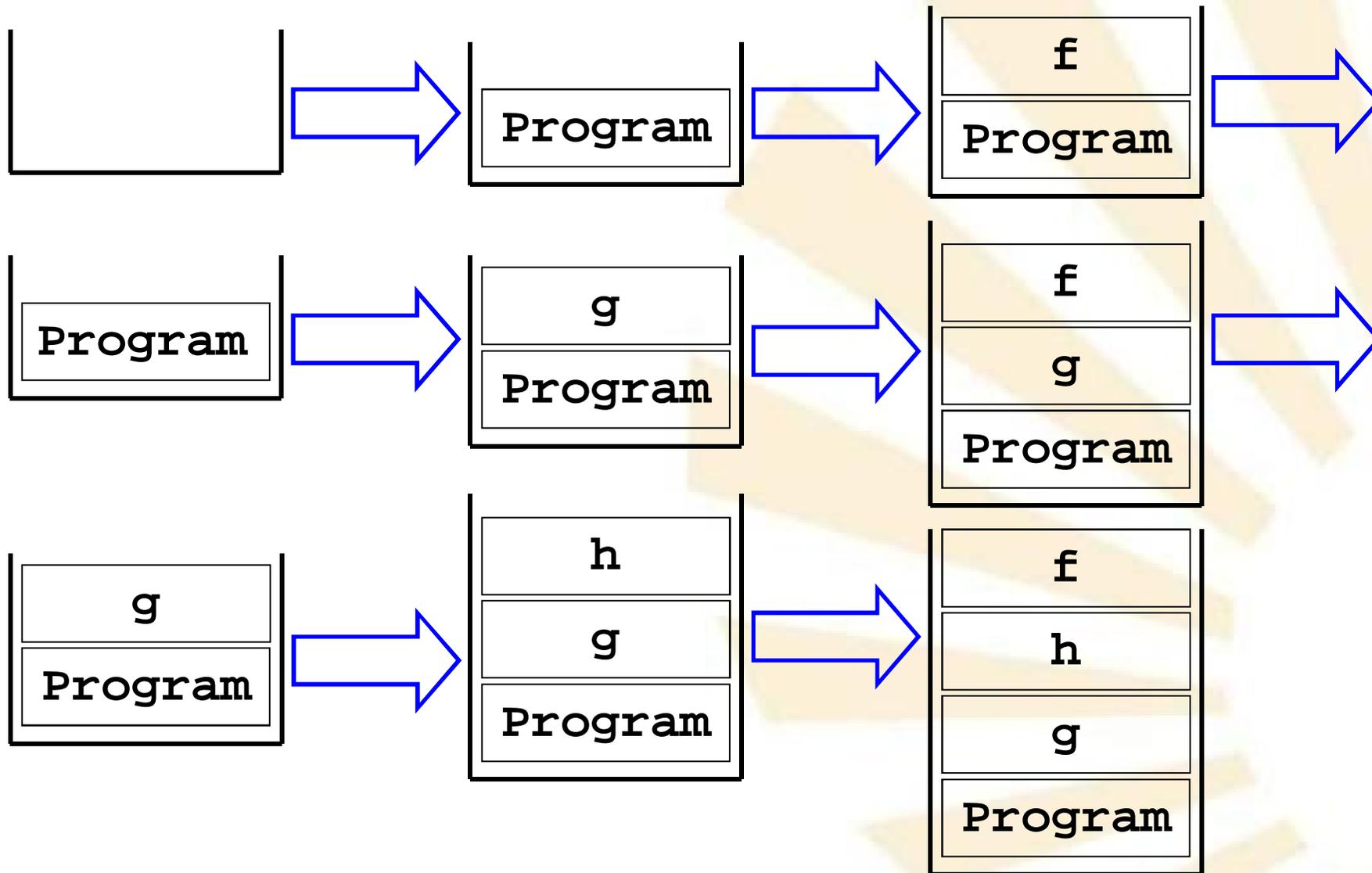
**Activation Tree**



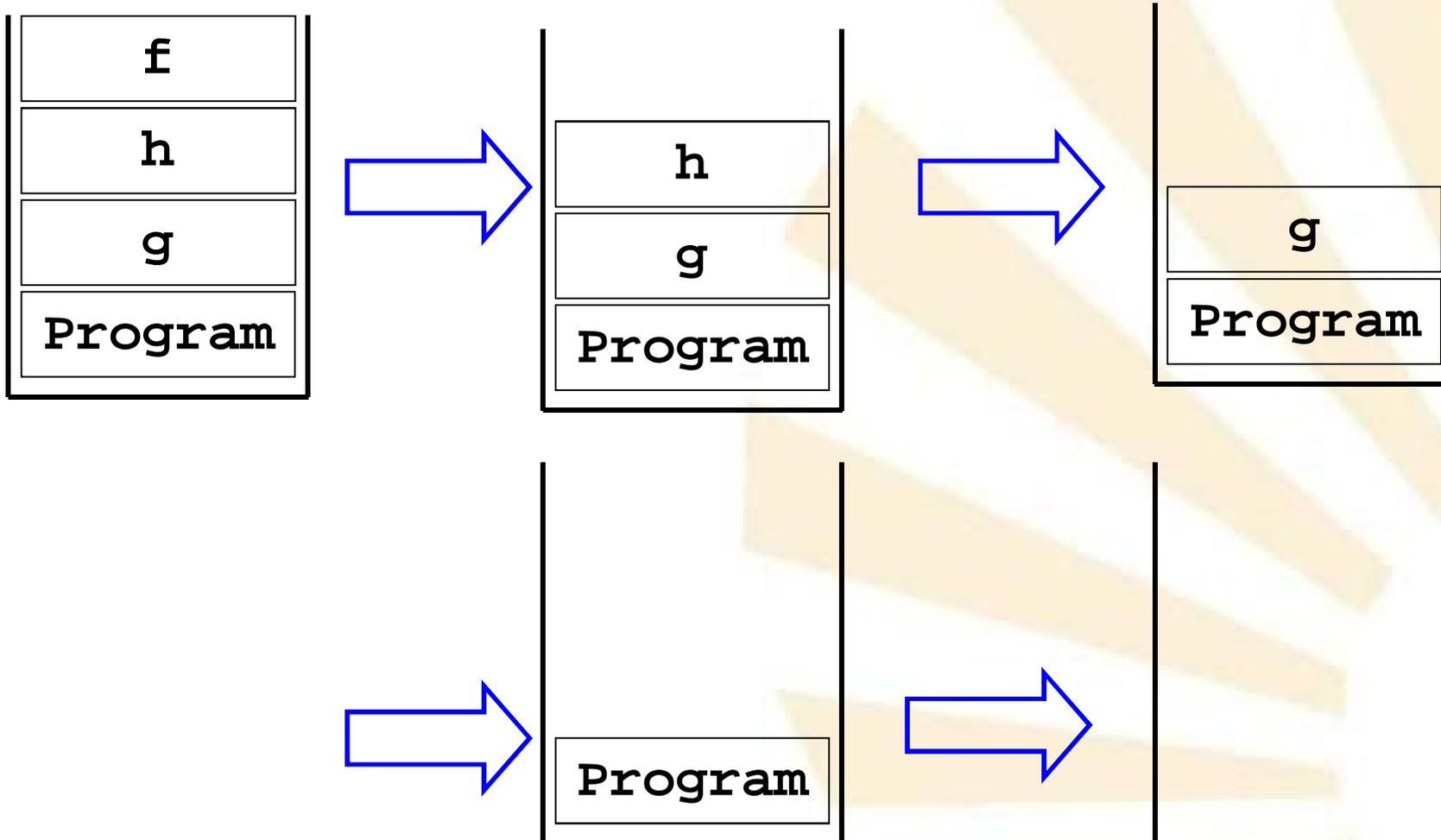
**Activation Stack**

**Activation Tree**





## Changes in the activation Stack (1 / 2)



## Changes in the activation Stack (2 / 2)

# Run with lexical scope

Program

Declaration of variable **x** (**x<sub>1</sub>**)

Declaration of **procedure f**

Use of **x**

Declaration of **procedure g**

Declaration of variable **x** (**x<sub>2</sub>**)

Declaration of **procedure h**

Use of **x**

Call to **f**

Call to **f**

Call to **h**

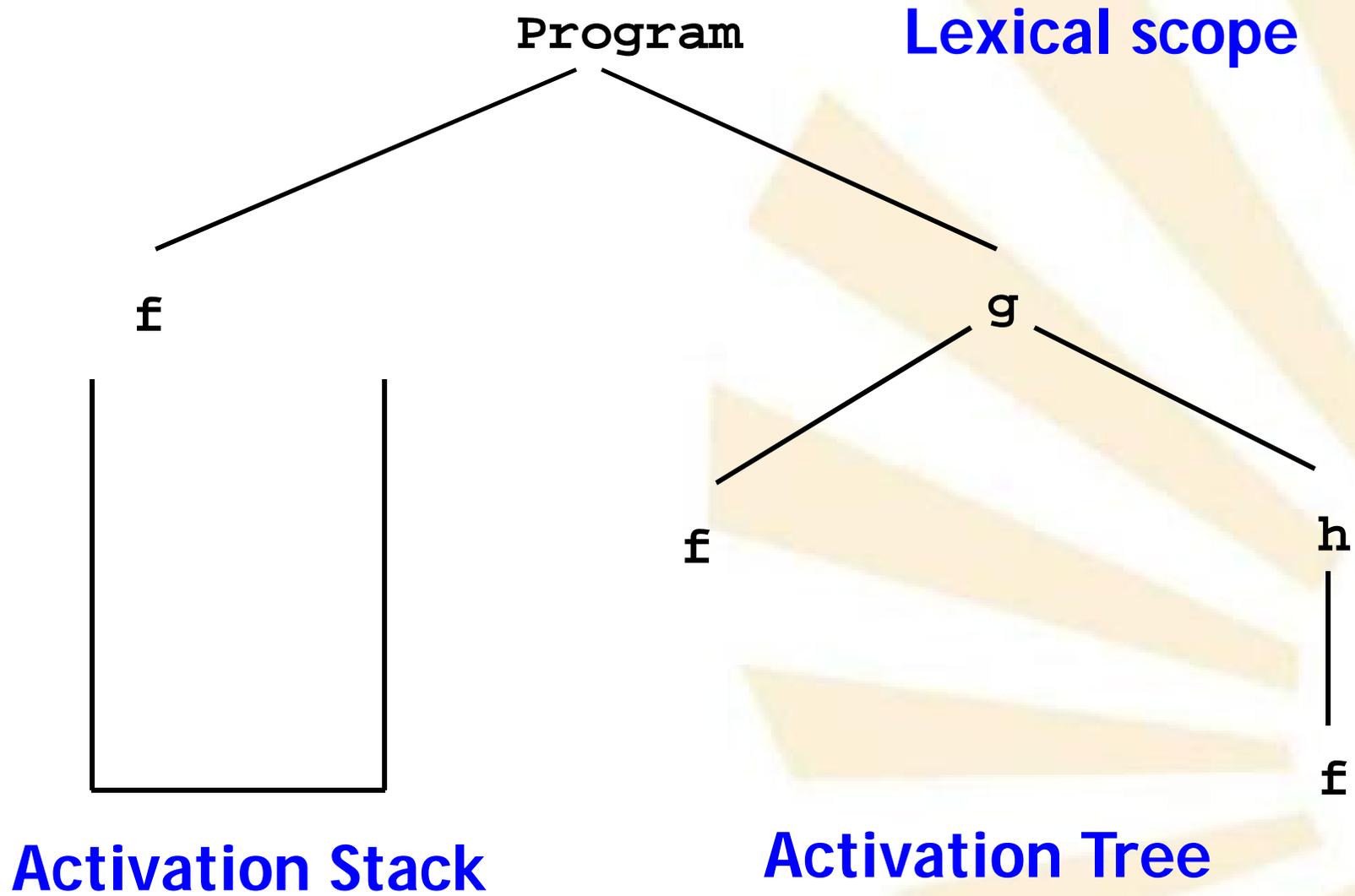
if condition = true then Call to **g**

else Use of **x**

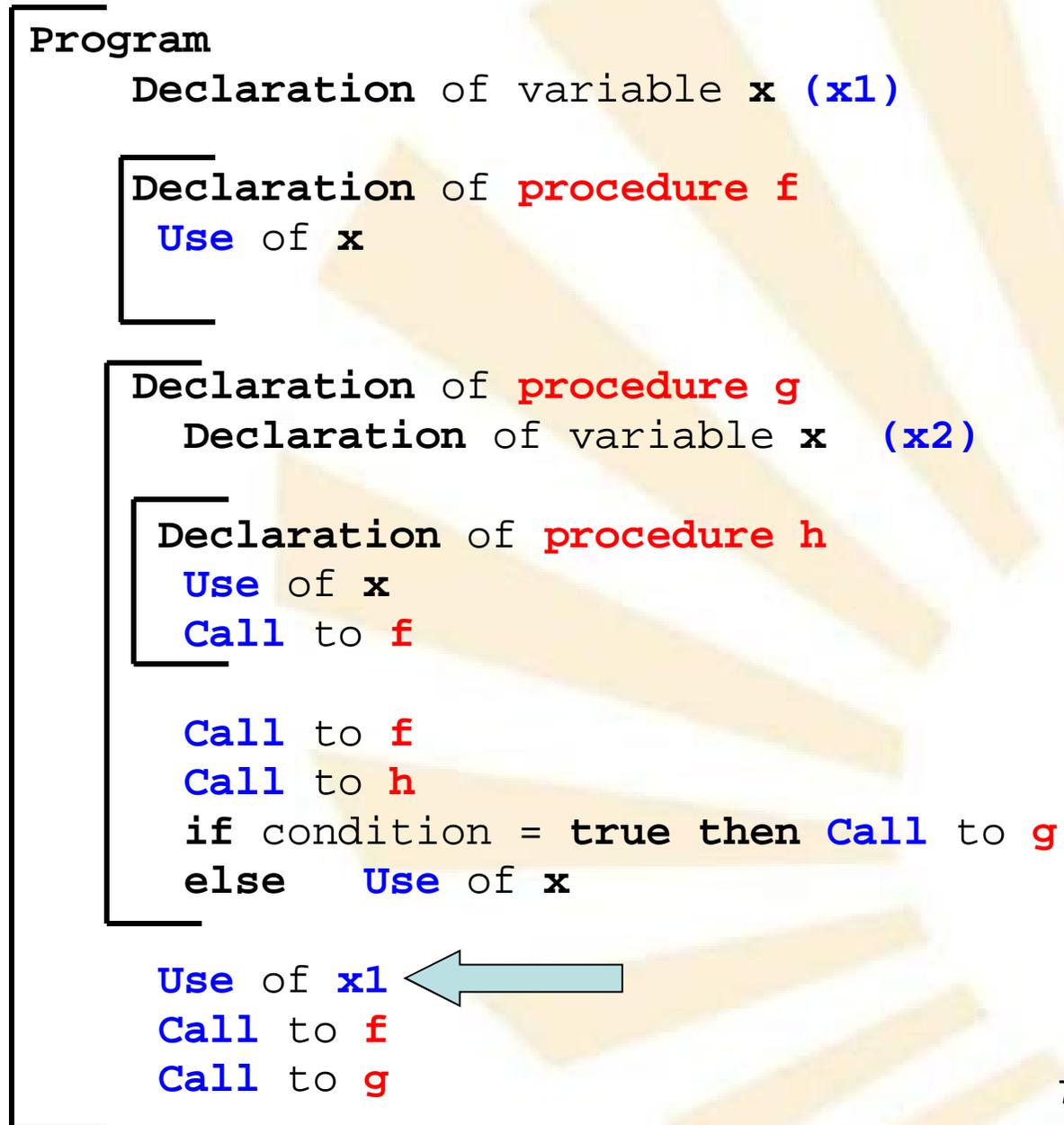
Use of **x**

Call to **f**

Call to **g**



# Run with lexical scope



Program

Lexical scope

- Use of *x1* of Program in Program

f

g

f

h

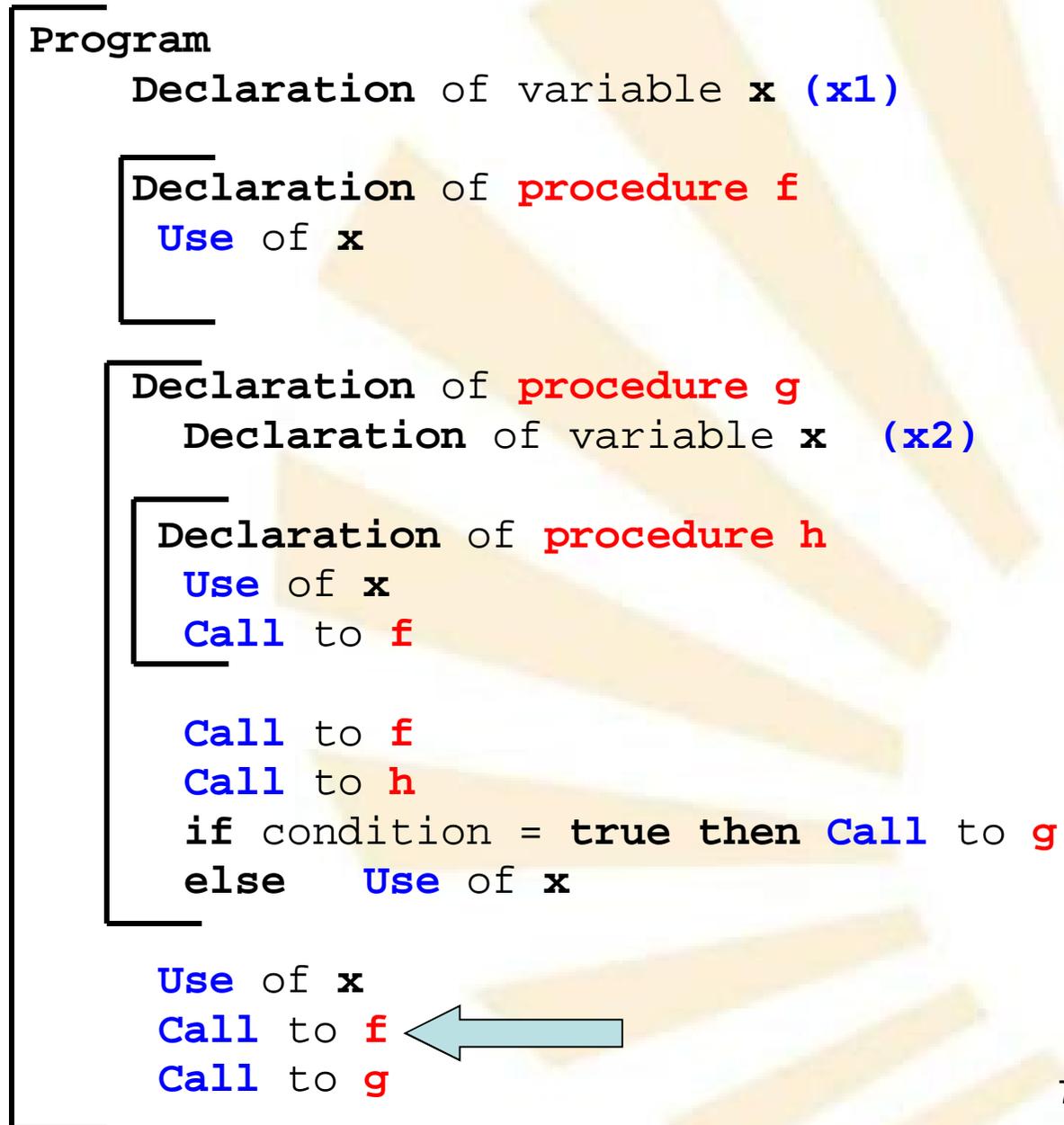
f

Program

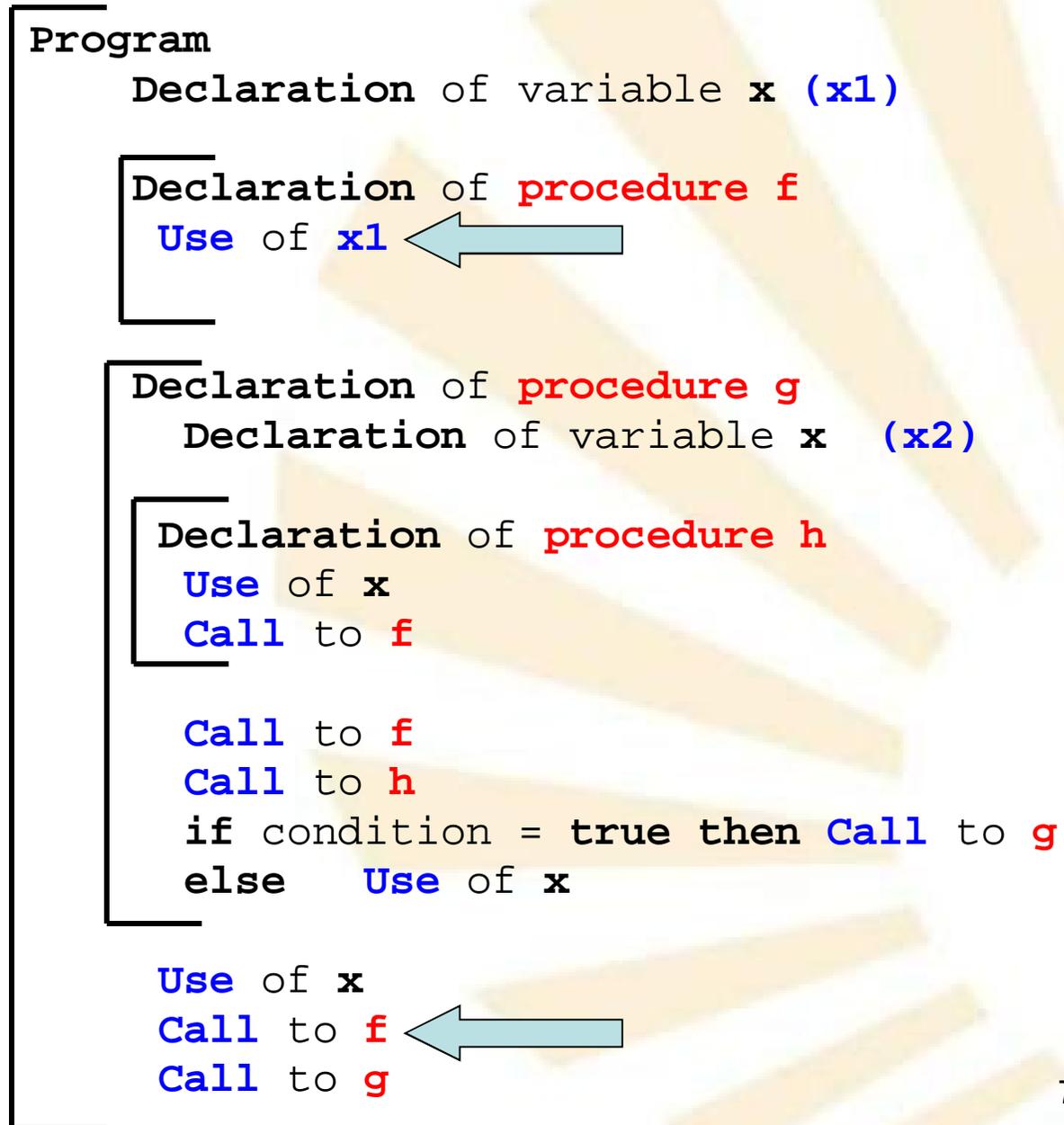
Activation Stack

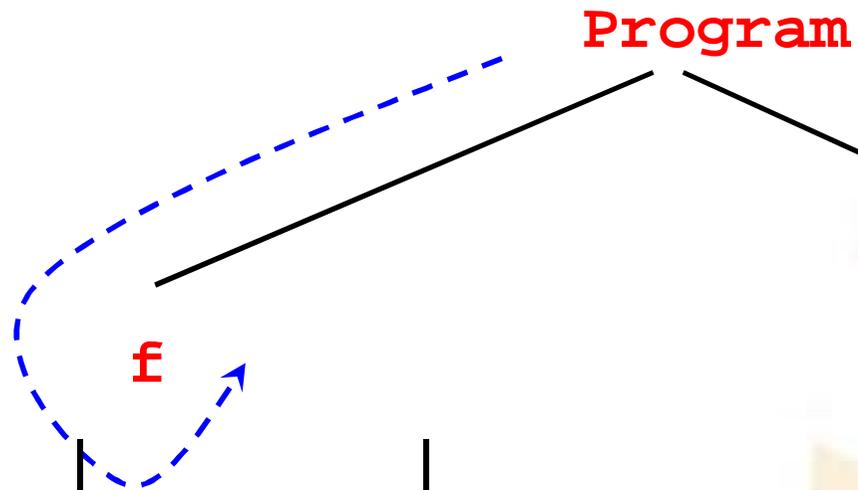
Activation Tree

# Run with lexical scope



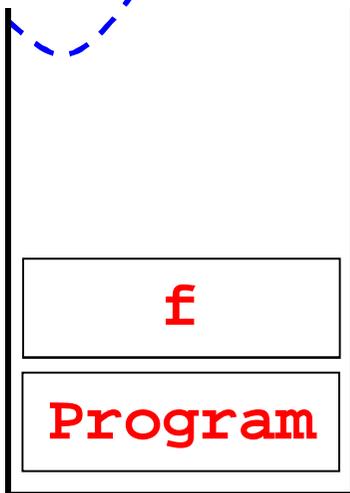
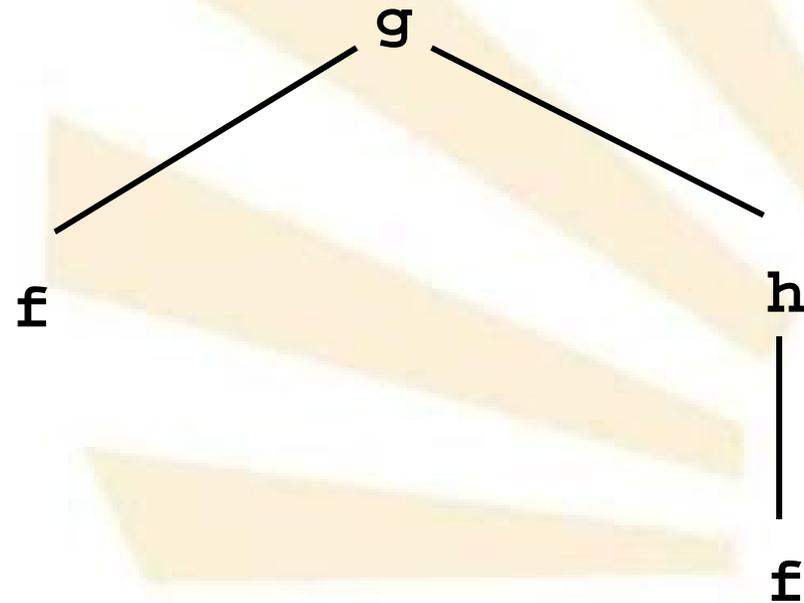
# Run with lexical scope





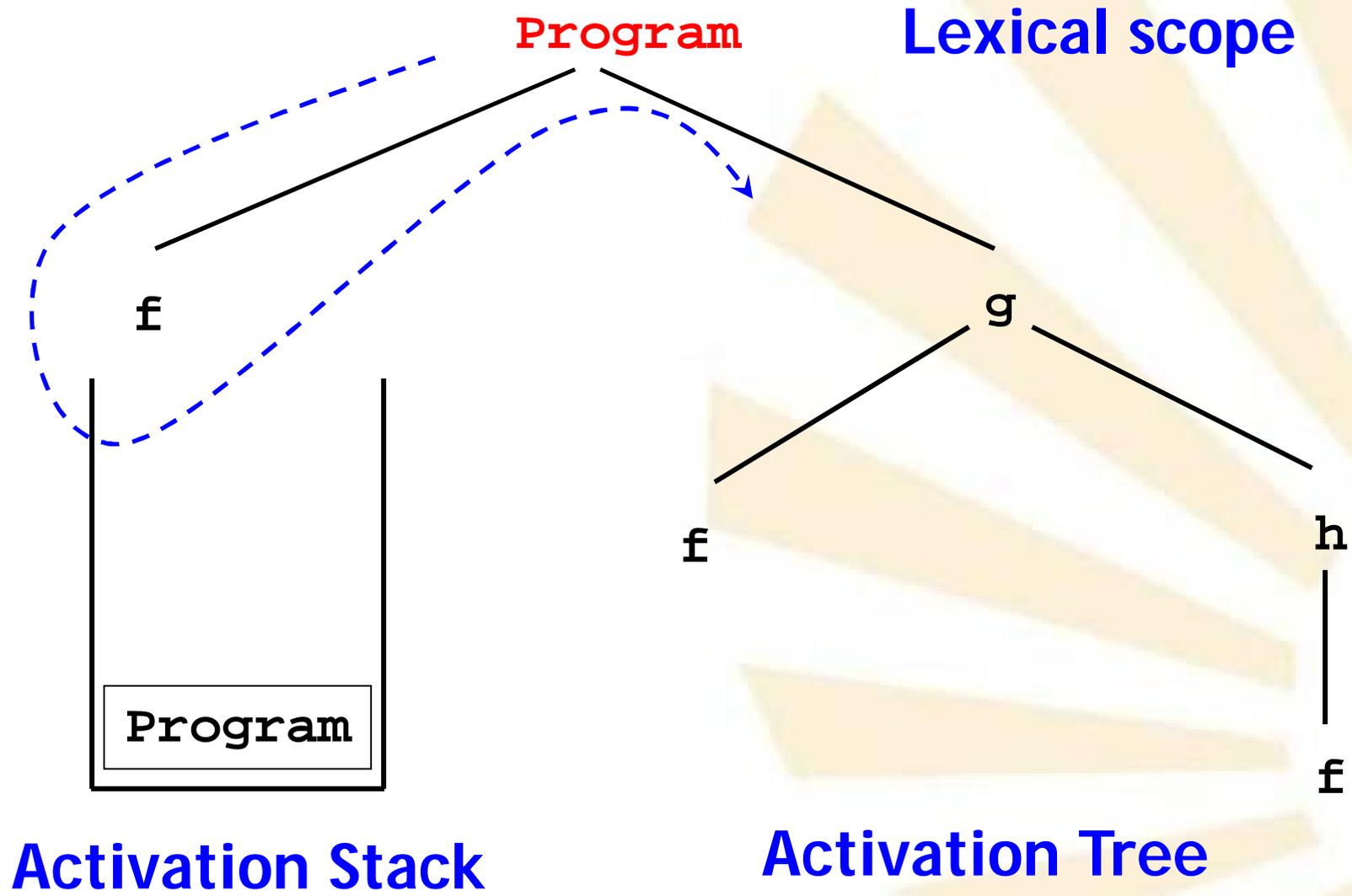
## Lexical scope

- Use of *x1* of *Program* in *f*

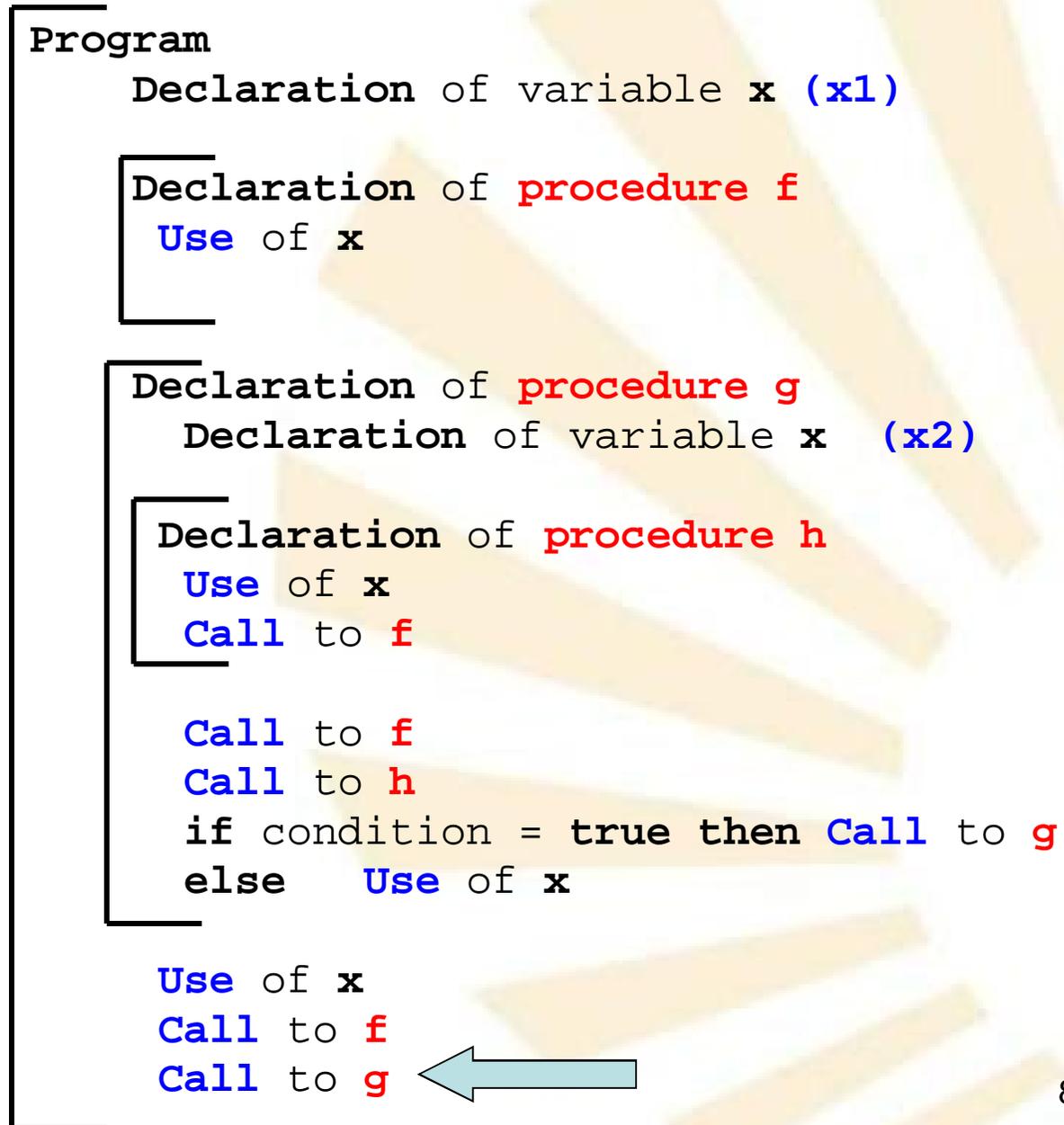


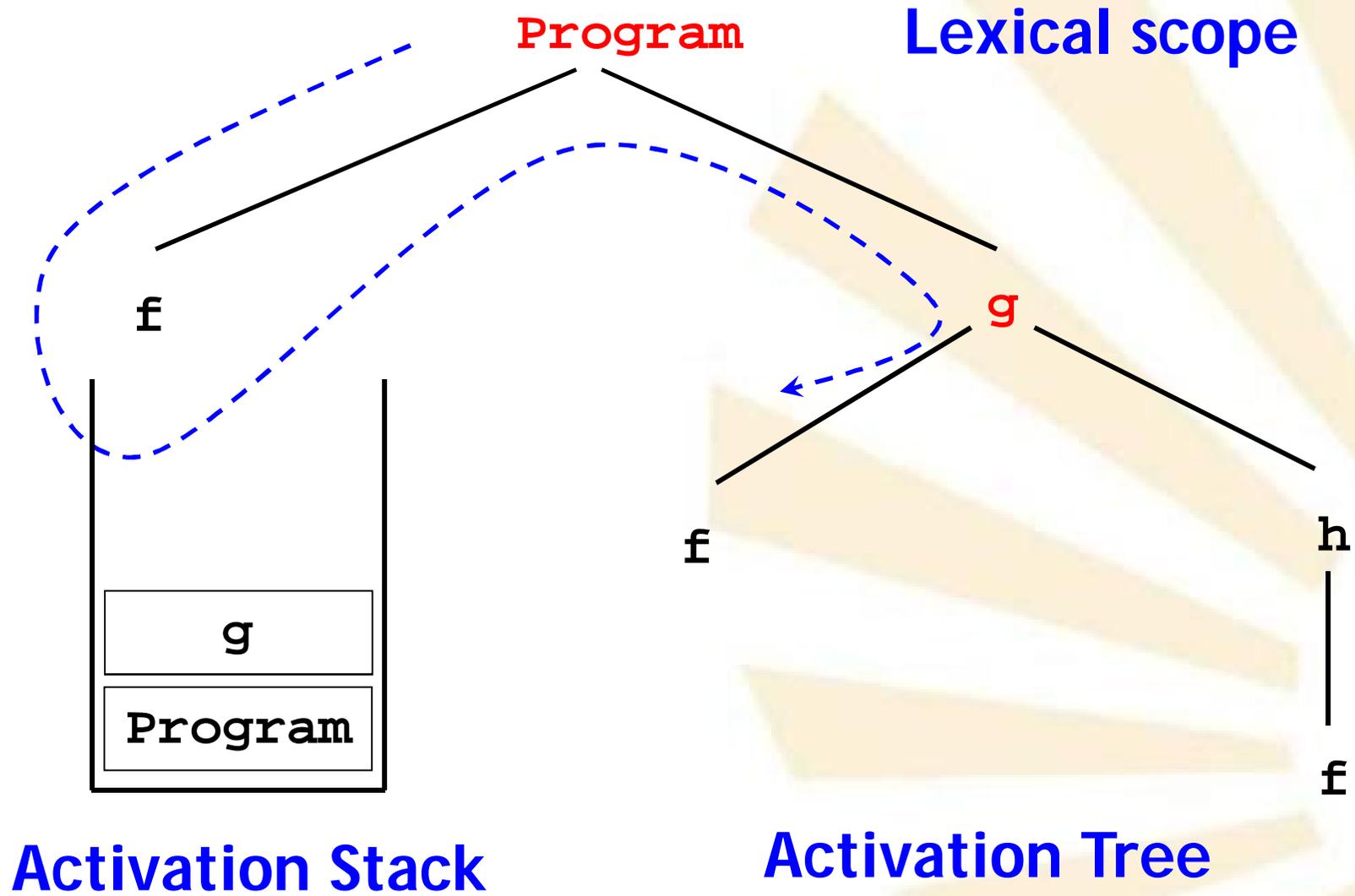
Activation Stack

Activation Tree

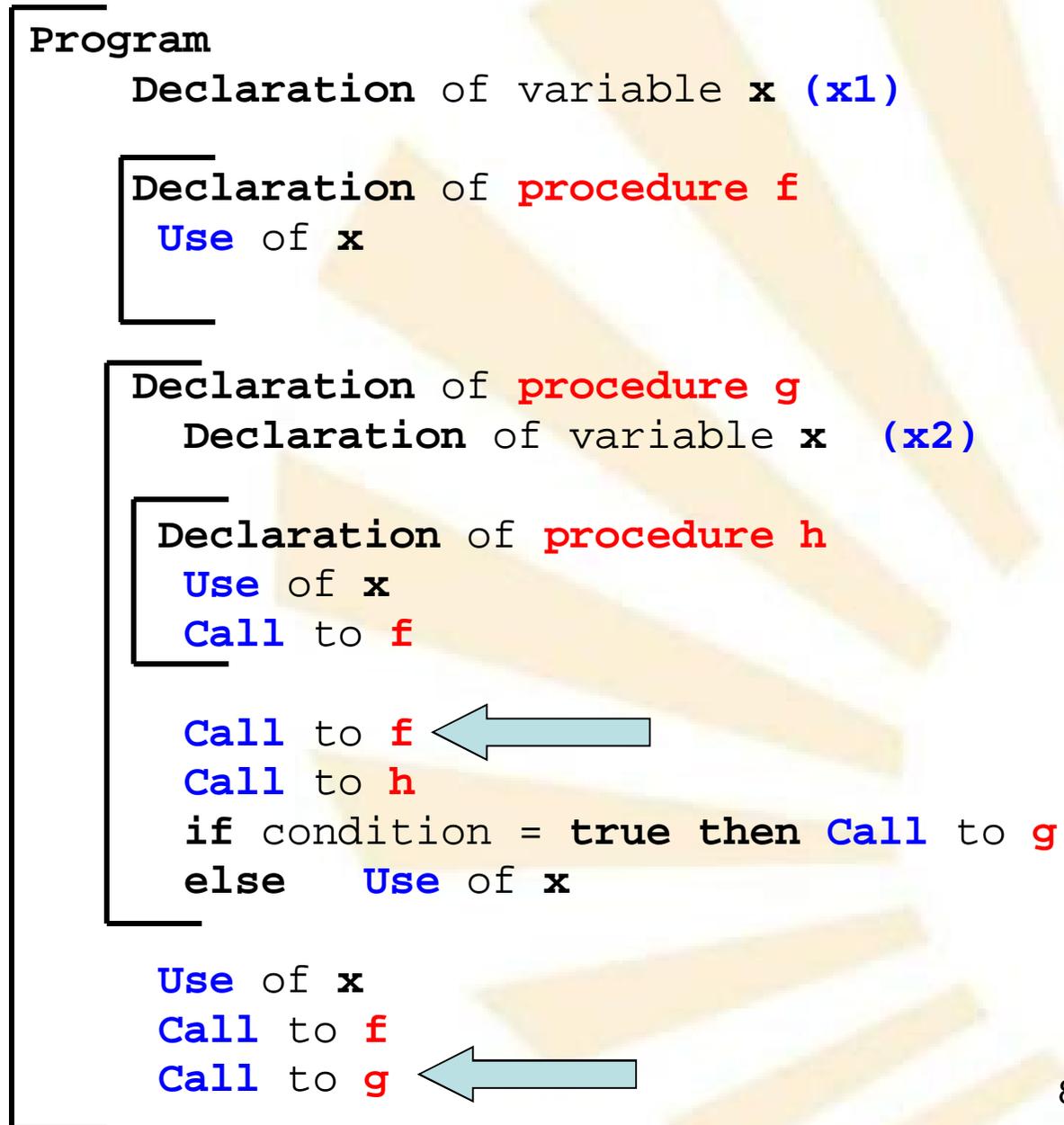


# Run with lexical scope

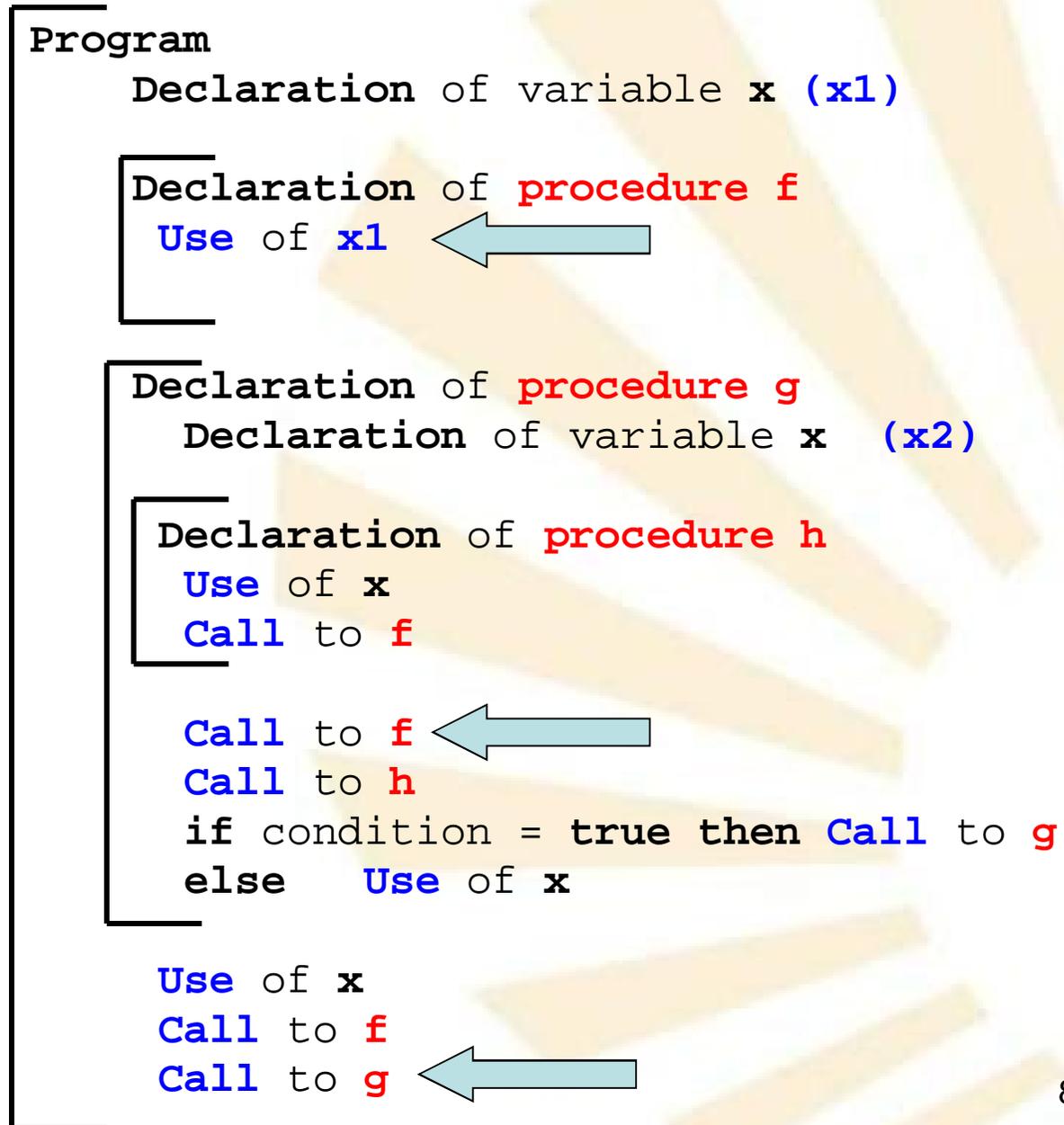


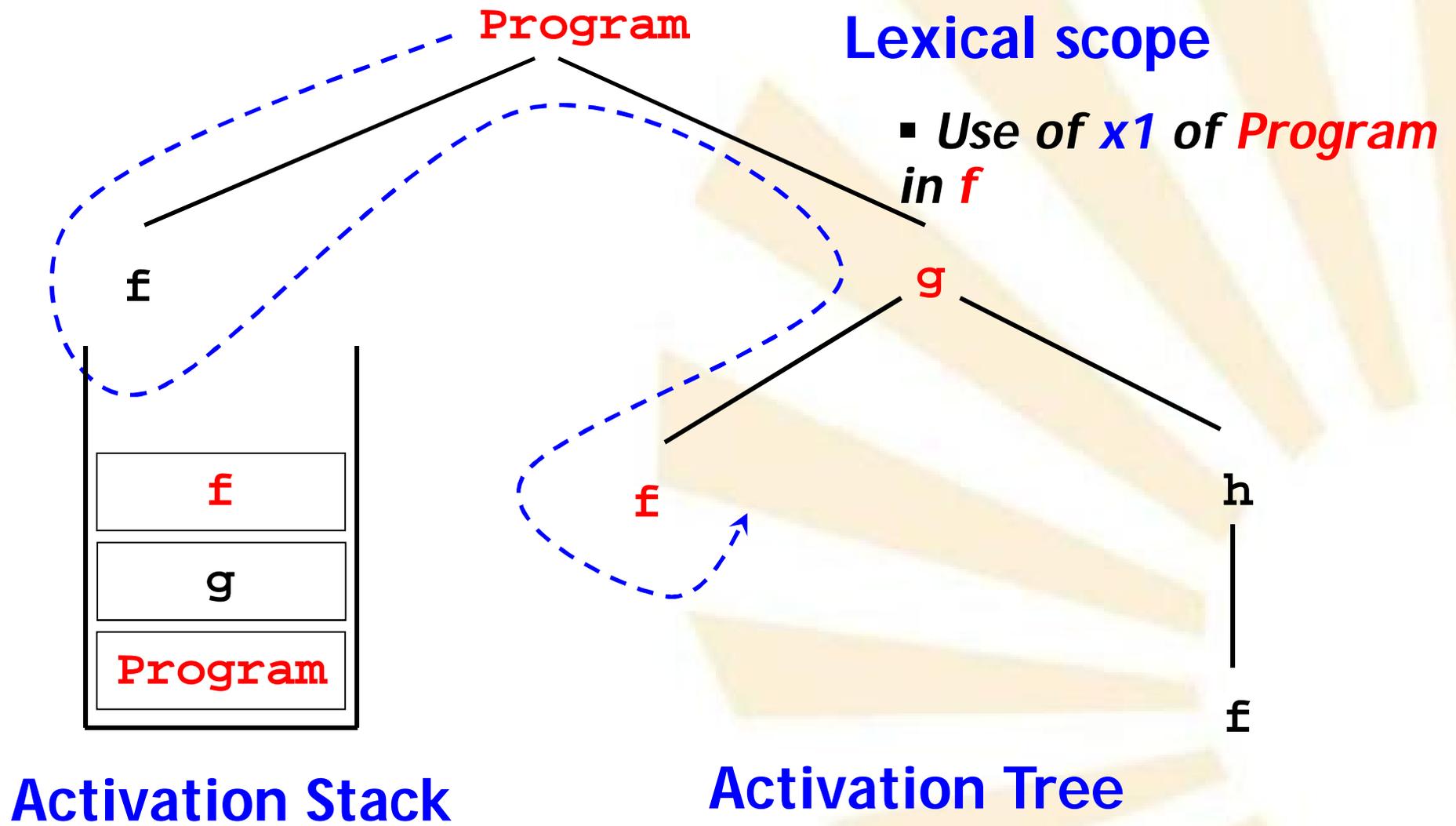


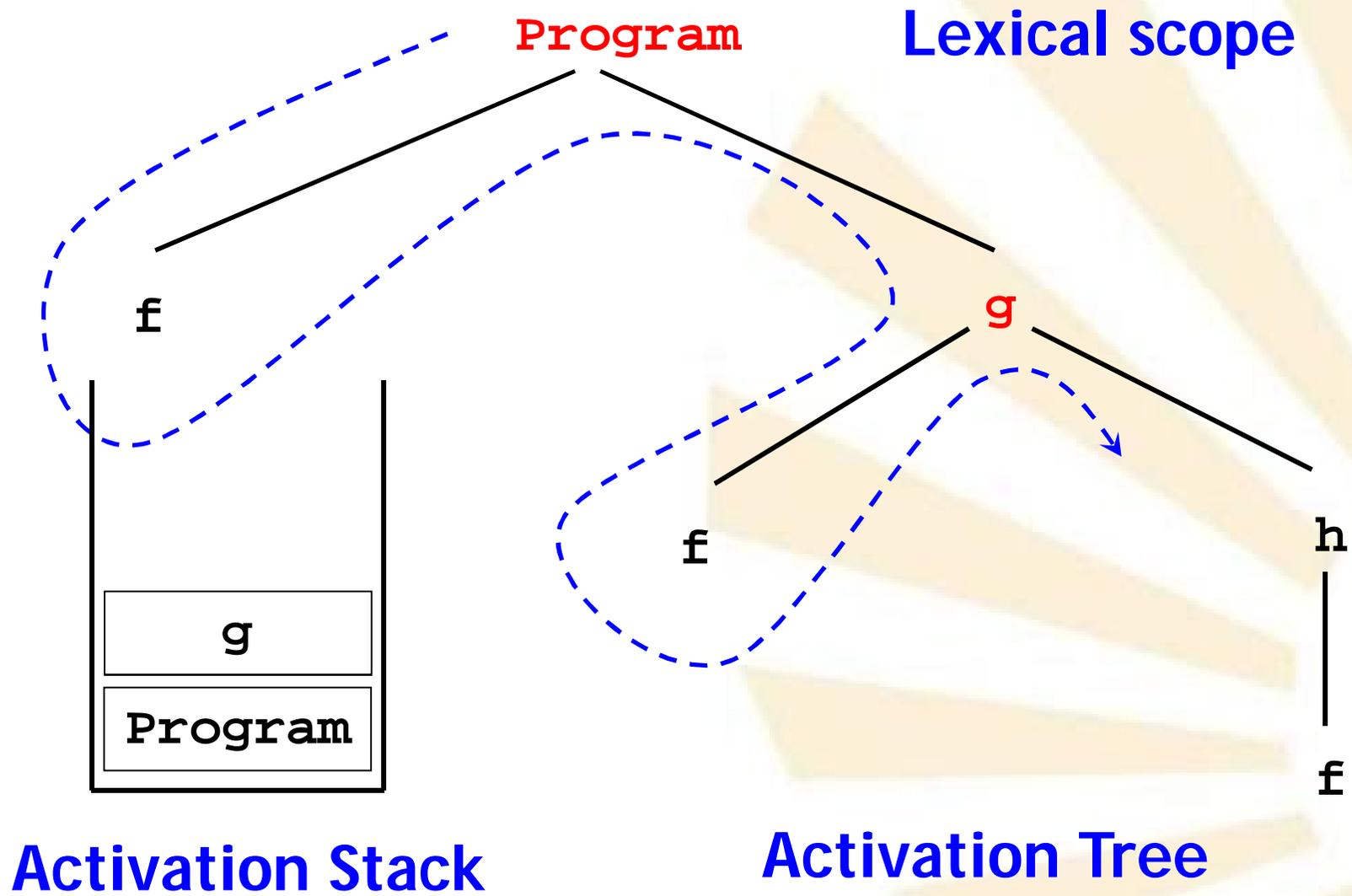
# Run with lexical scope



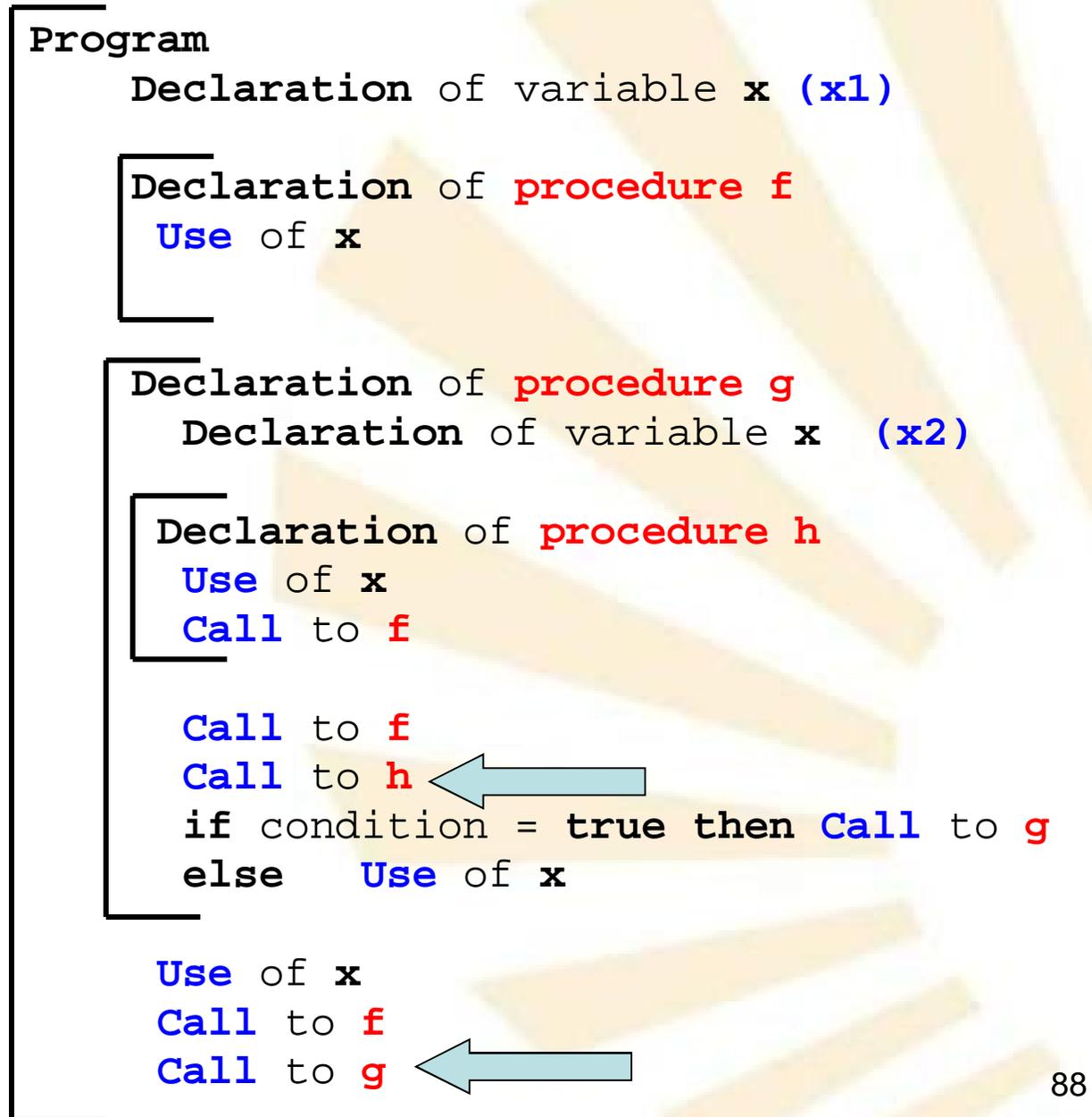
# Run with lexical scope



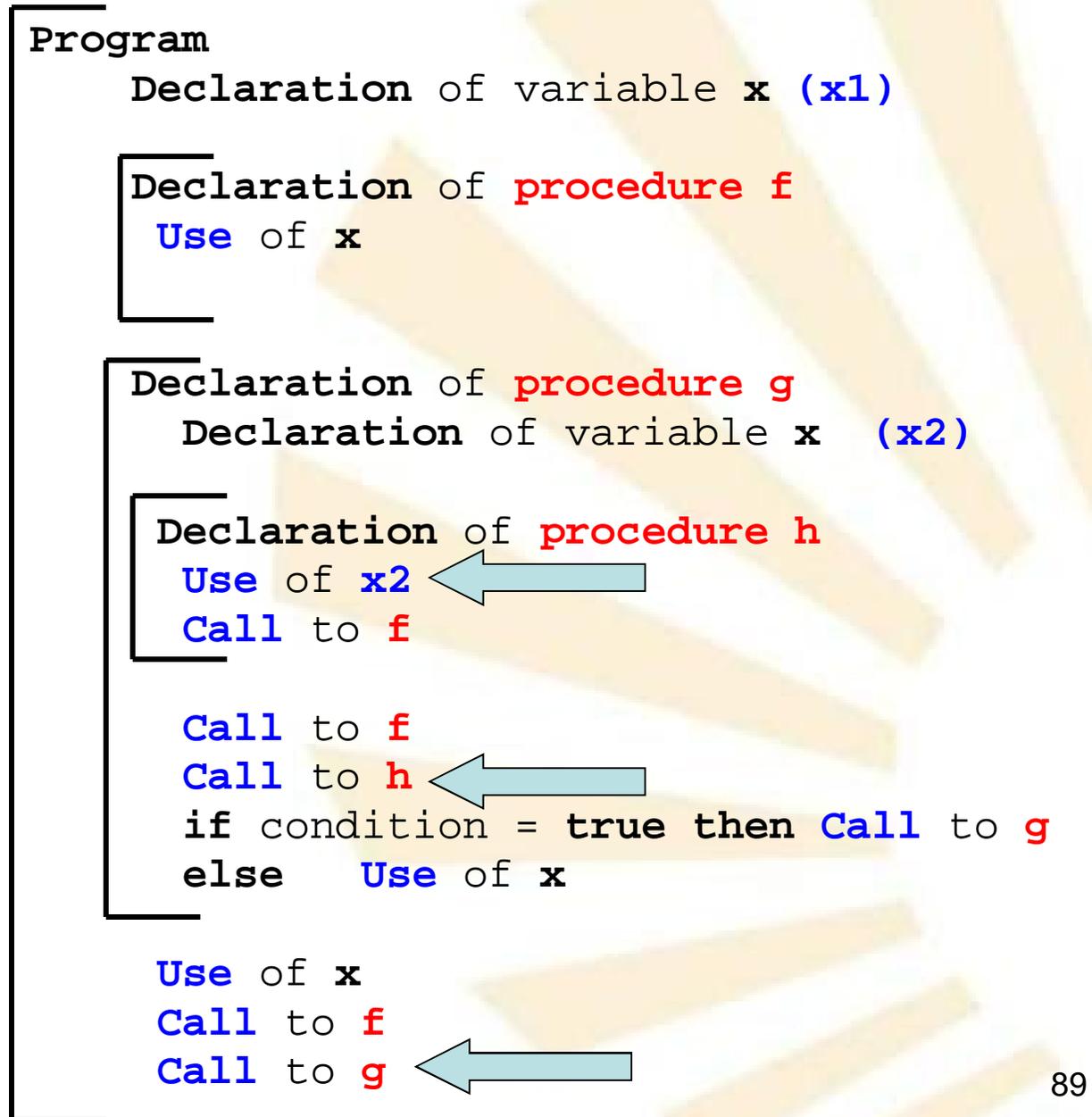




# Run with lexical scope



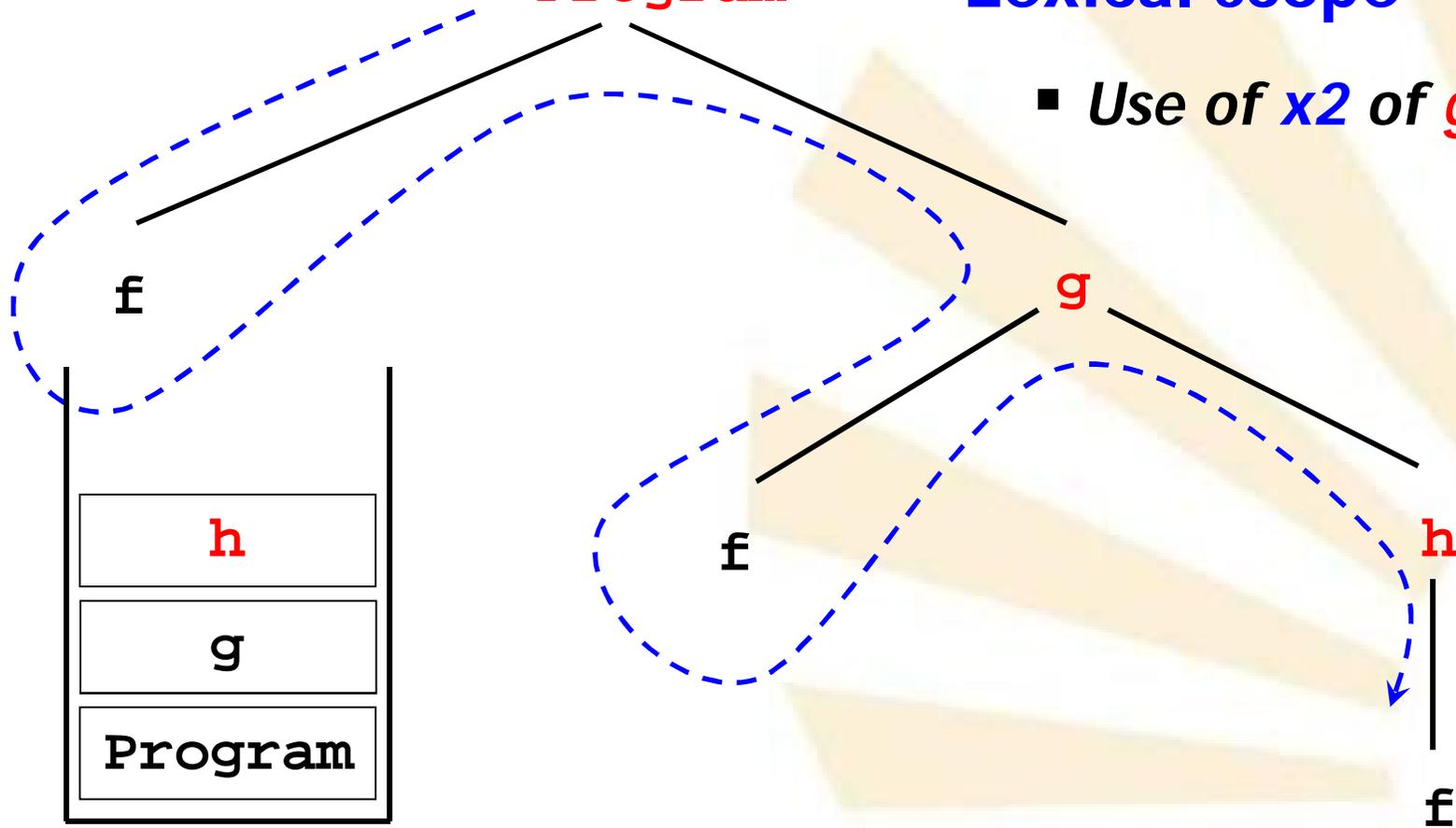
# Run with lexical scope



Program

Lexical scope

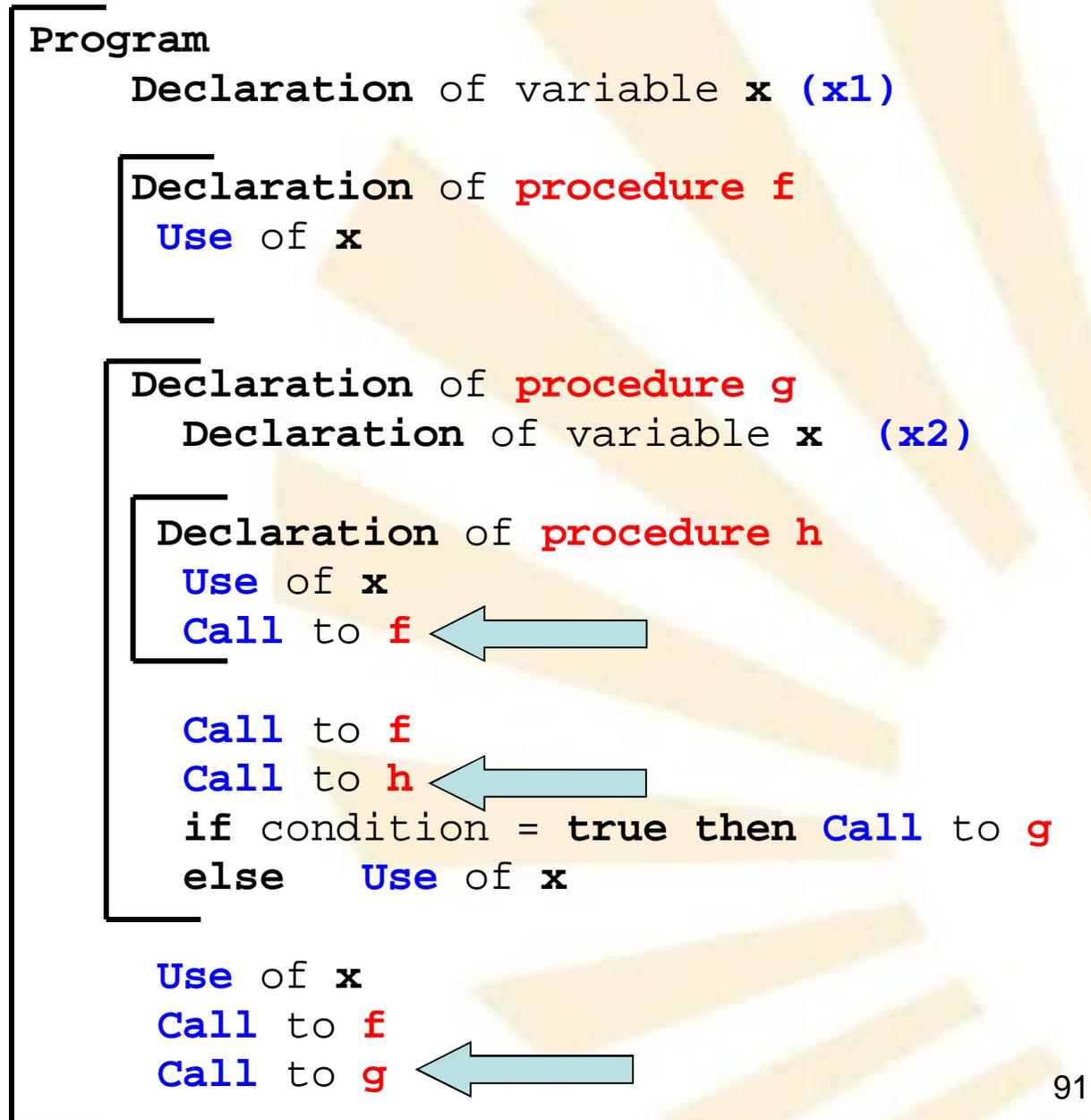
- Use of  $x_2$  of  $g$  in  $h$



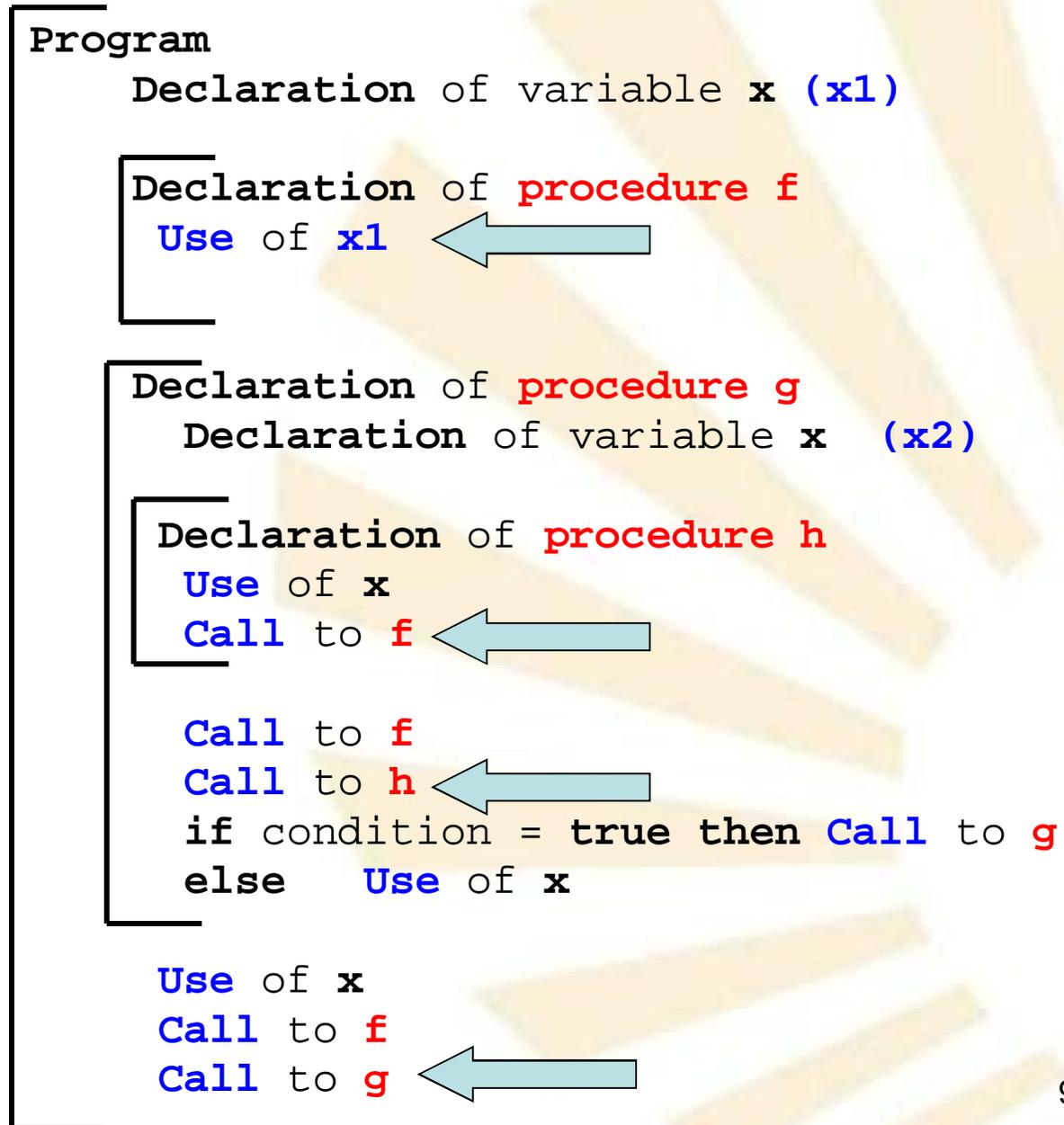
Activation Stack

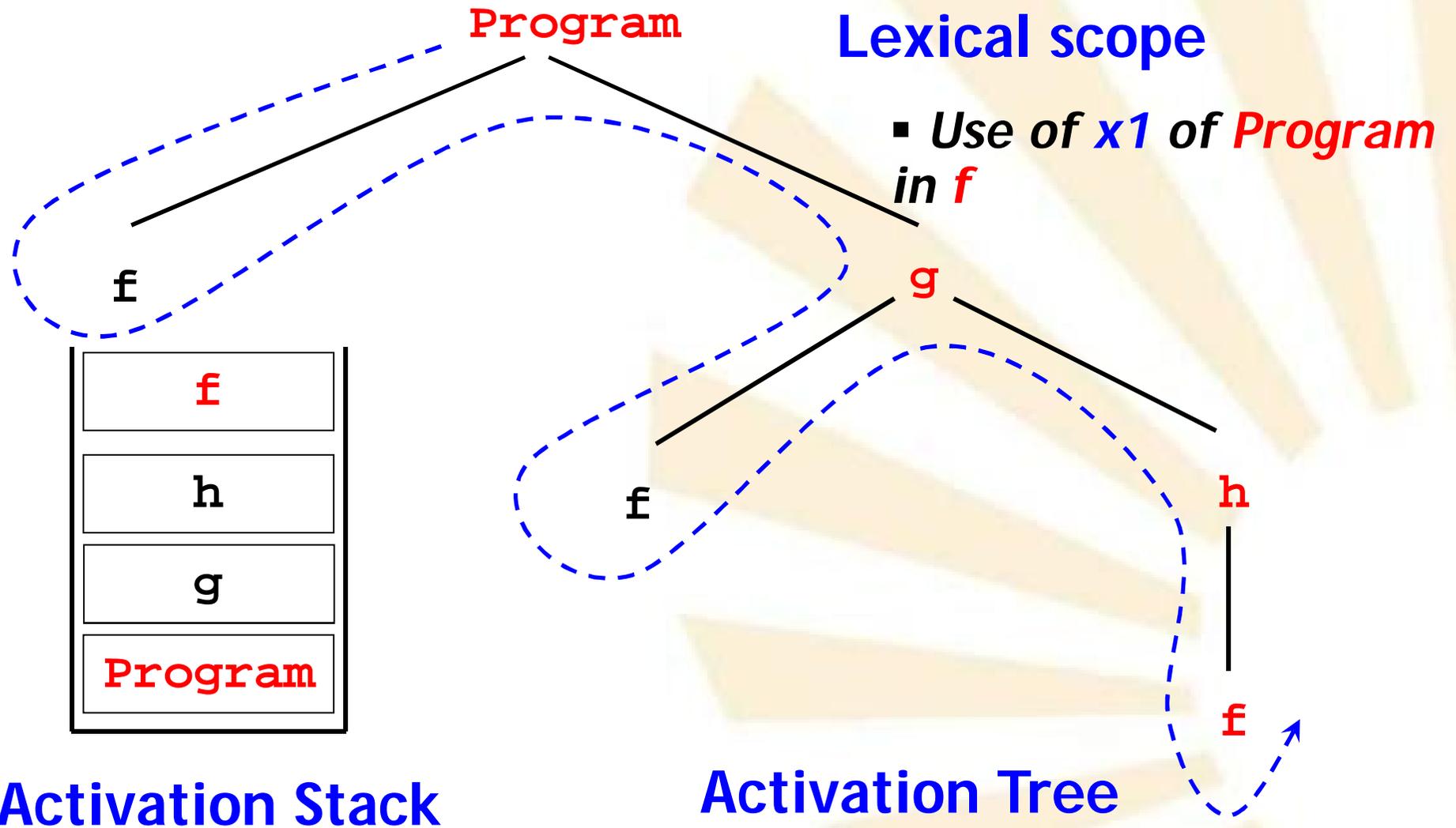
Activation Tree

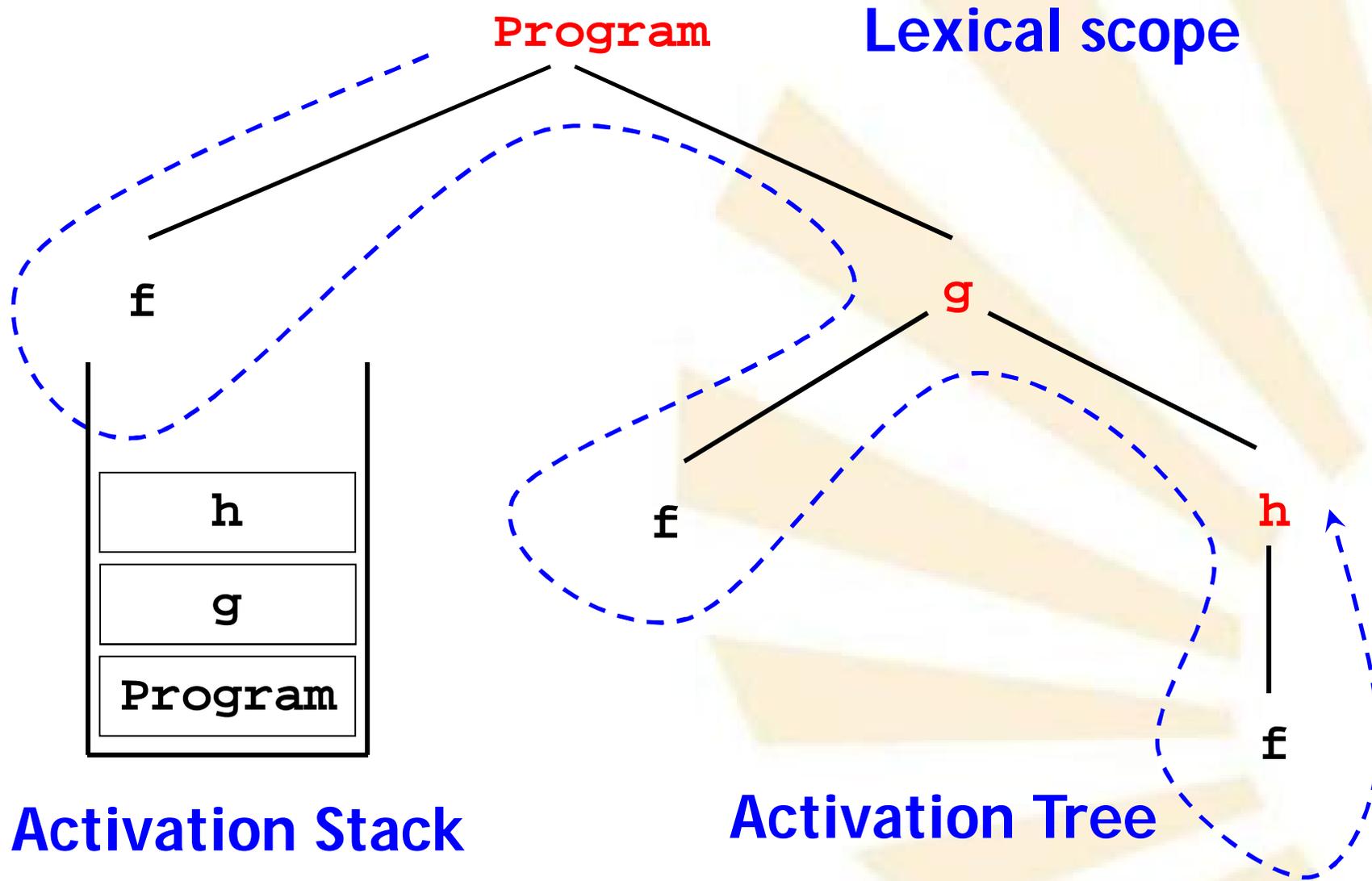
# Run with lexical scope



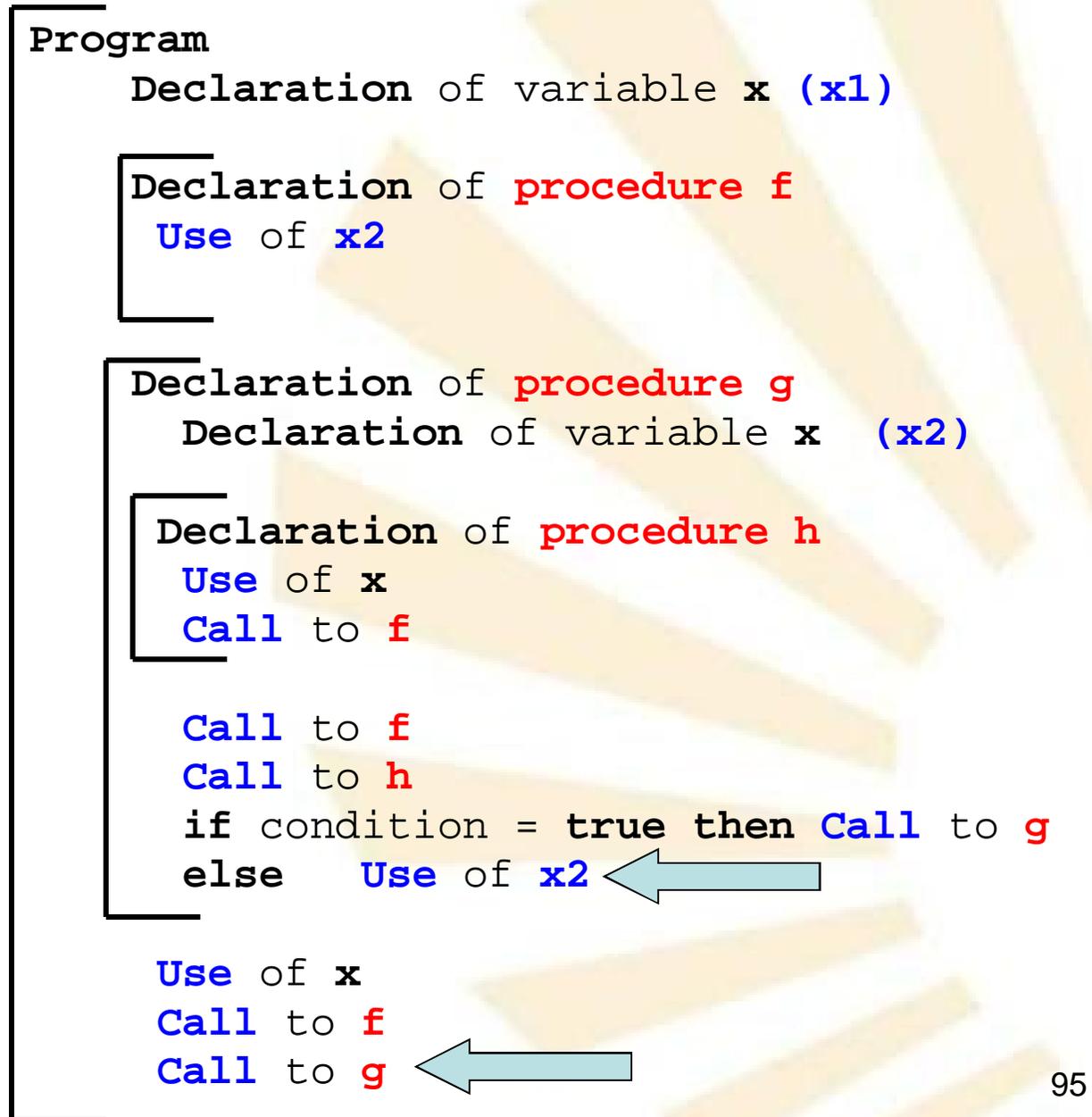
# Run with lexical scope





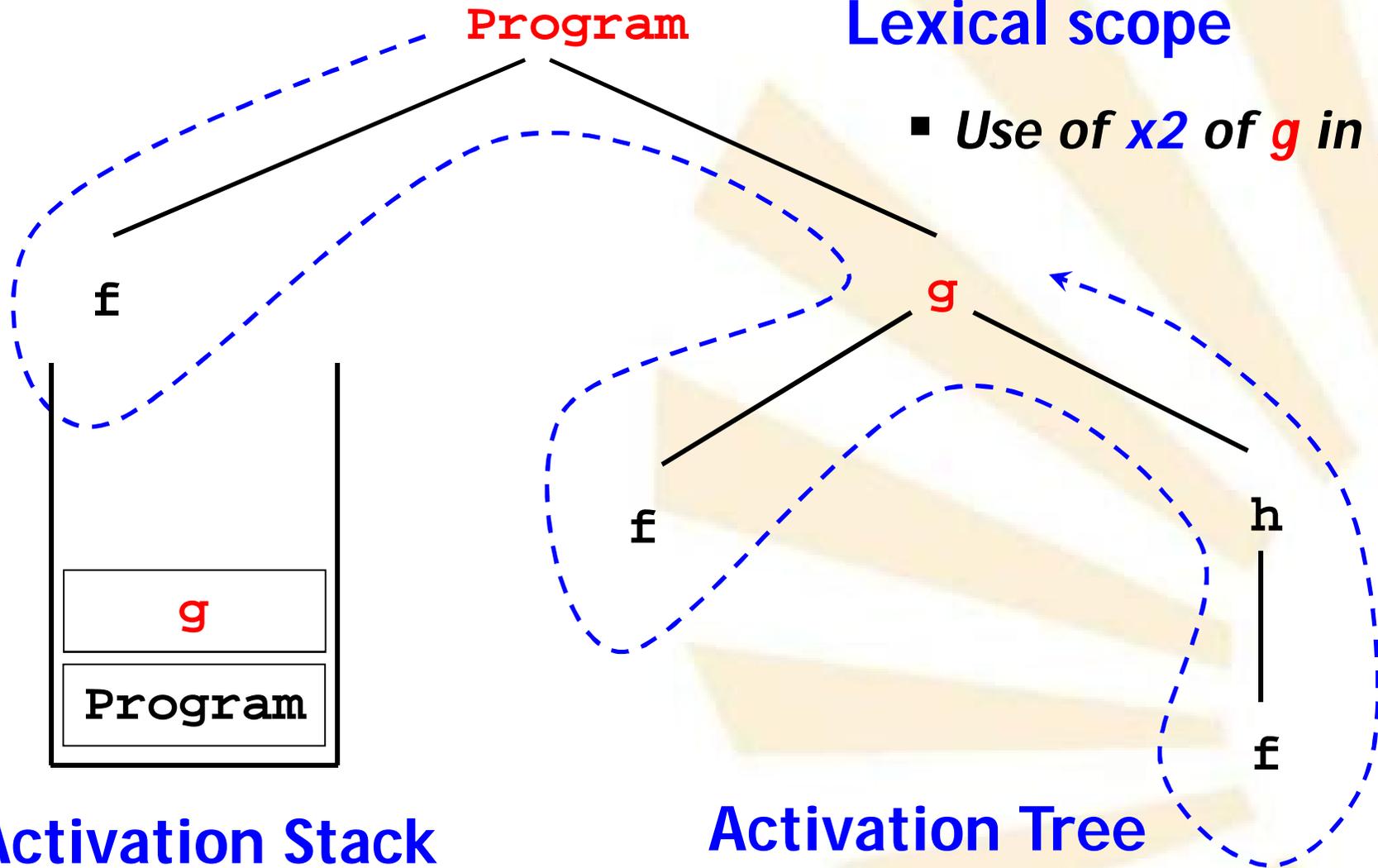


# Run with lexical scope



## Lexical scope

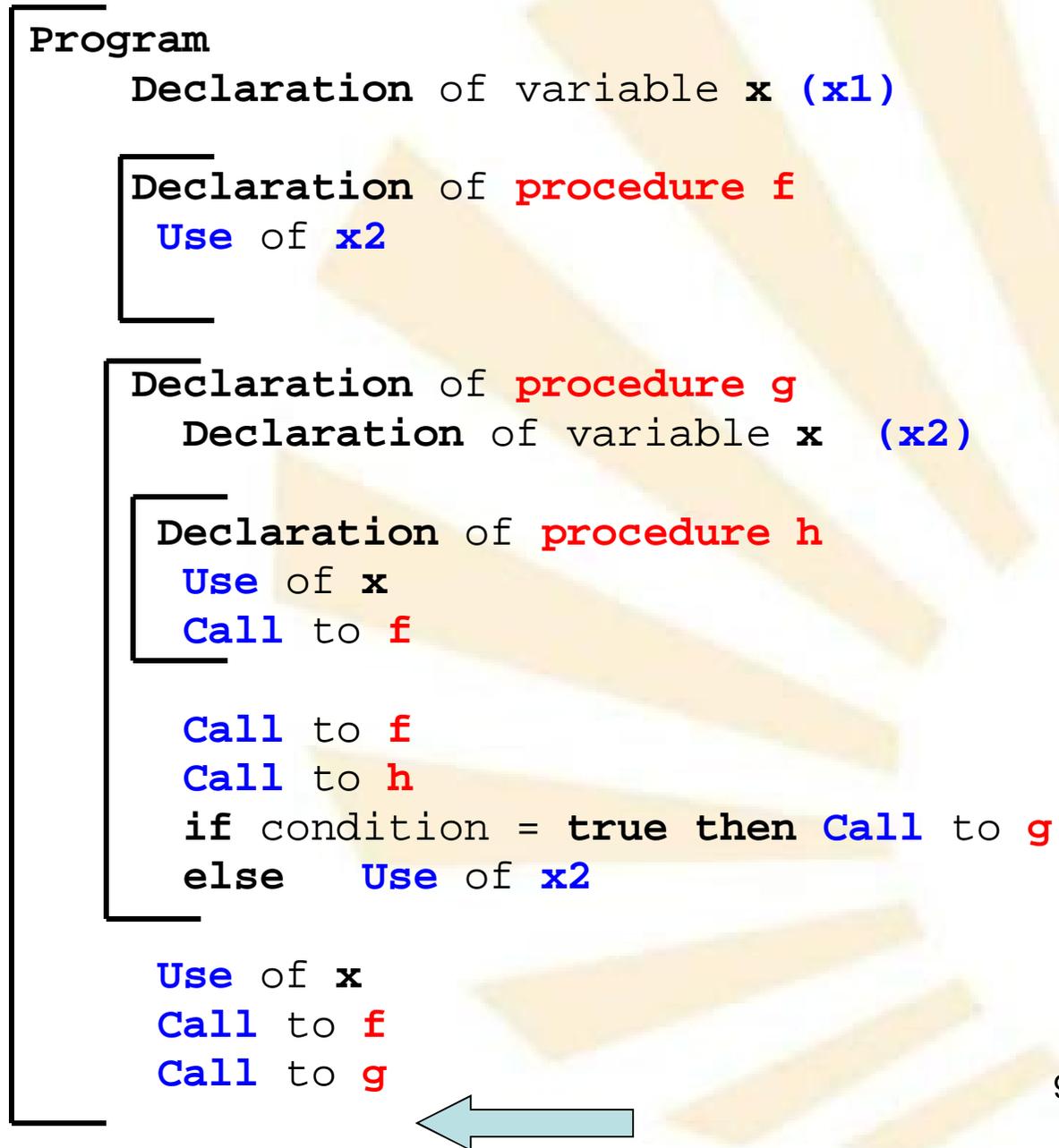
- Use of  $x_2$  of  $g$  in  $g$

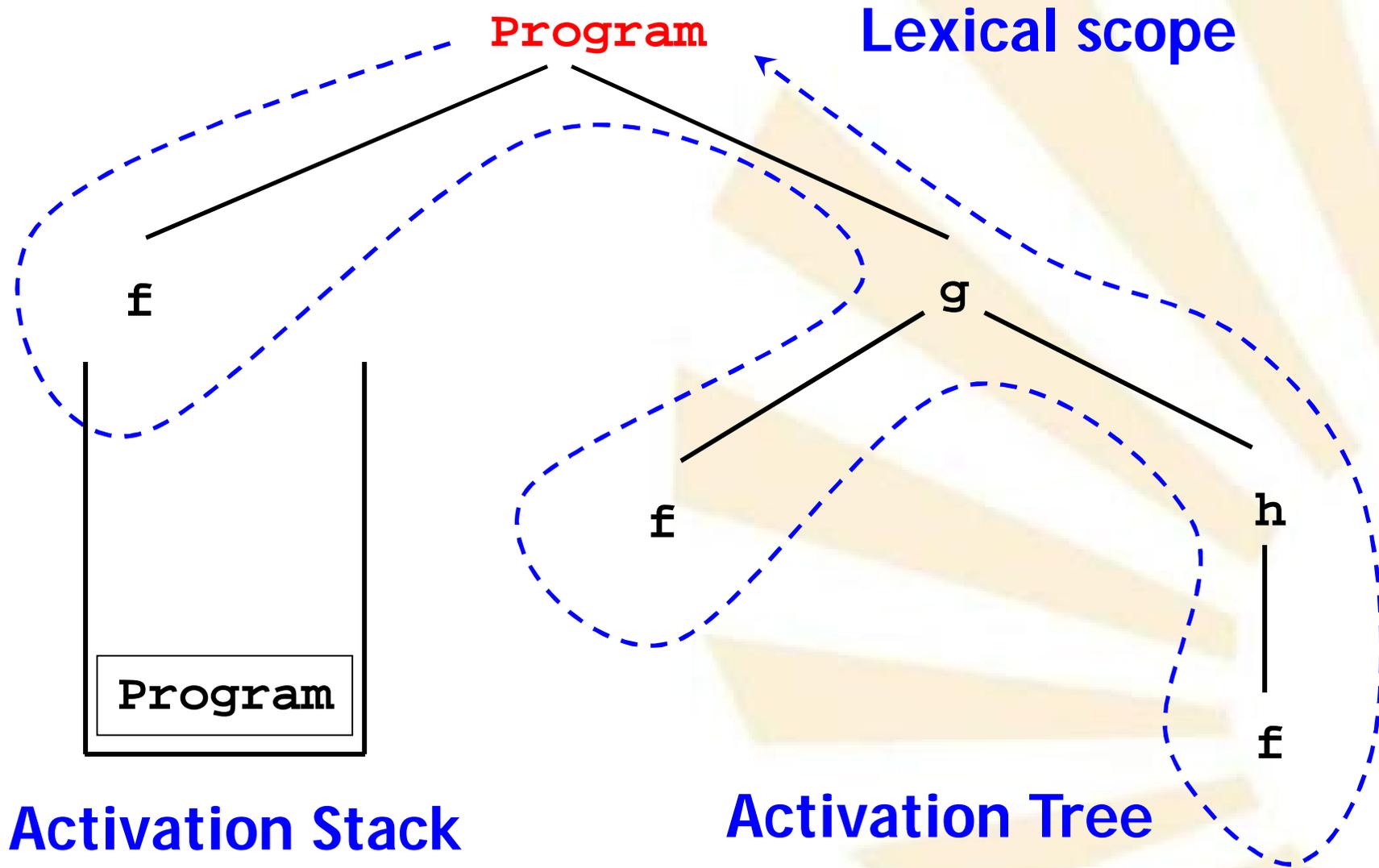


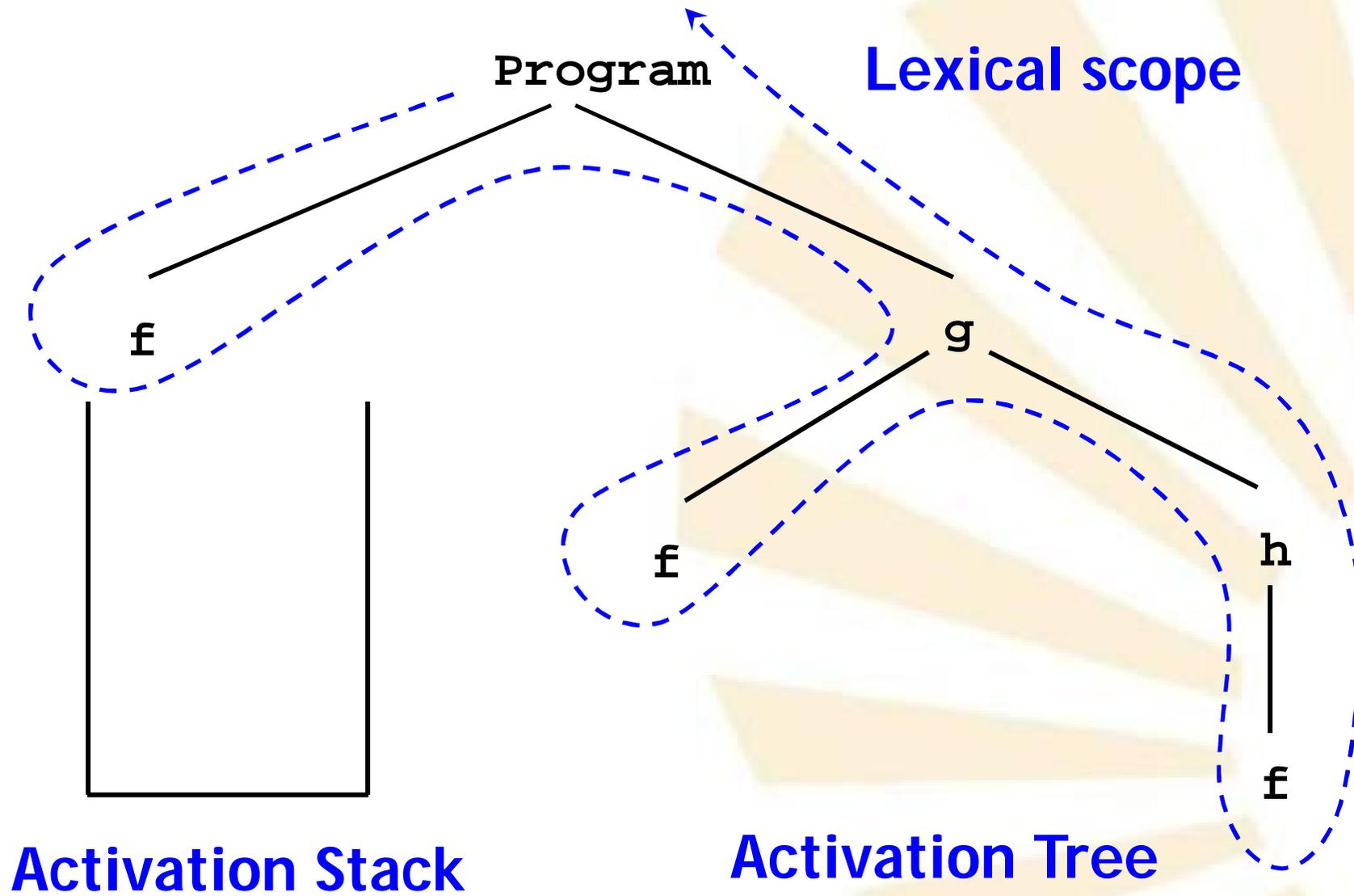
Activation Stack

Activation Tree

# Run with lexical scope







# Run with dynamical scope

Program

Declaration of variable **x** (**x<sub>1</sub>**)

Declaration of **procedure f**

Use of **x**

Declaration of **procedure g**

Declaration of variable **x** (**x<sub>2</sub>**)

Declaration of **procedure h**

Use of **x**

Call to **f**

Call to **f**

Call to **h**

if condition = true then Call to **g**

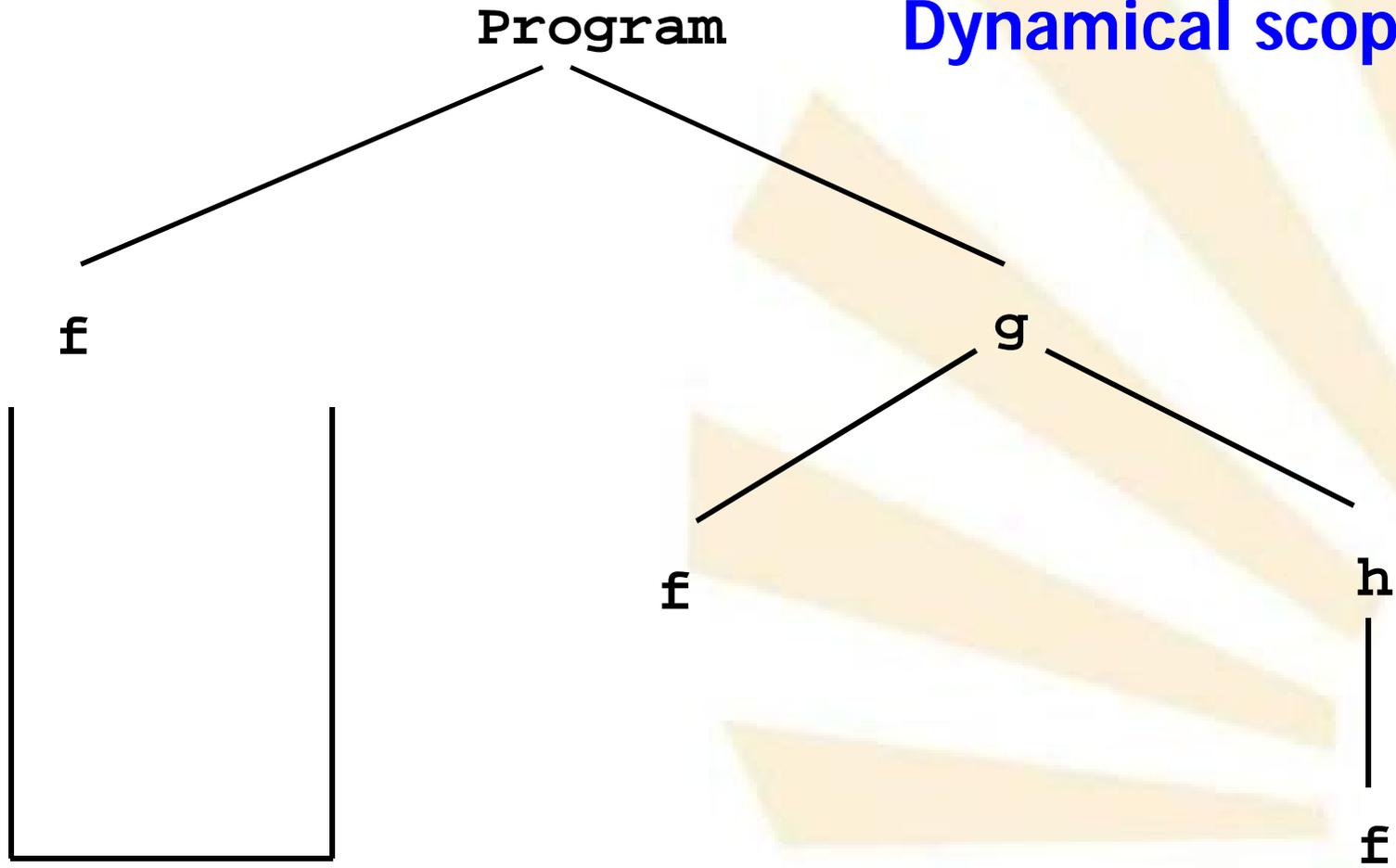
else Use of **x**

Use of **x**

Call to **f**

Call to **g**

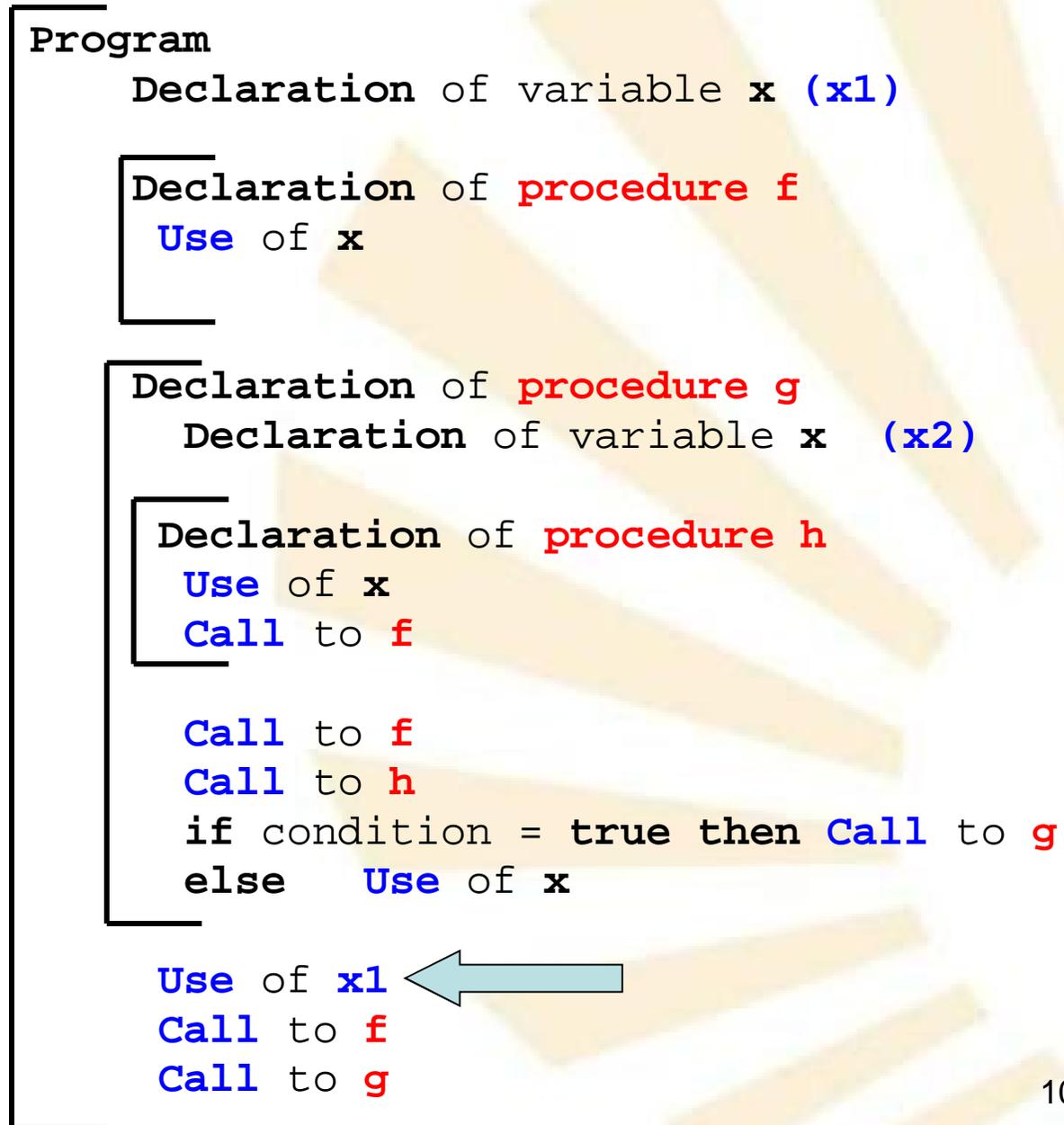
# Dynamical scope

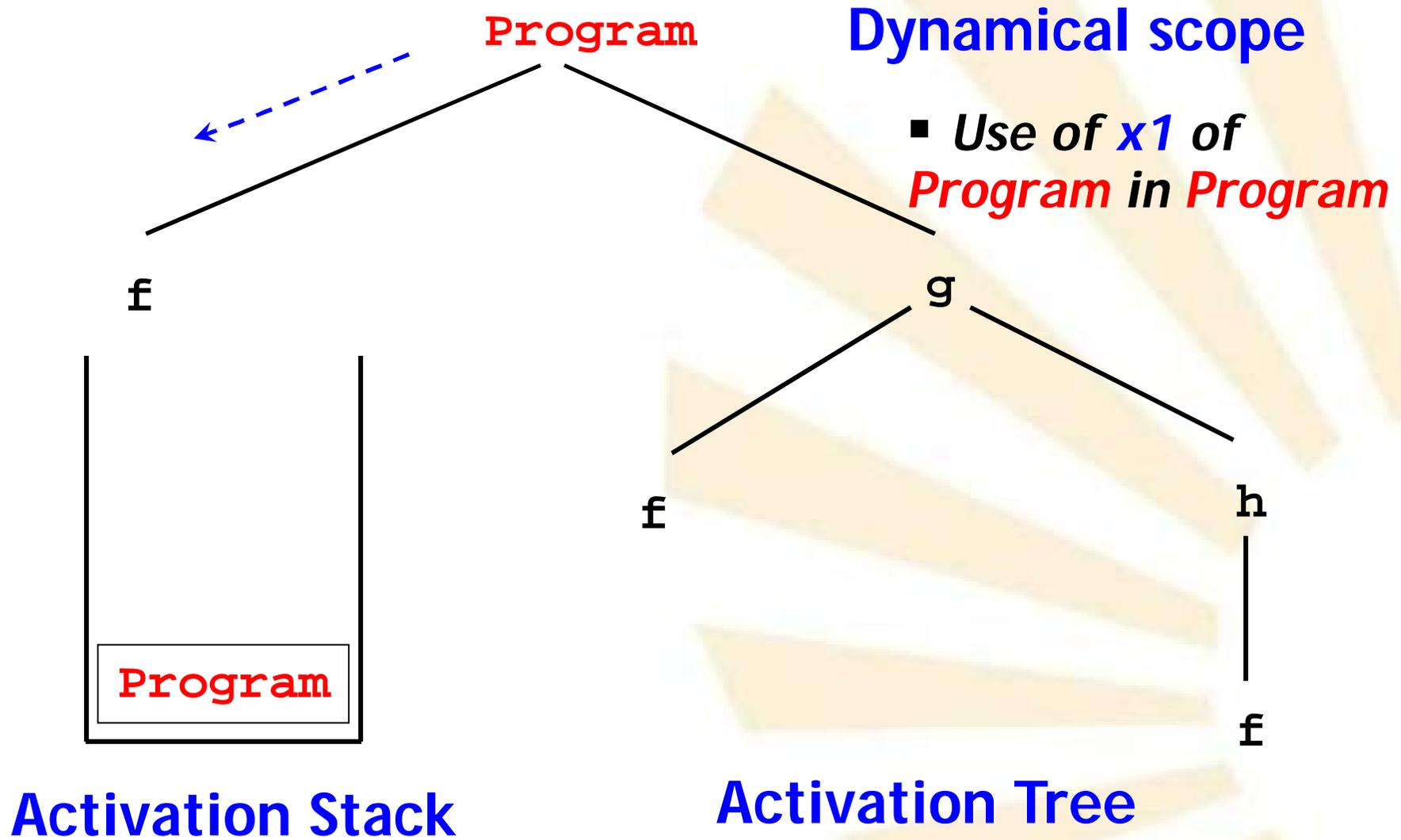


Activation Stack

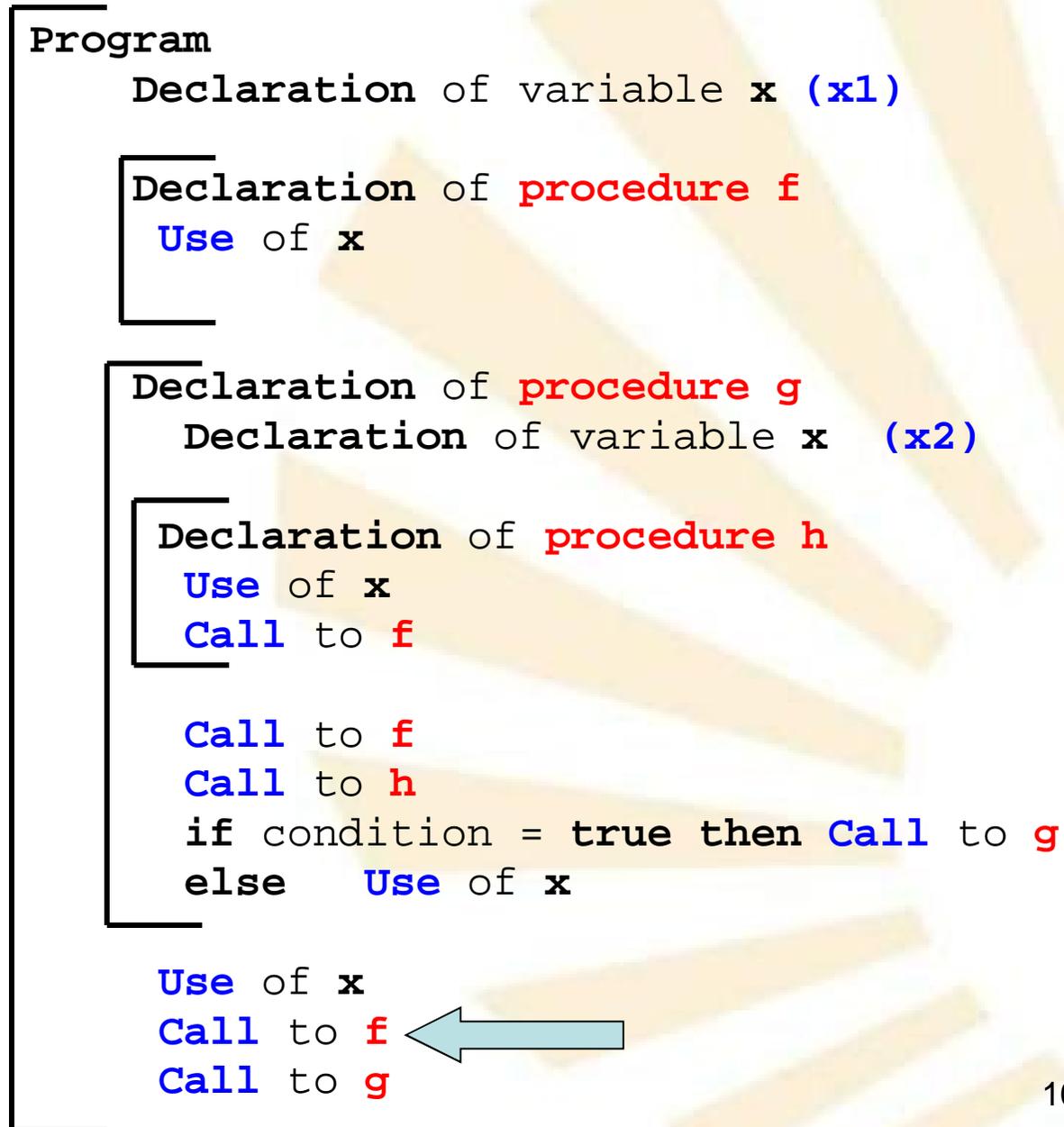
Activation Tree

Run with  
dynamical  
scope

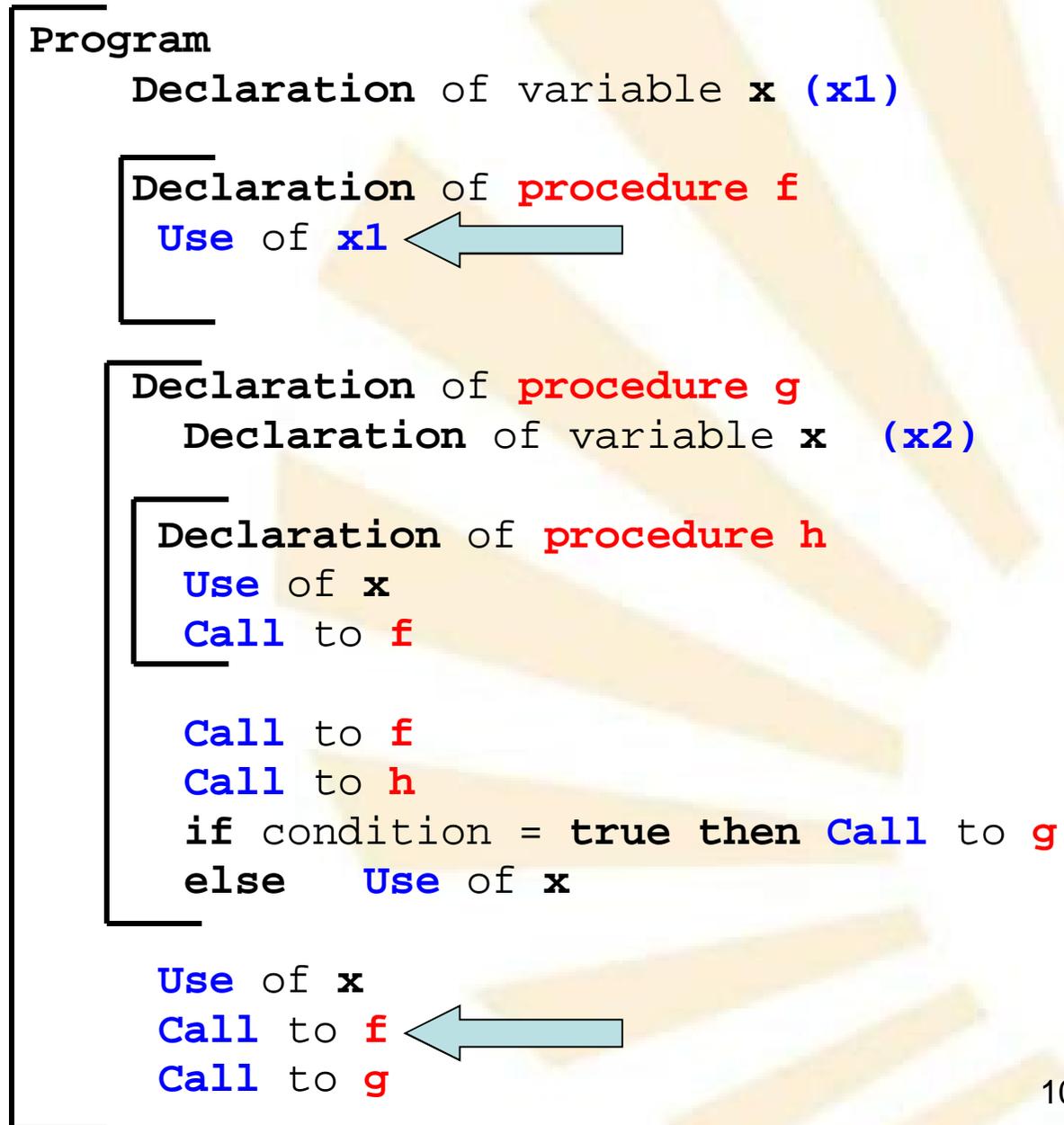


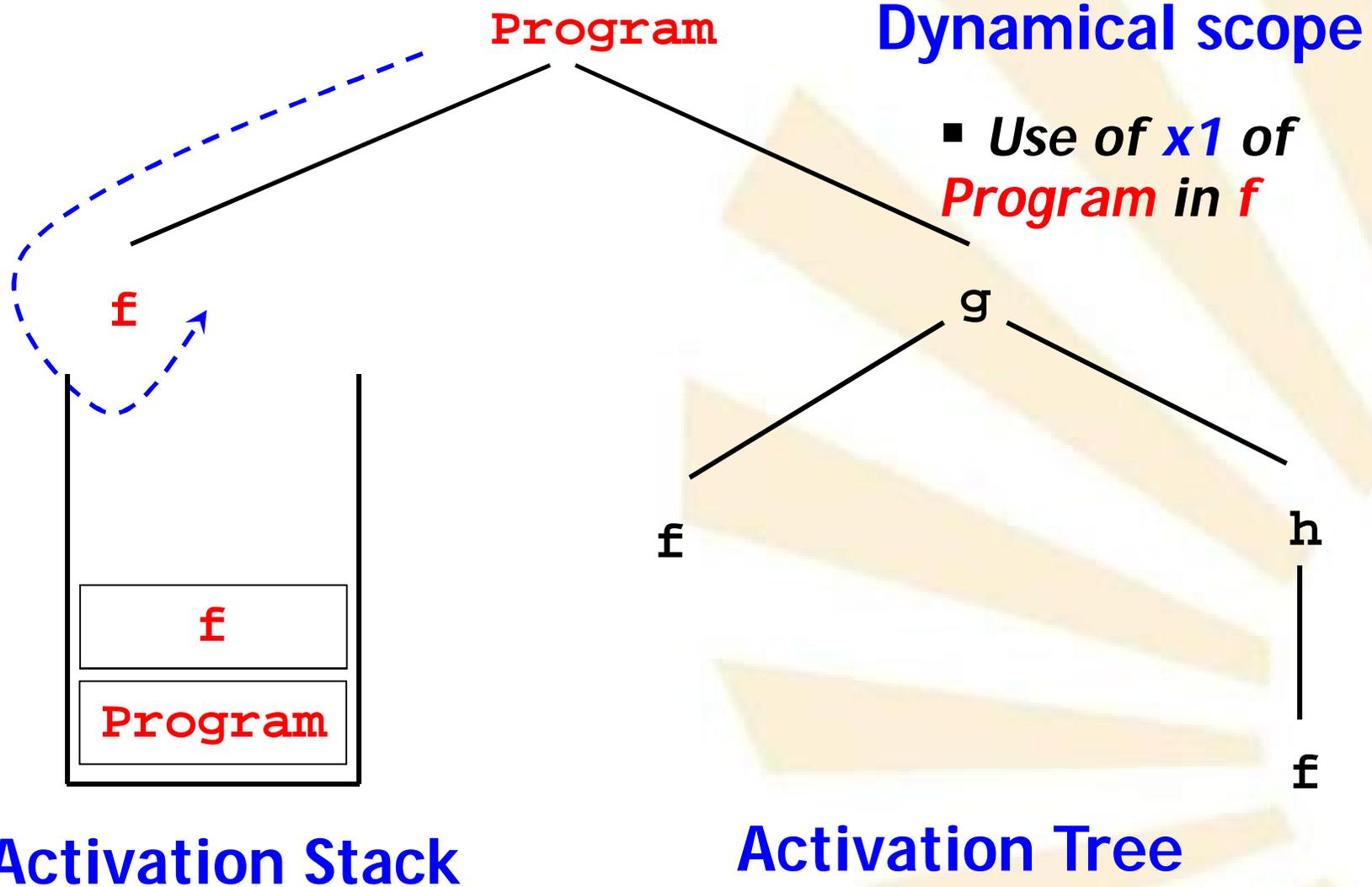


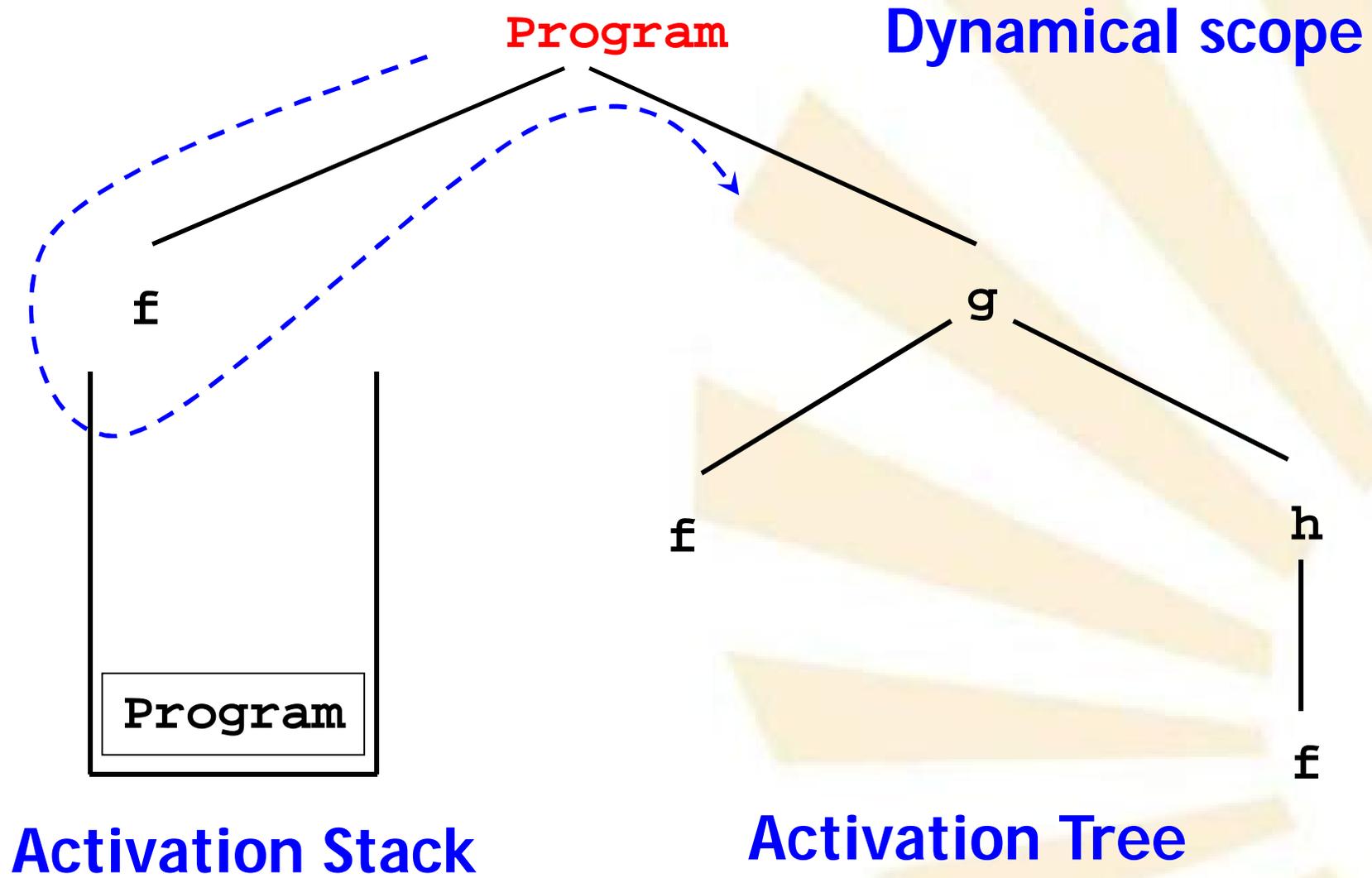
Run with  
dynamical  
scope



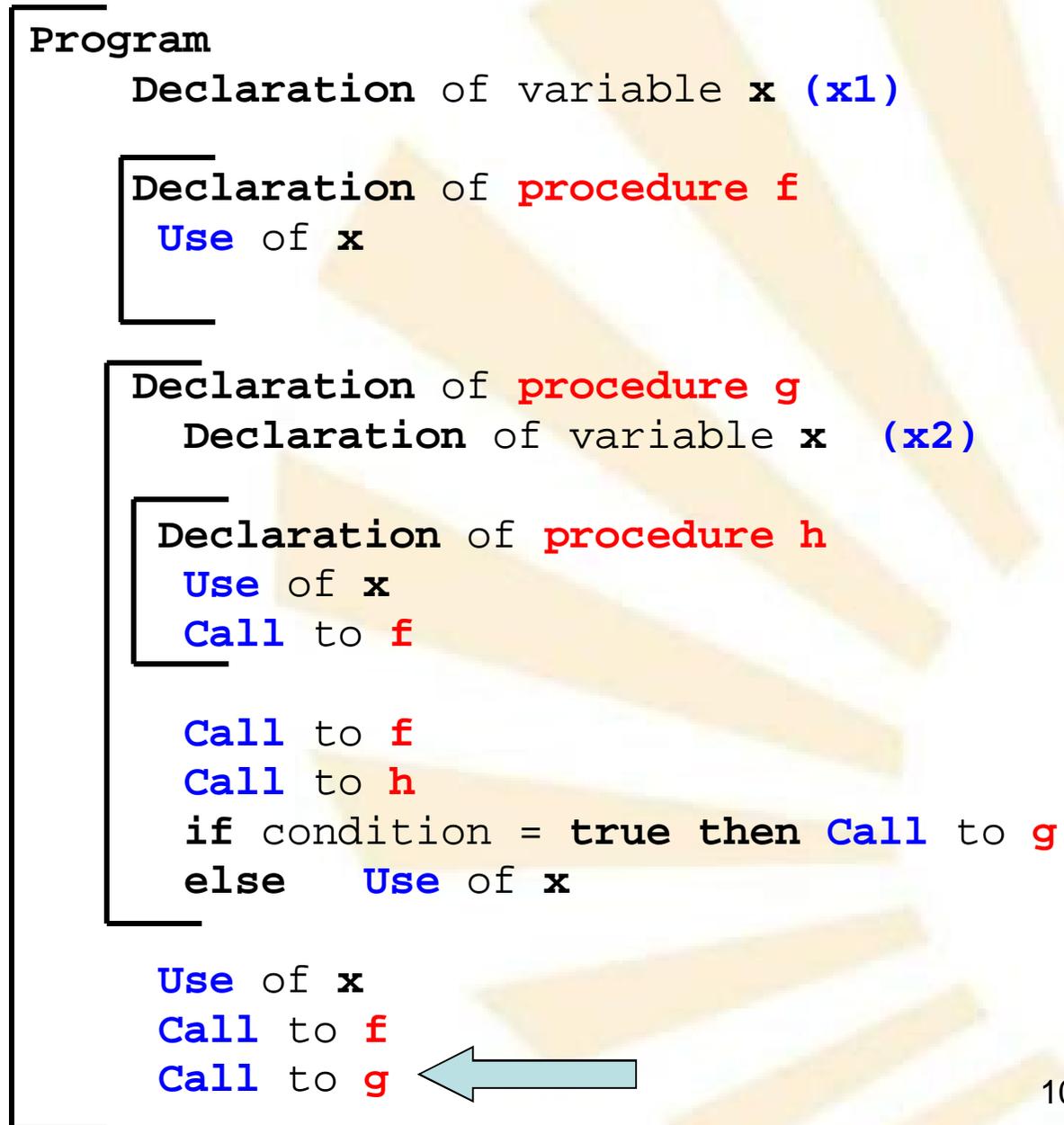
Run with  
dynamical  
scope

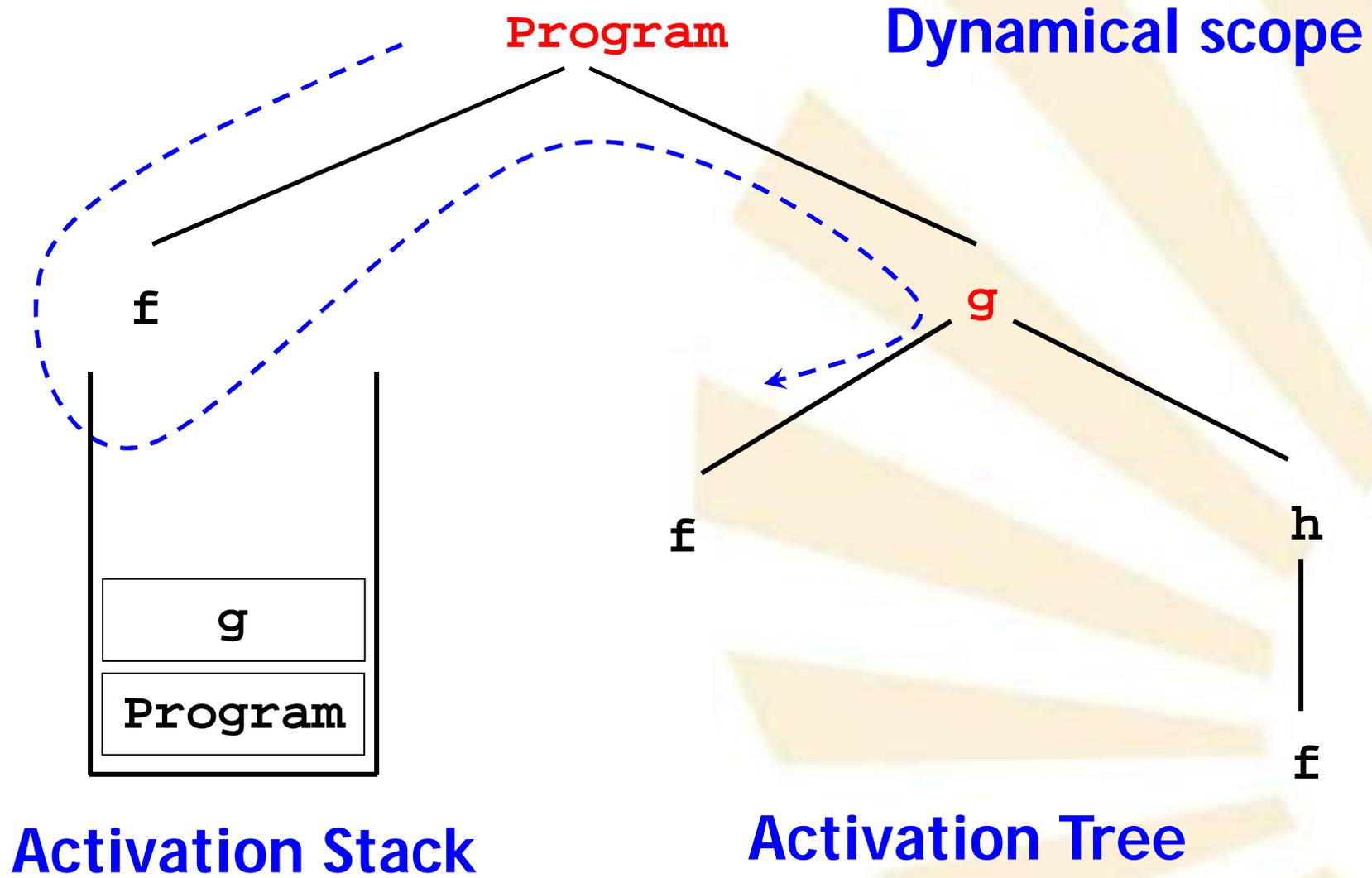




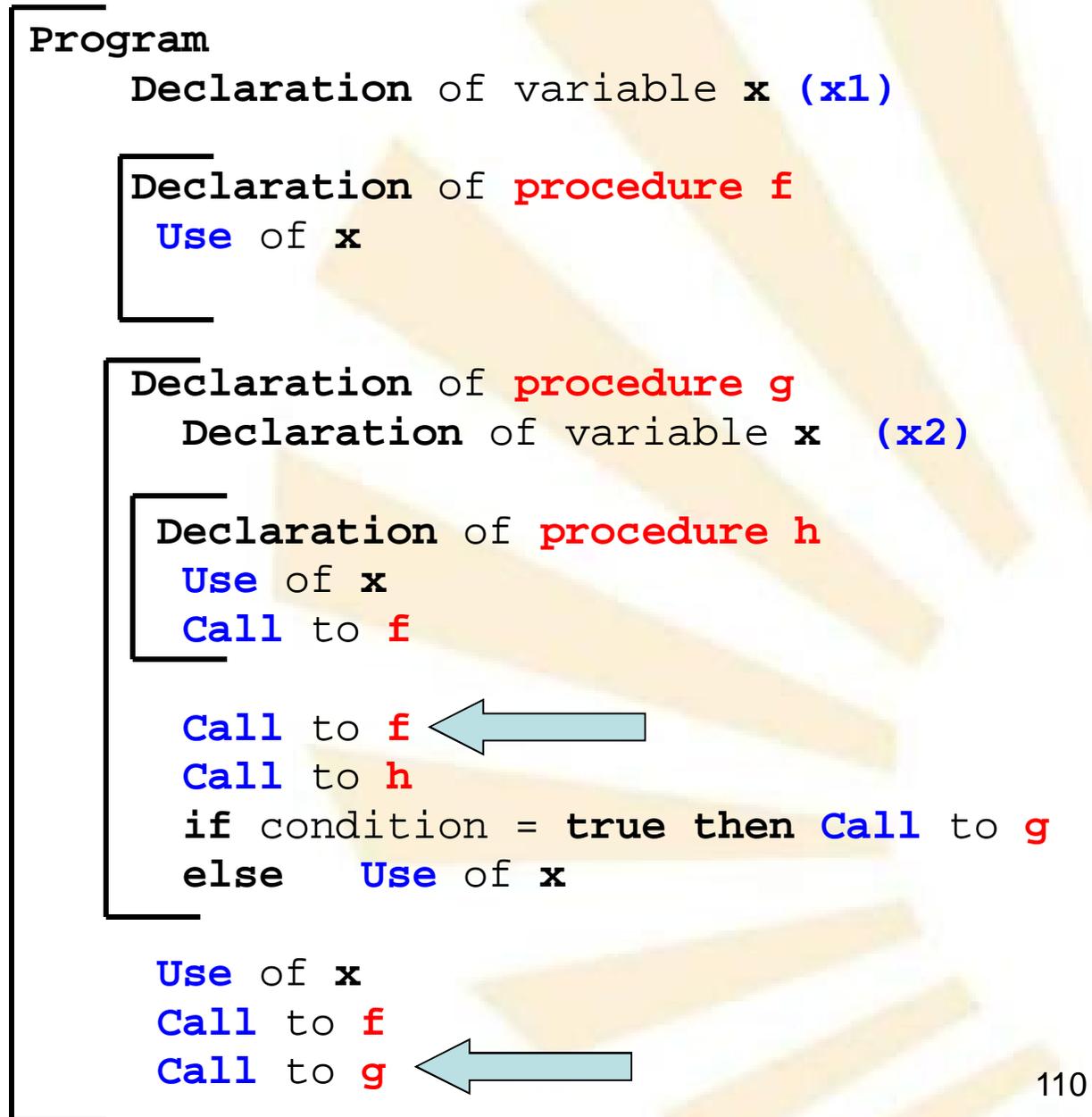


Run with  
dynamical  
scope

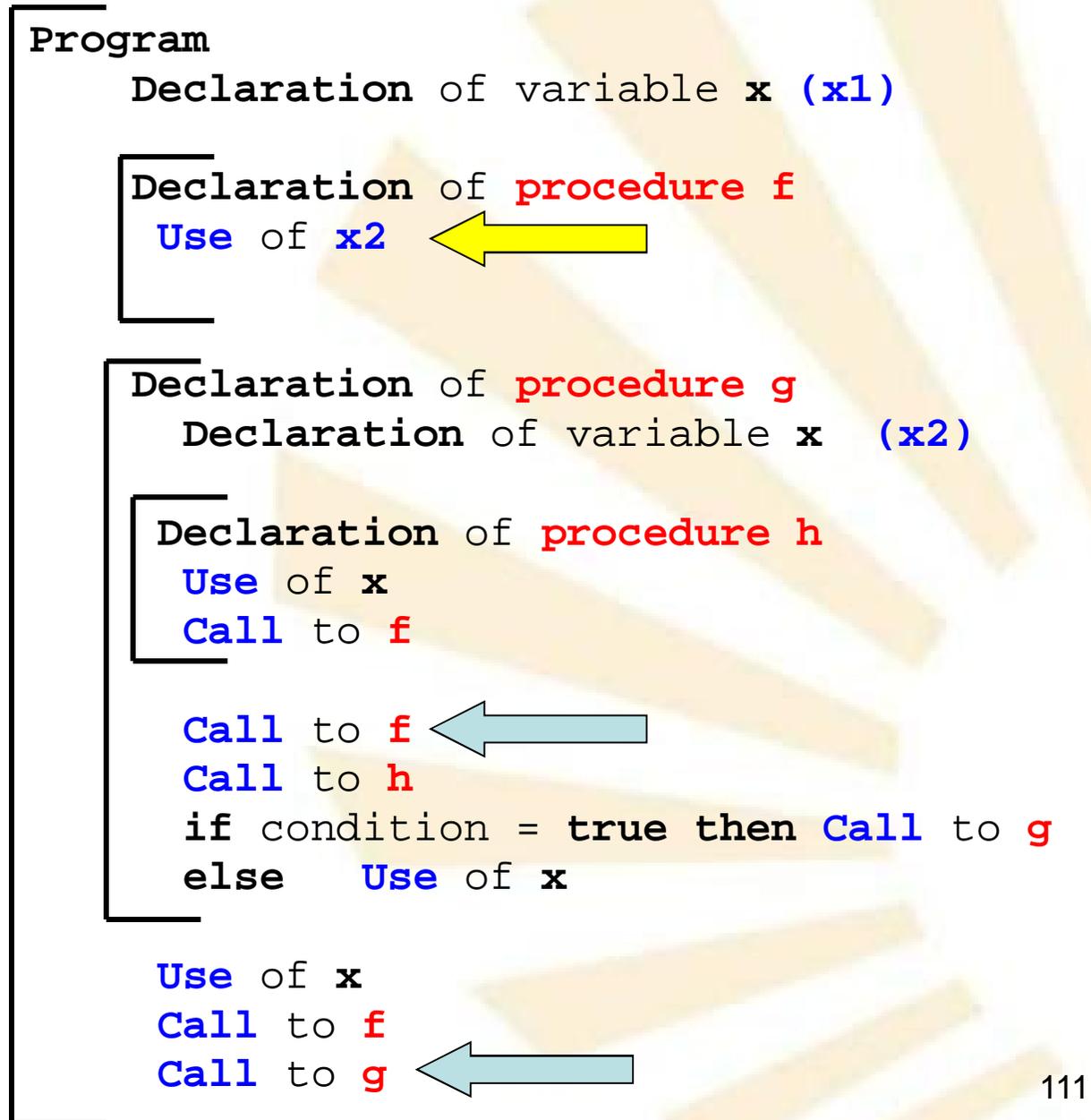




Run with  
dynamical  
scope

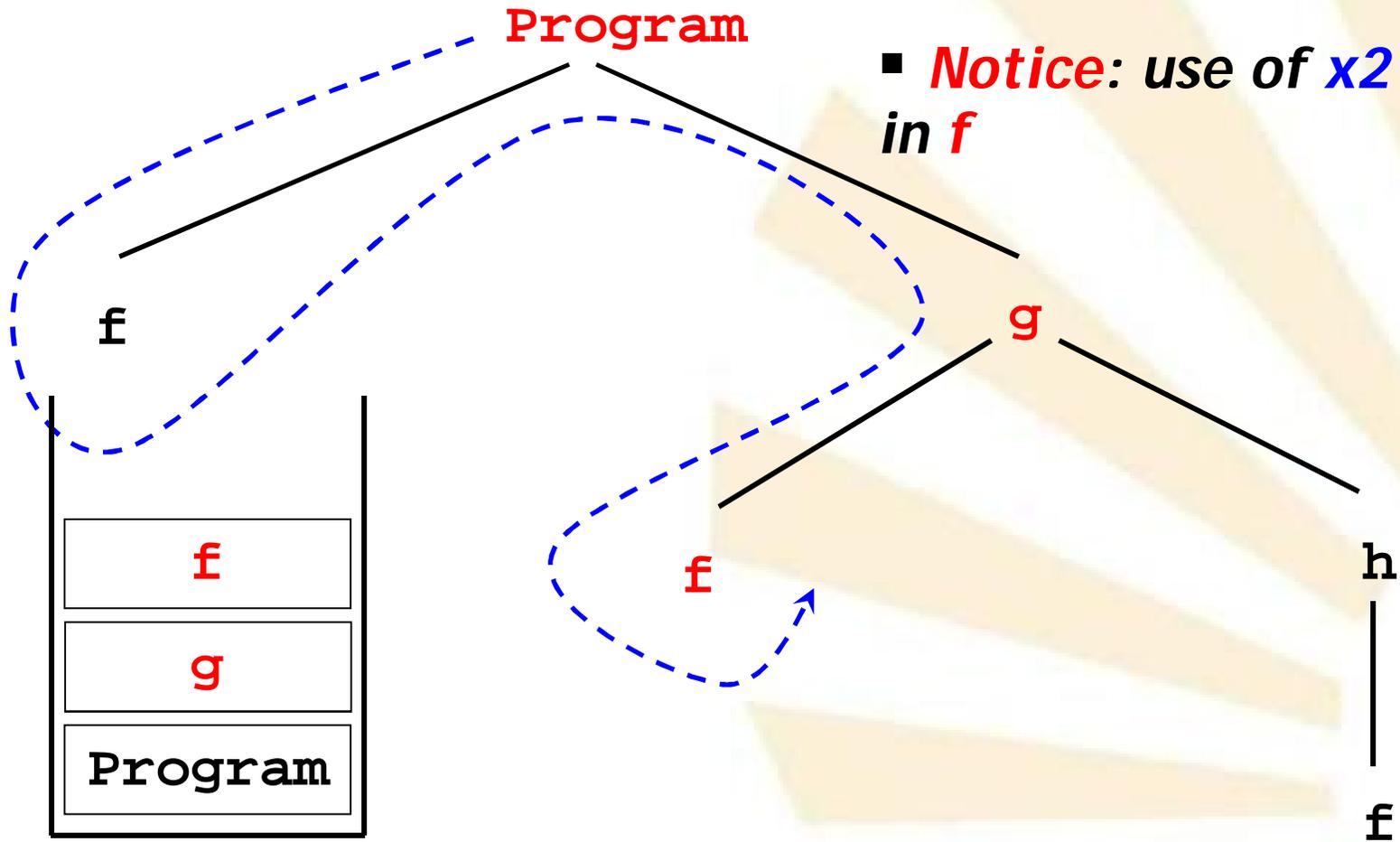


Run with  
dynamical  
scope



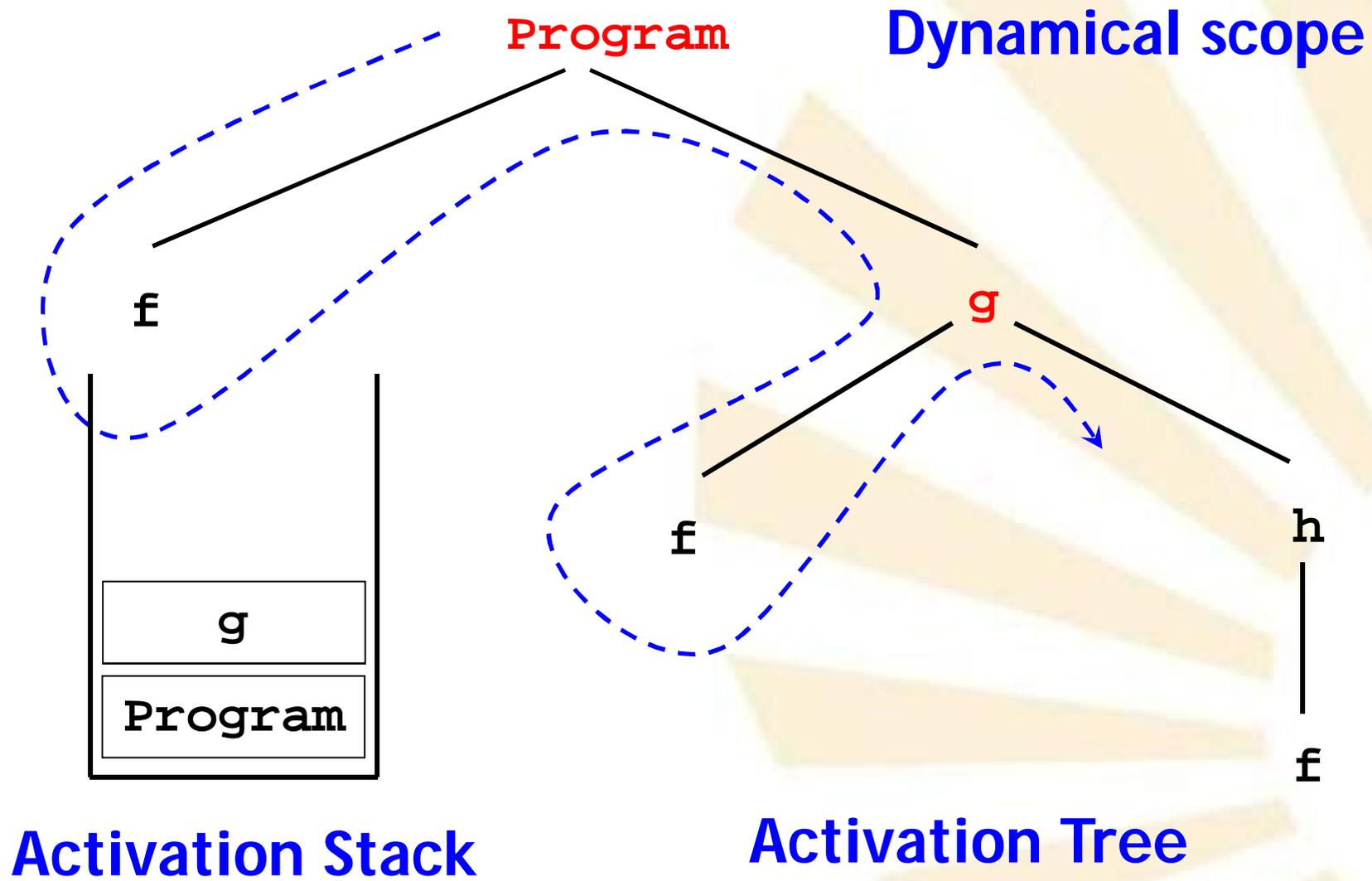
# Dynamical scope

- *Notice: use of  $x_2$  of  $g$  in  $f$*

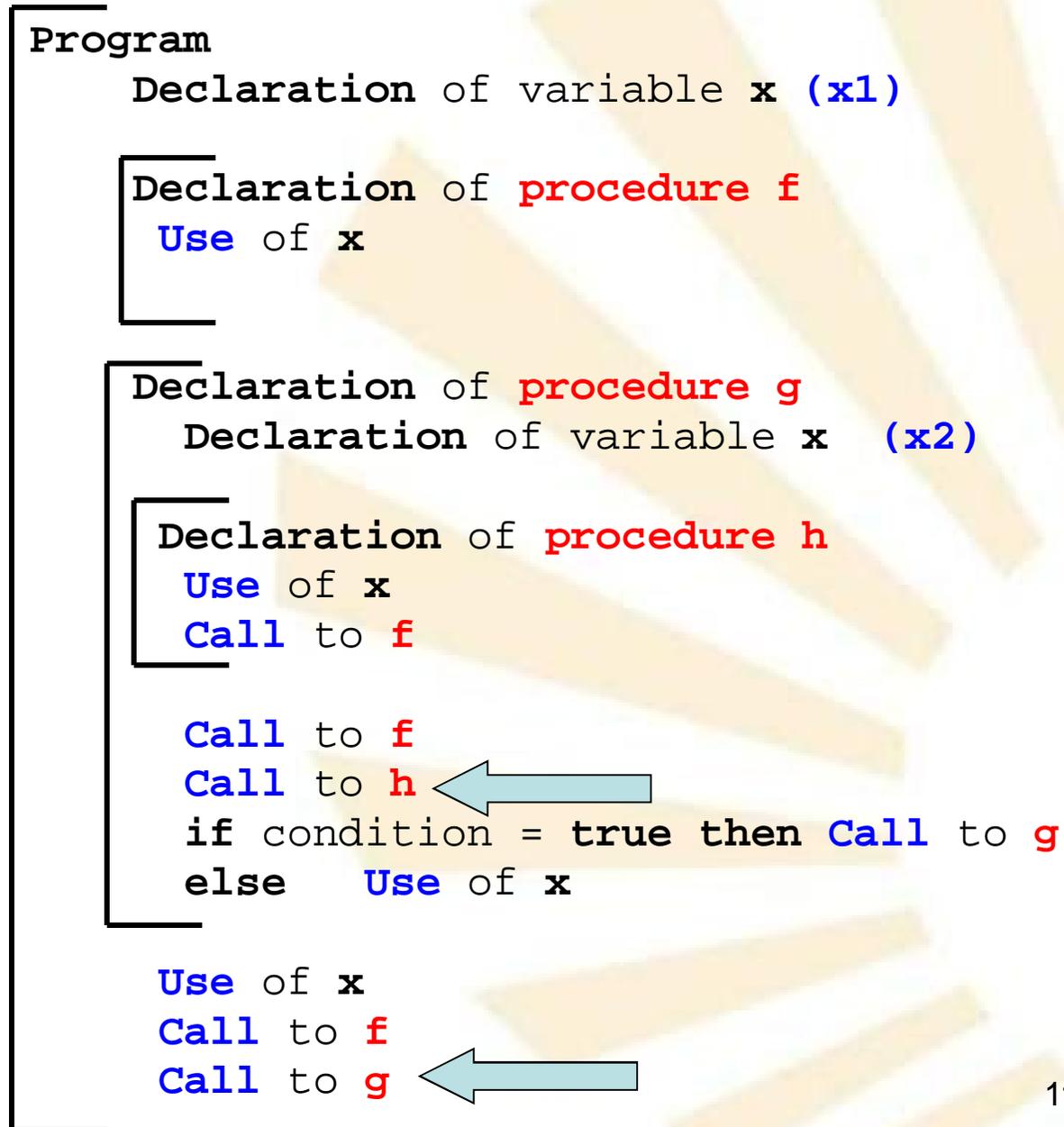


Activation Stack

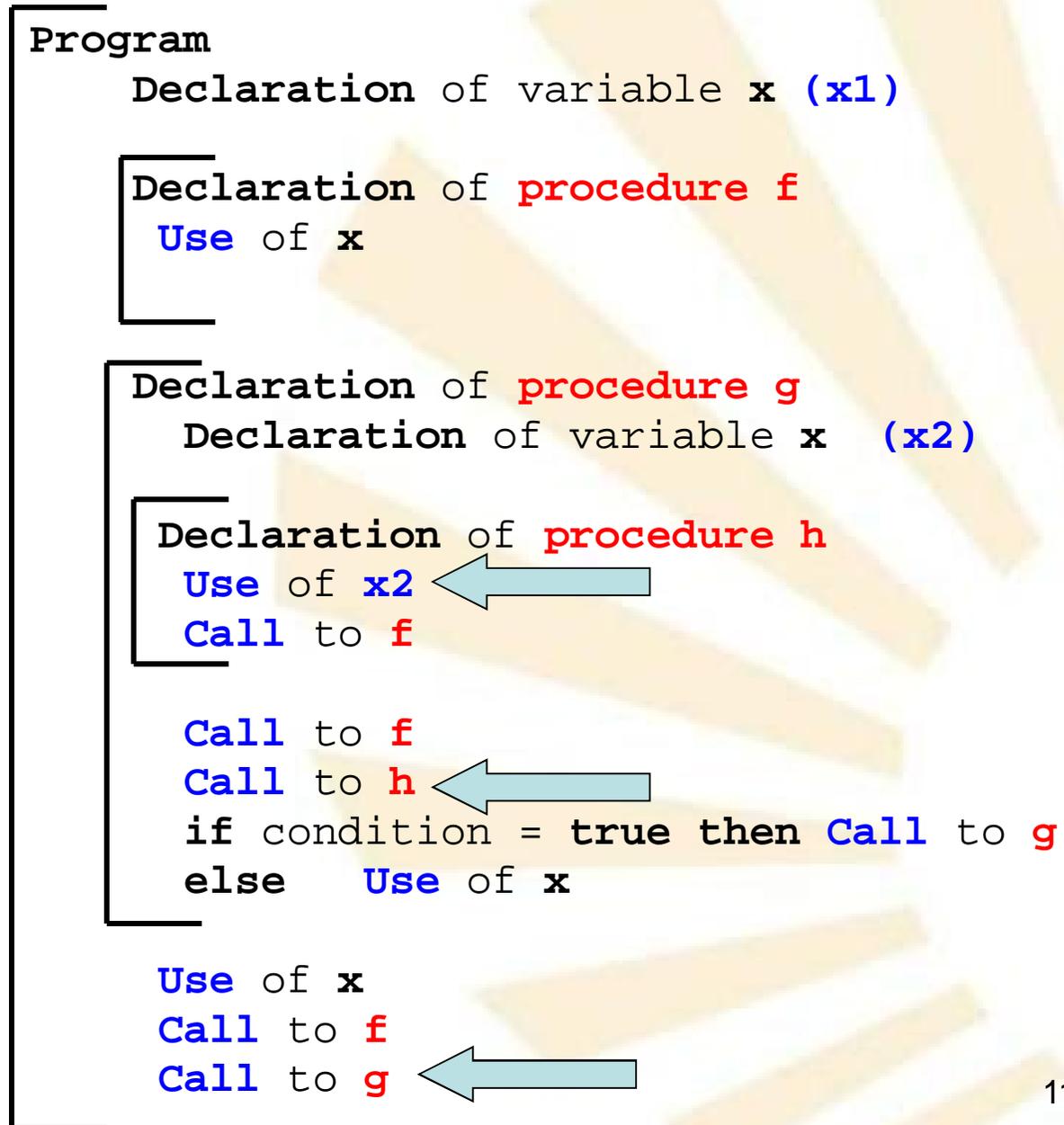
Activation Tree



Run with  
dynamical  
scope

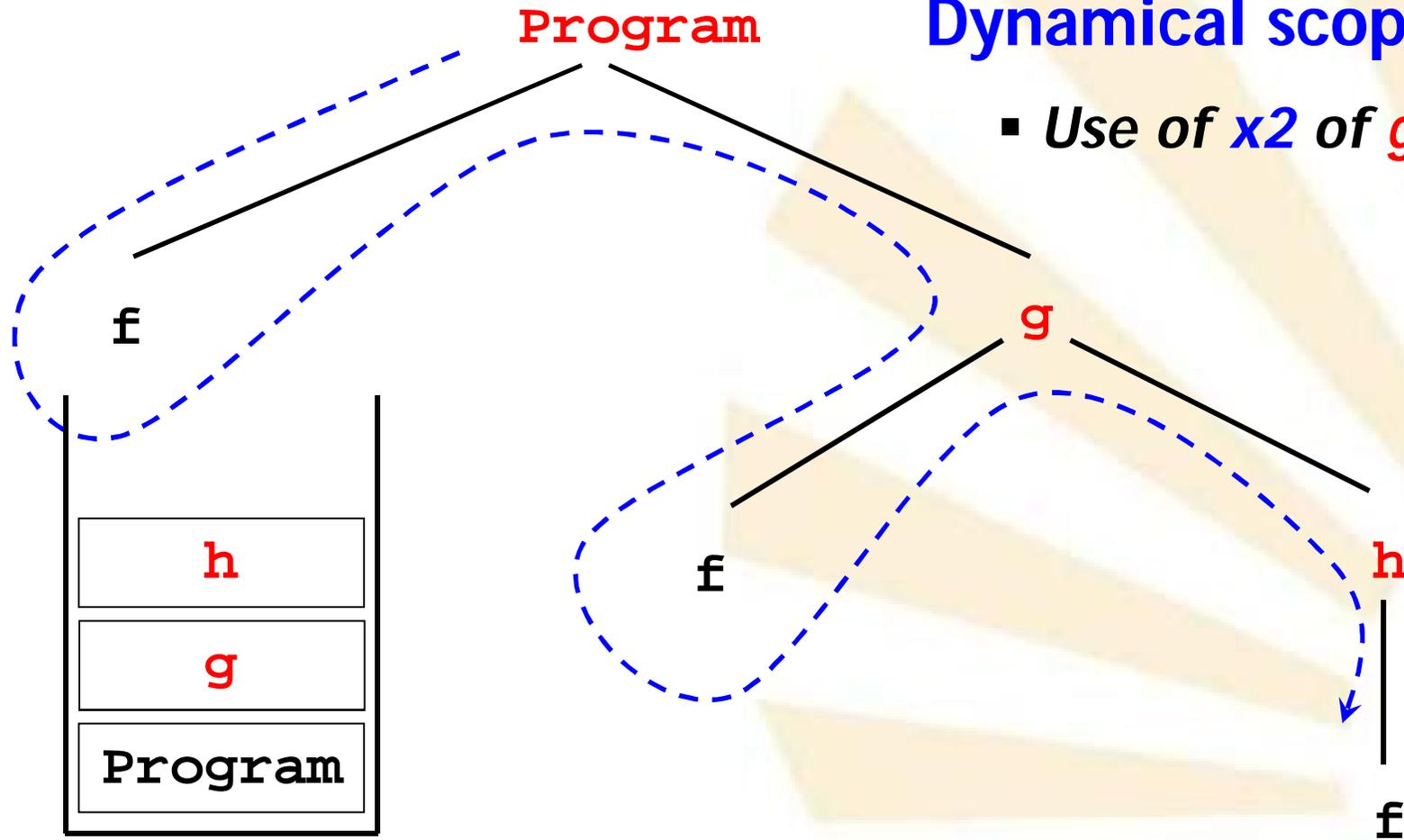


Run with  
dynamical  
scope



## Dynamical scope

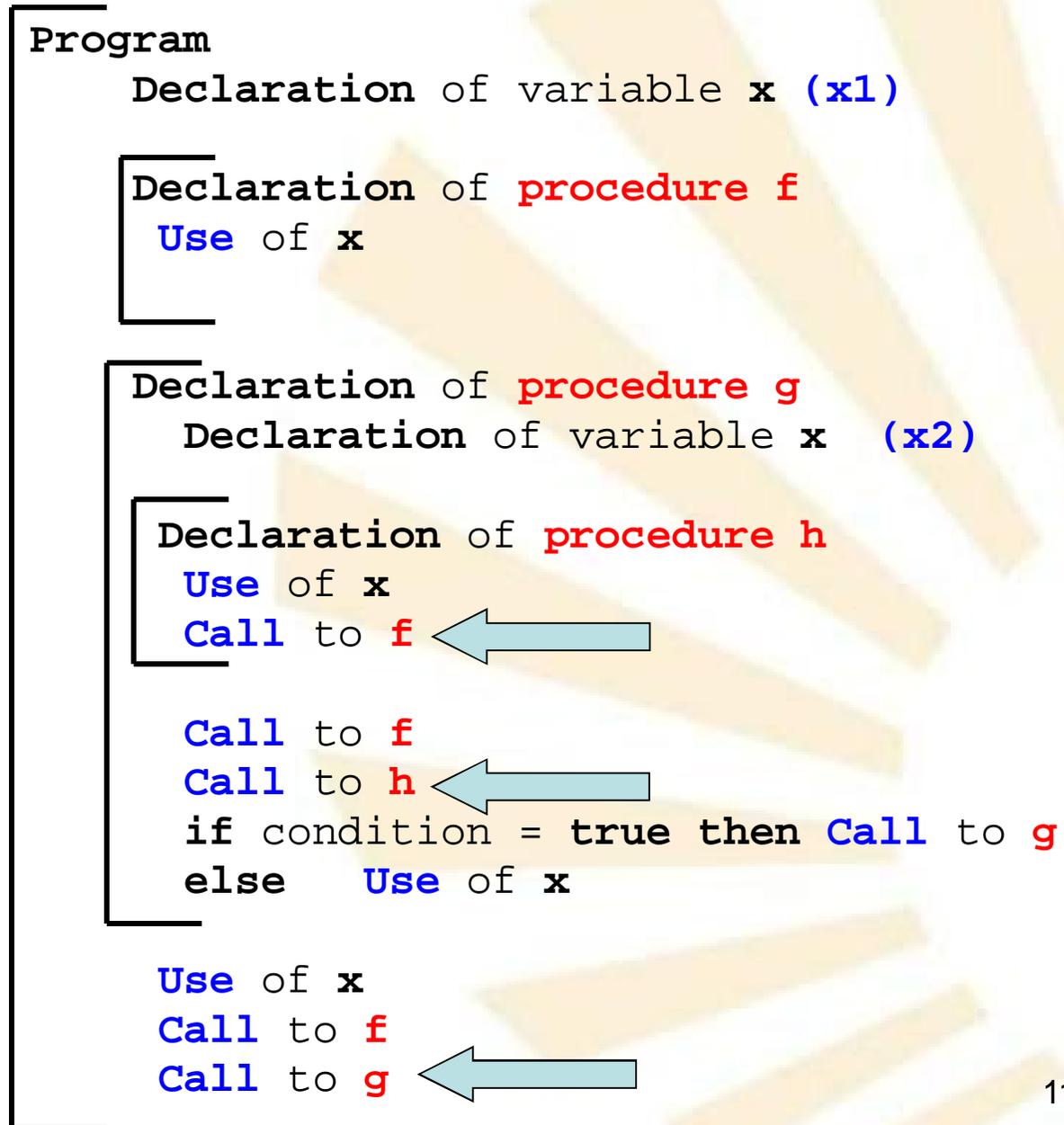
- Use of  $x_2$  of  $g$  in  $h$



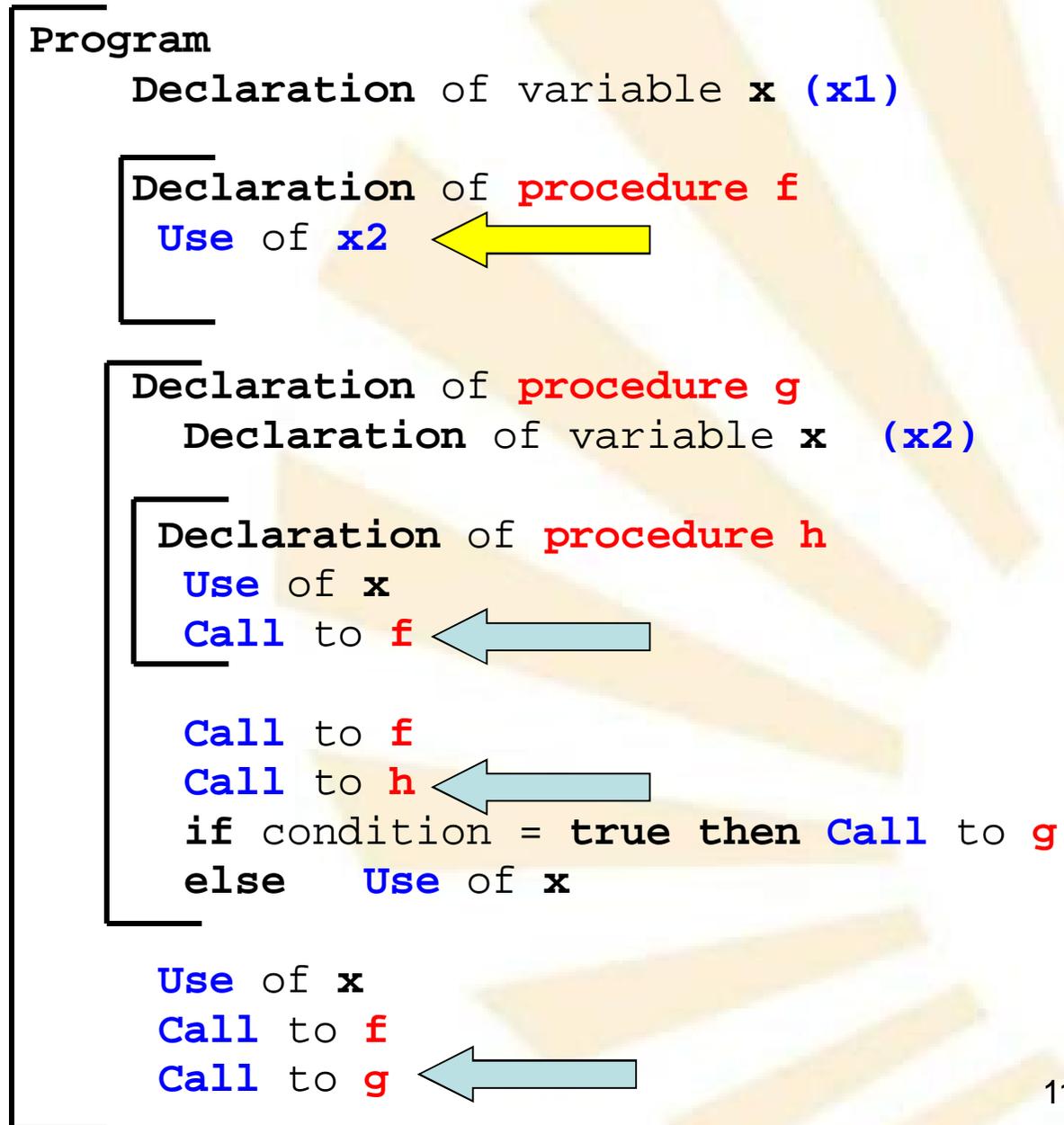
Activation Stack

Activation Tree

Run with  
dynamical  
scope

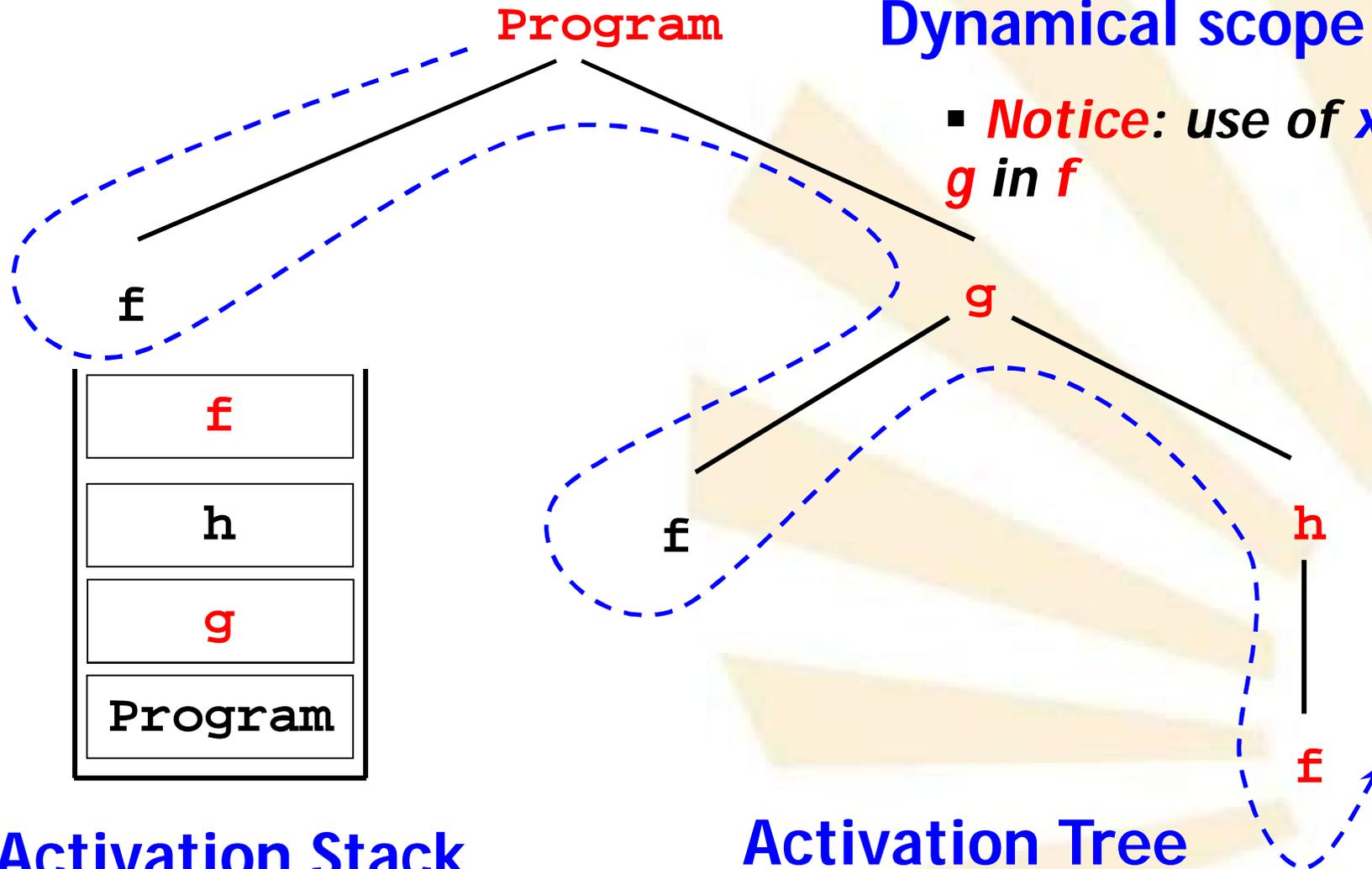


Run with  
dynamical  
scope



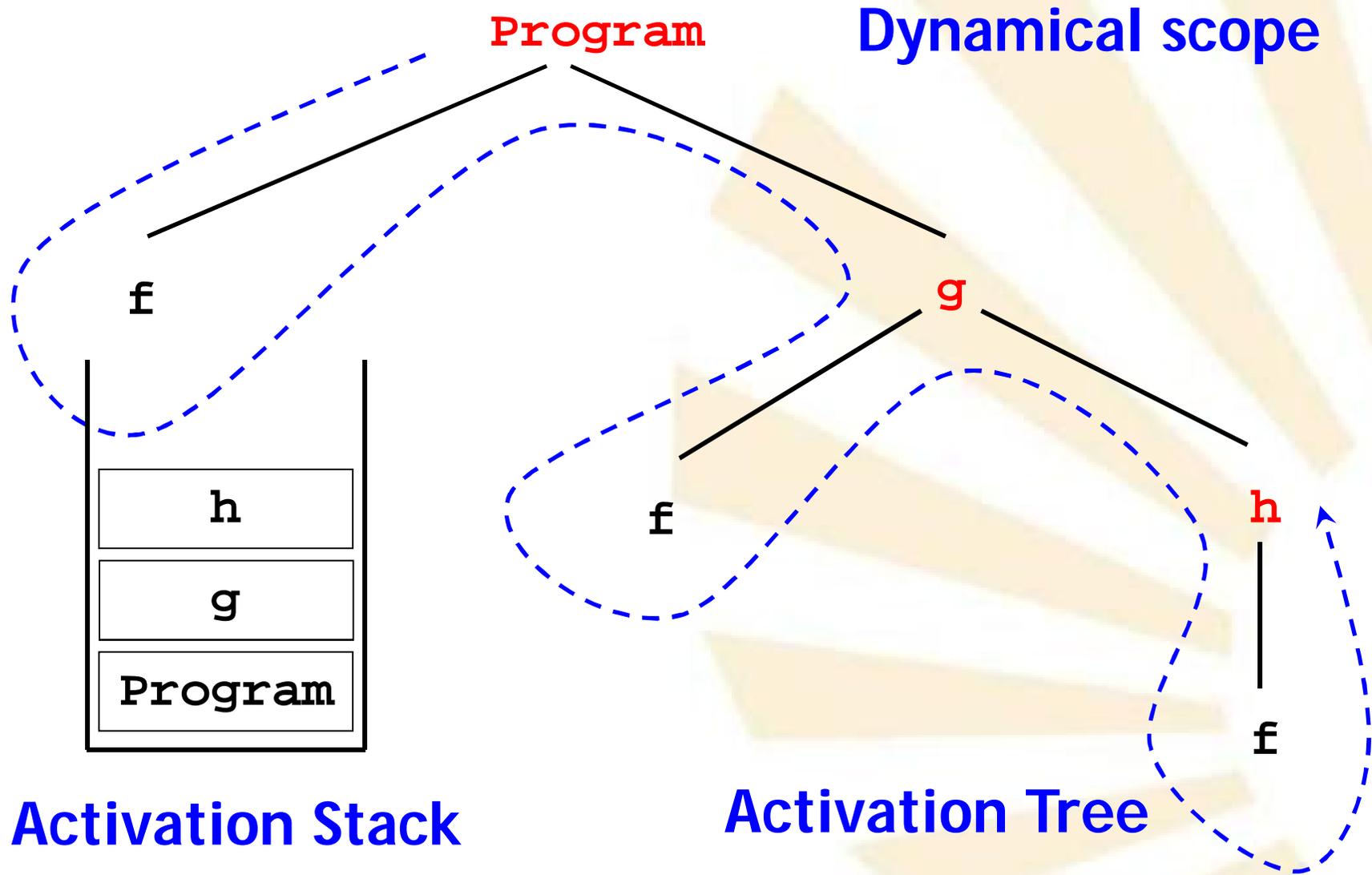
## Dynamical scope

- **Notice:** use of  $x_2$  of  $g$  in  $f$

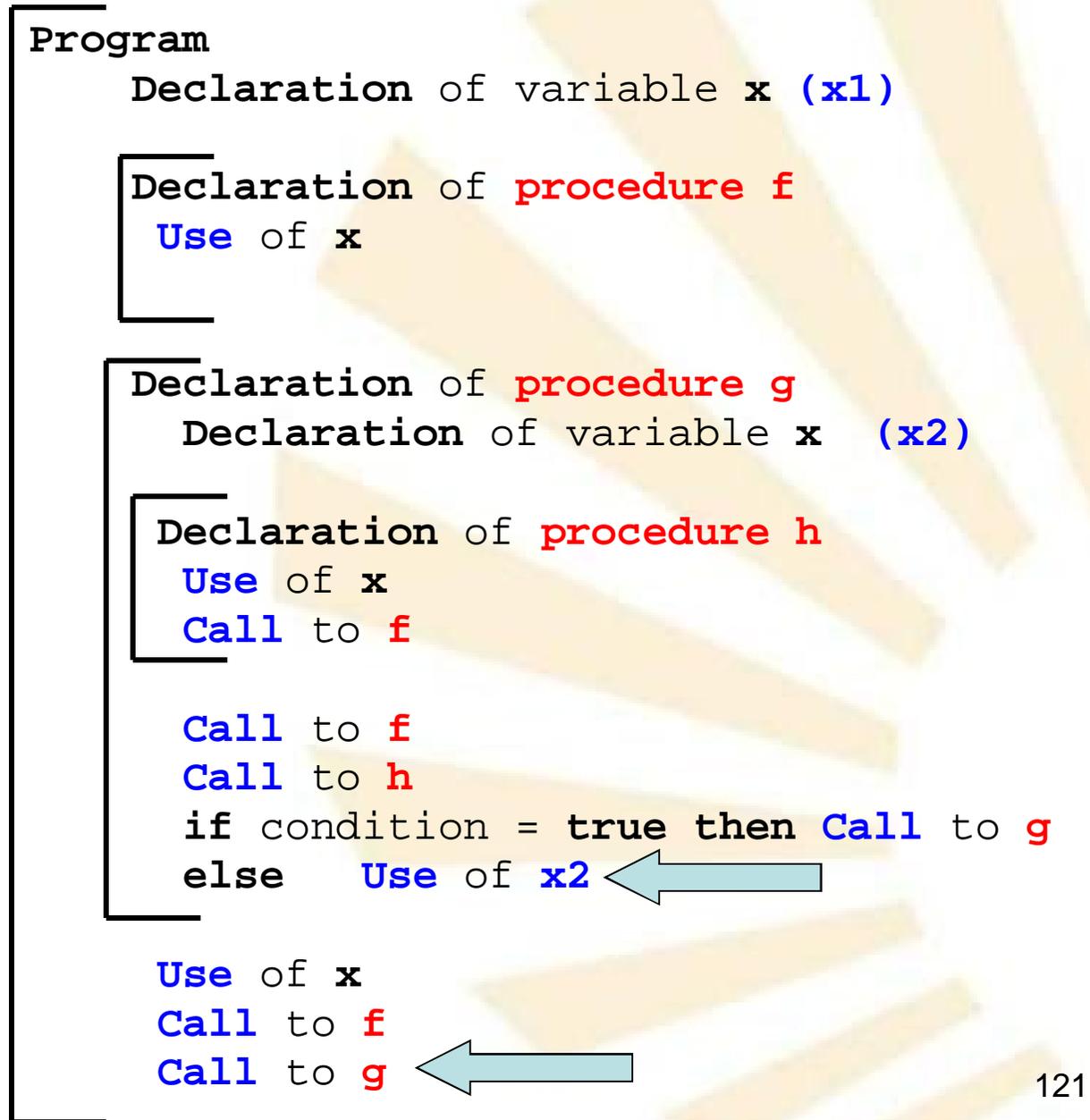


Activation Stack

Activation Tree

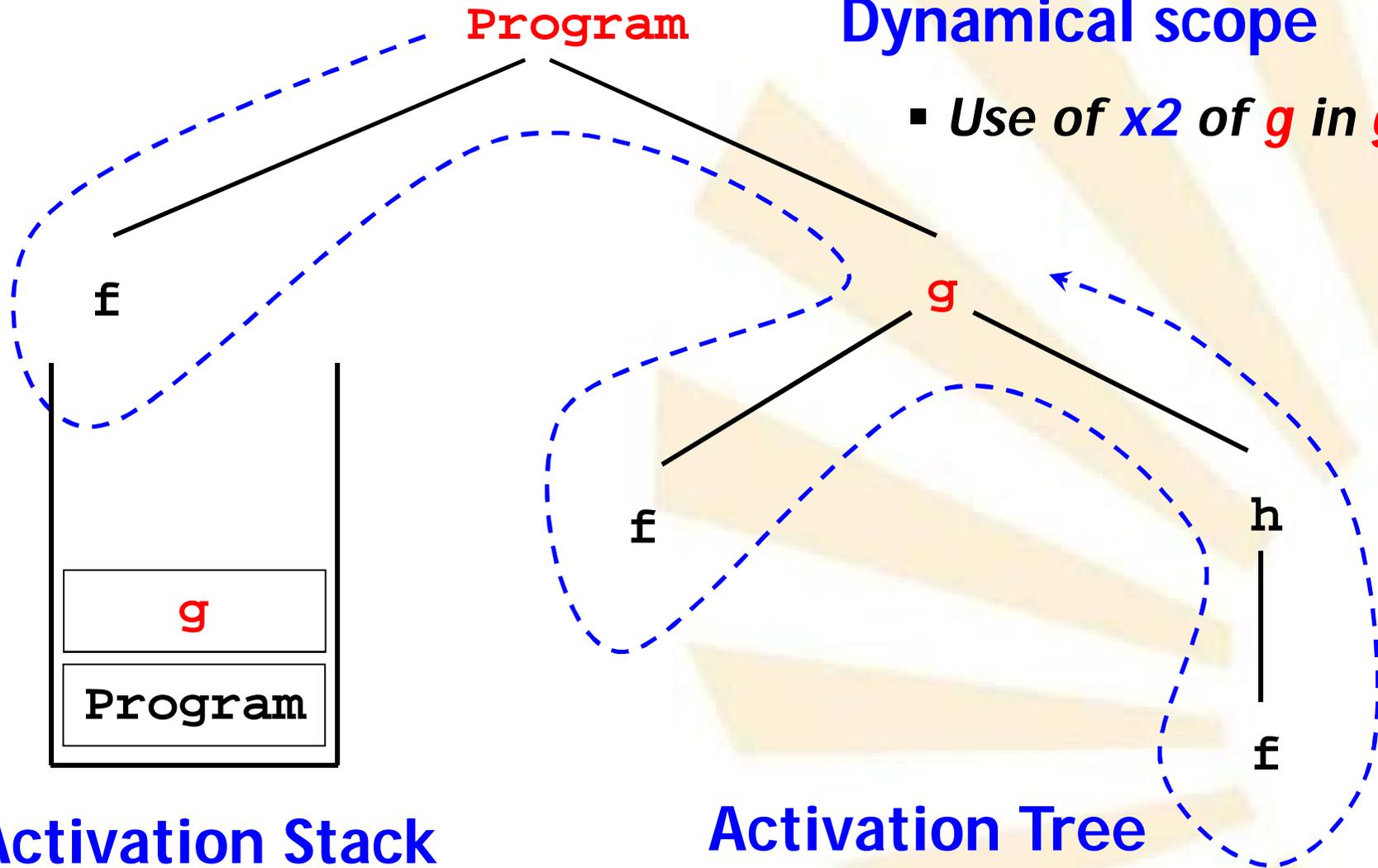


Run with  
dynamical  
scope



## Dynamical scope

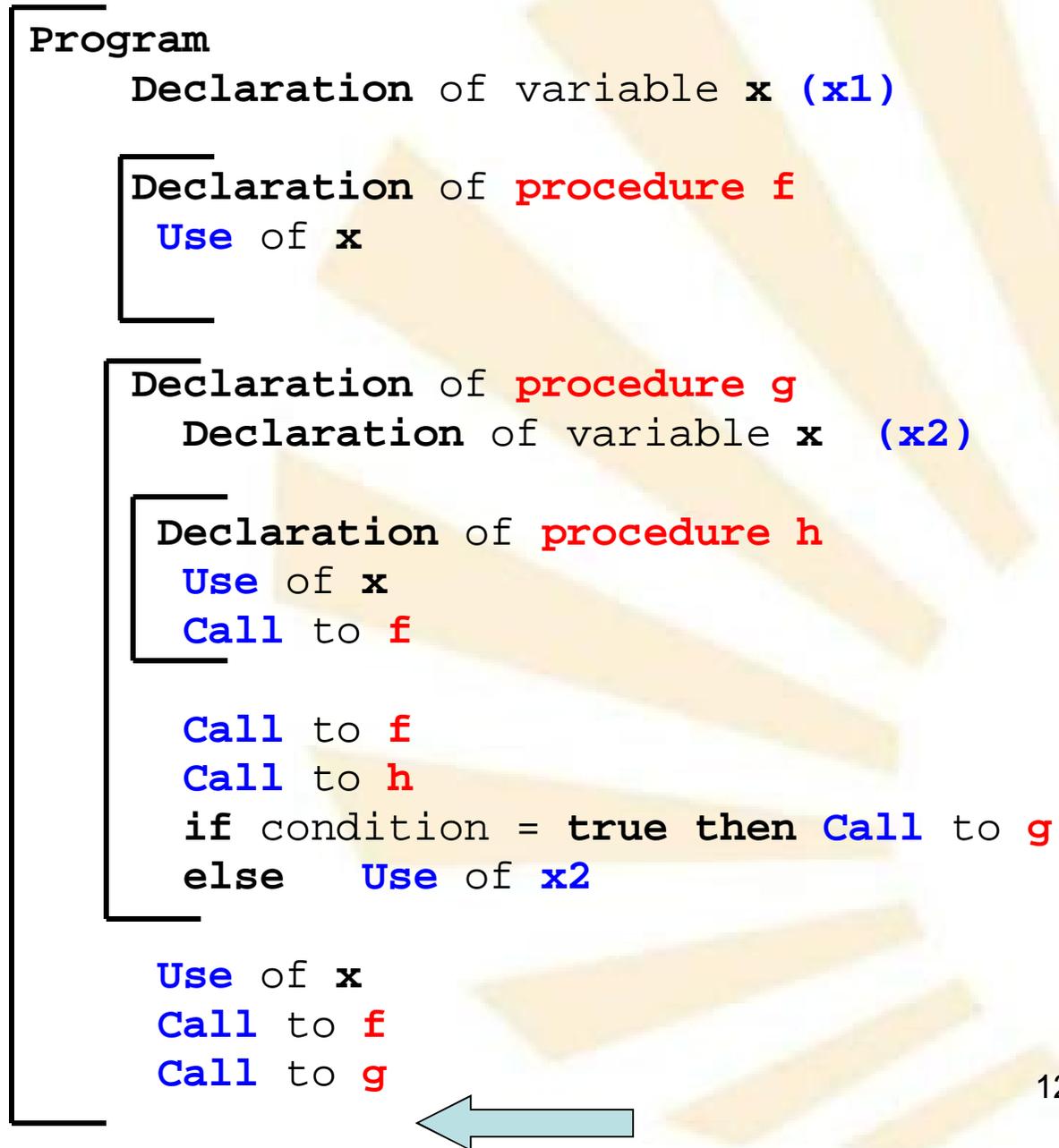
- Use of  $x_2$  of  $g$  in  $g$

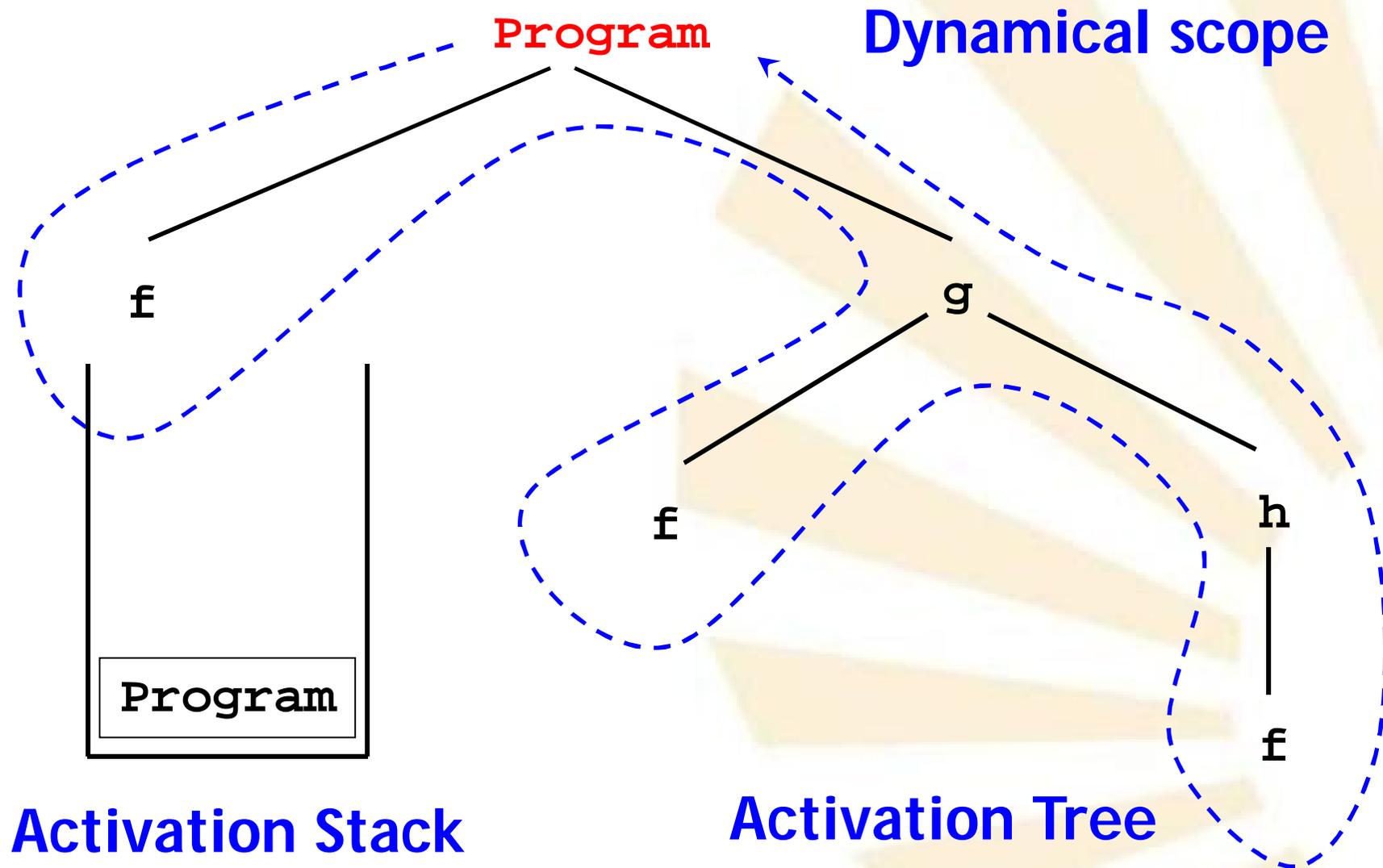


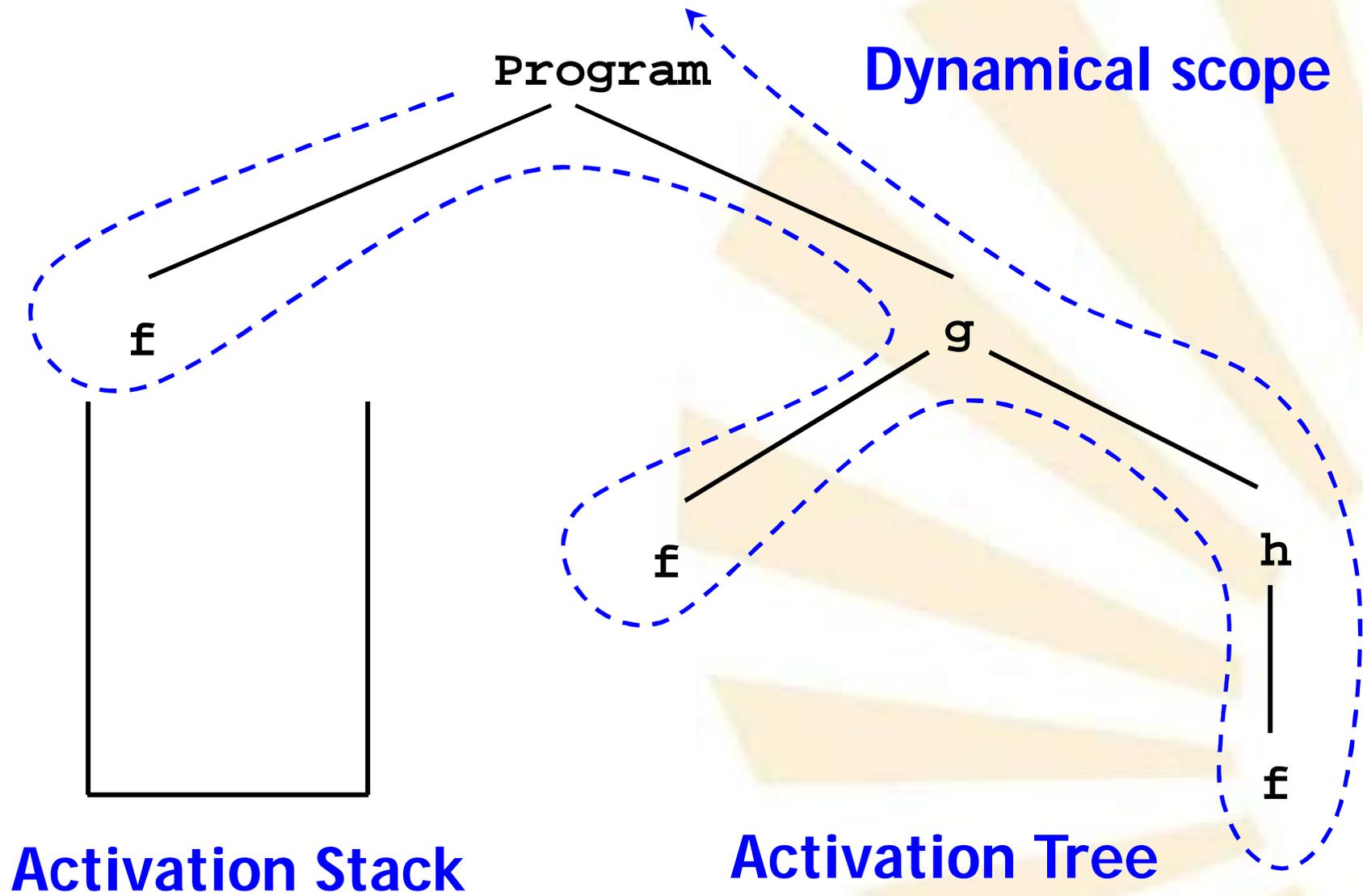
Activation Stack

Activation Tree

*Run with  
dynamical scope*







## 2. **Historic Summary of Scheme**

- ✓ LISP
- ✓ Compilation versus Interpretation
- ✓ Dynamically versus statically scope
- ✓ **Origin of Scheme**

## 2. Historic Summary of Scheme

### ✓ Origin of Scheme:

➤ Gerald Jay **Sussman** (MIT) and Guy Lewis **Steele** Jr.

➤ **Question:**

How would **LISP** be with **lexical** or **static scope** rules?

➤ **Answer:** new language → **Scheme**

▪ More **efficient** implementation of **recursion**

▪ **First class functions.**

▪ Rigorous **semantic** rules

➤ **Influence** on Common LISP: lexical scope rules

➤ ***Revised*<sup>5</sup> *Report on the Algorithmic Language Scheme***

## 2. Historic Summary of Scheme

✓ **Scheme:**

➤ **Structure of scheme programs**

- Sequence of
  - **definitions** of functions and variables
  - and **expressions**



**CÓRDOBA UNIVERSITY**

**SUPERIOR POLYTECHNIC SCHOOL**

**DEPARTMENT OF  
COMPUTER SCIENCE AND NUMERICAL ANALYSIS**



# **DECLARATIVE PROGRAMMING**

**COMPUTER ENGINEERING  
COMPUTATION ESPECIALITY**

**FOURTH YEAR**

**FIRST FOUR-MONTH PERIOD**



**Subject 1.- Introduction to Scheme language**